

OPTIMAL PACKINGS OF HAMILTON CYCLES IN GRAPHS OF HIGH MINIMUM DEGREE

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ABSTRACT. We study the number of edge-disjoint Hamilton cycles one can guarantee in a sufficiently large graph G on n vertices with minimum degree $\delta = (1/2 + \alpha)n$. For any constant $\alpha > 0$, we give an optimal answer in the following sense: let $\text{reg}_{\text{even}}(n, \delta)$ denote the degree of the largest even-regular spanning subgraph one can guarantee in a graph on n vertices with minimum degree δ . Then the number of edge-disjoint Hamilton cycles we find equals $\text{reg}_{\text{even}}(n, \delta)/2$. The value of $\text{reg}_{\text{even}}(n, \delta)$ is known for infinitely many values of n and δ . We also extend our results to graphs G of minimum degree $\delta \geq n/2$, unless G is close to the extremal constructions for Dirac's theorem. Our proof relies on a recent and very general result of Kühn and Osthus on Hamilton decomposition of robustly expanding regular graphs.

1. INTRODUCTION

Dirac's theorem [3] states that any graph on $n \geq 3$ vertices with minimum degree at least $n/2$ contains a Hamilton cycle. This degree condition is best possible. Surprisingly, though, the assertion of Dirac's theorem can be strengthened considerably: Nash-Williams [14] proved that the conditions of Dirac's theorem actually guarantee linearly many edge-disjoint Hamilton cycles.

Theorem 1. *Every graph on n vertices with minimum degree at least $n/2$ contains at least $\lfloor 5n/224 \rfloor$ edge-disjoint Hamilton cycles.*

Nash-Williams [15] initially conjectured that such a graph must contain at least $\lfloor n/4 \rfloor$ edge-disjoint Hamilton cycles, which would clearly be best possible. However, Babai observed that this trivial bound is very far from the truth (see [15]). Indeed, the following construction (which is based on Babai's argument) gives a graph G which contains at most $\lfloor (n+2)/8 \rfloor$ edge-disjoint Hamilton cycles. The graph G consists of one empty vertex class A of size $2m$, one vertex class B of size $2m+2$ containing a perfect matching and no other edges, and all possible edges between A and B . Thus G has order $n = 4m+2$ and minimum degree $2m+1$. Any Hamilton cycle in G must contain at least two edges of the perfect matching in B , so G contains at most $\lfloor (m+1)/2 \rfloor$ edge-disjoint Hamilton cycles.

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The above question of Nash-Williams naturally extends to graphs of higher minimum degree: suppose that $n/2 \leq \delta \leq n-1$. *How many edge-disjoint Hamilton cycles can one guarantee in a graph G on n vertices with minimum degree δ ?*

Clearly, as δ increases, one expects to find more edge-disjoint Hamilton cycles. However, the above construction shows that the trivial bound of $\lfloor \delta/2 \rfloor$ cannot always be attained. A less trivial bound is provided by the largest even-regular spanning subgraph in G . More precisely, let $\text{reg}_{\text{even}}(G)$ be the largest degree of an even-regular spanning subgraph of G . Then let

$$\text{reg}_{\text{even}}(n, \delta) := \min\{\text{reg}_{\text{even}}(G) : |G| = n, \delta(G) = \delta\}.$$

Clearly, in general we cannot guarantee more than $\text{reg}_{\text{even}}(n, \delta)/2$ edge-disjoint Hamilton cycles in a graph of order n and minimum degree δ . In fact, we conjecture this bound can always be attained.

Conjecture 2. *Suppose G is a graph on n vertices with minimum degree $\delta \geq n/2$. Then G contains at least $\text{reg}_{\text{even}}(n, \delta)/2$ edge-disjoint Hamilton cycles.*

Our main result confirms this conjecture exactly, as long as n is large and δ is slightly larger than $n/2$.

Theorem 3. *For every $\varepsilon > 0$, there exists an integer $n_0 = n_0(\varepsilon)$ such that every graph G on $n \geq n_0$ vertices with $\delta(G) \geq (1/2 + \varepsilon)n$ contains at least $\text{reg}_{\text{even}}(n, \delta(G))/2$ edge-disjoint Hamilton cycles.*

In fact, we even show that if G is not close to the extremal example, then G contains significantly more than the required number of edge-disjoint Hamilton cycles (see Lemma 24). Our proof of Theorem 3 is based on a recent result (Theorem 9) of Kühn and Osthus [11, 12], which states that every “robustly expanding” regular (di)graph has a Hamilton decomposition. In [12], a straightforward argument was already used to derive Conjecture 2 for $\delta \geq (2 - \sqrt{2} + \varepsilon)n$ (see Section 3.2). Our extension of this result to $\delta \geq (1/2 + \varepsilon)n$ involves new ideas.

Subsequently, Csaba, Kühn, Lo, Osthus and Treglown [2] have proved Conjecture 2 for large n , by solving the case when δ is allowed to be close to $n/2$. The proof relies on Theorem 3 and Theorem 5. (The latter provides a stability result when δ is close to $n/2$.)

Earlier, Christofides, Kühn and Osthus [1] used the regularity lemma to prove an approximate version of Theorem 3. Hartke and Seacrest [5] were able to improve this result while avoiding the use of the regularity lemma (but still with the same restriction on δ). This enabled them to omit the condition that G has to be very large. They also gave significantly better error bounds.

Accurate bounds on $\text{reg}_{\text{even}}(n, \delta)$ are known. Note that the complete bipartite graph whose vertex classes are almost equal shows that $\text{reg}_{\text{even}}(n, \delta) = 0$ for $\delta < n/2$. Katerinis [8] considered the case when $\delta = n/2$. His result was independently generalised to larger values of δ in [1] (see [12] for a summarised version) and by Hartke, Martin and Seacrest [4]. The following bounds are from [4].

Theorem 4. *Suppose that $n, \delta \in \mathbb{N}$ and $n/2 \leq \delta < n$. Then*

$$(1) \quad \frac{\delta + \sqrt{n(2\delta - n) + 8}}{2} - \varepsilon \leq \text{reg}_{\text{even}}(n, \delta) \leq \frac{\delta + \sqrt{n(2\delta - n)}}{2} + \frac{4}{\sqrt{n(2\delta - n) + 4}}.$$

where $0 < \varepsilon \leq 2$ is chosen to make the left hand side of (1) an even integer.

Note that (1) always yields at most two possible values for $\text{reg}_{\text{even}}(n, \delta)$ and even determines it exactly for many values of the parameters n and δ . For example, (1) determines $\text{reg}_{\text{even}}(n, n/2)$ (e.g. in the case when n is divisible by 8 it is $n/4$). The bounds in [1] also give at most two possible values. The lower bound in (1) is based on Tutte's factor theorem [17]. The upper bound is obtained by a natural generalization of Babai's construction (see Section 3.1 for a description).

Our second result concerns the case of Conjecture 2 where we allow δ to be close to $n/2$. In this case, we obtain the following 'stability result': if $\delta(G) = (1/2 + o(1))n$, then Conjecture 2 holds for large n as long as G has suitable expansion properties. In this case, we even obtain significantly more than the required number of edge-disjoint Hamilton cycles again. These expansion properties fail only when G is very close to the extremal examples for Dirac's theorem.

Theorem 5. *For every $0 < \eta < 1/8$, there exist $\varepsilon > 0$ and an integer n_0 such that every graph G on $n \geq n_0$ vertices with $(1/2 - \varepsilon)n \leq \delta(G) \leq (1/2 + \varepsilon)n$ satisfies one of the following:*

- (i) *There exists $A \subseteq V(G)$ with $|A| = \lfloor n/2 \rfloor$ and such that $e(A) \leq \eta n^2$.*
- (ii) *There exists $A \subseteq V(G)$ with $|A| = \lfloor n/2 \rfloor$ and such that $e(A, \bar{A}) \leq \eta n^2$.*
- (iii) *G contains at least $\max\{\text{reg}_{\text{even}}(n, \delta(G))/2, n/8\} + \varepsilon n$ edge-disjoint Hamilton cycles.*

Note that if G satisfies (i) then $e(A, \bar{A})$ must be roughly $n^2/4$, i.e. G is close to $K_{n/2, n/2}$ with possibly some edges added to one of the vertex classes. If G satisfies (ii), then both $e(A)$ and $e(\bar{A})$ must be roughly $n^2/8$, i.e. G is close to the union of two equal-sized cliques.

Although Conjecture 2 is optimal for the class of graphs on n vertices and minimum degree δ , it will not be optimal for every graph in the class – some graphs G will contain far more than $\text{reg}_{\text{even}}(n, \delta)/2$ edge-disjoint Hamilton cycles. The following conjecture accounts for this and would be best possible for every single graph G . Note that it is far stronger than Conjecture 2.

Conjecture 6. *Suppose G is a graph on n vertices with minimum degree $\delta(G) \geq n/2$. Then G contains at least $\text{reg}_{\text{even}}(G)/2$ edge-disjoint Hamilton cycles.*

For $\delta \geq (2 - \sqrt{2} + \varepsilon)n$, this conjecture was proved in [12], based on the main result of [11]. It would already be very interesting to obtain an approximate version of Conjecture 6, i.e. a set of $(1 - \varepsilon)\text{reg}_{\text{even}}(G)/2$ edge-disjoint Hamilton cycles under the assumption that $\delta(G) \geq (1 + \varepsilon)n/2$.

As a very special case, Conjecture 6 would imply the long-standing 'Hamilton factorization' conjecture of Nash-Williams [15, 16]: any d -regular graph on at most $2d$ vertices contains $\lfloor d/2 \rfloor$ edge-disjoint Hamilton cycles. Jackson [16] raised the

same conjecture independently, and proved a partial result. This was improved to an approximate version of the conjecture in [1]. The best current result towards the Hamilton factorization conjecture is due to Kühn and Osthus [12] (again as a corollary of their main result in [11]). Note that the set of Hamilton cycles guaranteed by Theorem 7 actually forms a Hamilton decomposition.

Theorem 7. *For every $\varepsilon > 0$ there exists an integer n_0 such that every d -regular graph on $n \geq n_0$ vertices for which $d \geq (1/2 + \varepsilon)n$ is even contains $d/2$ edge-disjoint Hamilton cycles.*

Frieze and Krivelevich conjectured that the trivial bound of $\lfloor \delta(G)/2 \rfloor$ edge-disjoint Hamilton cycles is in fact correct for random graphs. Indeed, the results of several authors (mainly Krivelevich and Samotij [10] as well as Knox, Kühn and Osthus [9]) can be combined to show that for all $0 \leq p \leq 1$, the binomial random graph $G_{n,p}$ contains $\lfloor \delta(G_{n,p})/2 \rfloor$ edge-disjoint Hamilton cycles with high probability. Some further related results can be found in [6, 11, 12].

2. NOTATION

Given a graph G , we write $V(G)$ for its vertex set, $E(G)$ for its edge set, $e(G) := |E(G)|$ for the number of its edges and $|G|$ for the number of its vertices. Given $X \subseteq V(G)$, we write $G - X$ for the graph formed by deleting all vertices in X and $G[X]$ for the subgraph of G induced by X . We will also write $\bar{X} := V(G) \setminus X$ when it is unambiguous to do so. Given disjoint sets $X, Y \subseteq V(G)$, we write $G[X, Y]$ for the bipartite subgraph induced by X and Y . If G and G' are two graphs, we write $G \dot{\cup} G'$ for the graph on $V(G) \dot{\cup} V(G')$ with edge set $E(G) \dot{\cup} E(G')$. If $V(G) = V(G')$, we also write $G + G'$ for the graph on $V(G)$ with edge set $E(G) \cup E(G')$. An r -factor of a graph G is a spanning r -regular subgraph of G . If H is an r -factor of G and r is even then we also call H an *even factor* of G .

If G is an undirected graph, we write $\delta(G)$ for the minimum degree of G , $\Delta(G)$ for the maximum degree of G and $d(G)$ for the average degree of G . Whenever $X, Y \subseteq V(G)$, we write $e_G(X, Y)$ for the number of all those edges which have one endvertex in X and the other in Y . We write $e_G(X)$ for the number of edges in $G[X]$, and $e'_G(X, Y) := e_G(X, Y) + e_G(X \cap Y)$. Thus $e'_G(X, Y)$ is the number of ordered pairs (x, y) of vertices such that $x \in X$, $y \in Y$ and $xy \in E(G)$. Given a vertex x of G , we write $d_G(x)$ for the degree of x in G . We often omit the subscript G if this is unambiguous. Also, if $A \subseteq V(G)$ and the graph G is clear from the context, we sometimes write $d_A(x)$ for the number of neighbours of x in A . If G is a digraph, we write $\delta^+(G)$ for the minimum outdegree of G and $\delta^-(G)$ for the minimum indegree of G .

In order to simplify the presentation, we omit floors and ceilings and treat large numbers as integers whenever this does not affect the argument. The constants in the hierarchies used to state our results have to be chosen from right to left. More precisely, if we claim that a result holds whenever $0 < 1/n \ll a \ll b \ll c \leq 1$ (where n is the order of the graph or digraph), then this means that there are non-decreasing functions $f : (0, 1] \rightarrow (0, 1]$, $g : (0, 1] \rightarrow (0, 1]$ and $h : (0, 1] \rightarrow (0, 1]$ such that the result holds for all $0 < a, b, c \leq 1$ and all $n \in \mathbb{N}$ with $b \leq f(c)$, $a \leq g(b)$ and $1/n \leq h(a)$.

We will not calculate these functions explicitly. Hierarchies with more constants are defined in a similar way.

Whenever $x \in \mathbb{R}$ we shall write $x_+ := \max\{x, 0\}$. We will write $a = x \pm \varepsilon$ as shorthand for $x - \varepsilon \leq a \leq x + \varepsilon$, and $a \neq x \pm \varepsilon$ as shorthand for the statement that either $a < x - \varepsilon$ or $a > x + \varepsilon$.

3. PROOF OUTLINE AND FURTHER NOTATION

3.1. The extremal graph. We start by defining a graph $G_{n,\delta,\text{ext}}$ on n vertices which is extremal for Theorem 4 in the sense that $G_{n,\delta,\text{ext}}$ has minimum degree δ but the largest degree of an even factor of $G_{n,\delta,\text{ext}}$ is at most the right hand side of (1). Given $\delta > n/2$, let Δ be the smallest integer such that $\Delta(\delta + \Delta - n)$ is even and $\Delta \geq (n + \sqrt{n(2\delta - n)})/2$. Partition the vertex set of $G_{n,\delta,\text{ext}}$ into two classes A and B , with $|B| = \Delta$ and $|A| = n - \Delta$. Let $G_{n,\delta,\text{ext}}[A]$ be empty, let $G_{n,\delta,\text{ext}}[B]$ be any $(\delta + \Delta - n)$ -regular graph, and let $G_{n,\delta,\text{ext}}[A, B]$ be the complete bipartite graph. Clearly $\delta(G_{n,\delta,\text{ext}}) = \delta$. Moreover, if H is a factor of $G_{n,\delta,\text{ext}}$, then one can show that $d(H)$ is at most the right hand side of (1) (see [4] for details). In particular, $G_{n,\delta,\text{ext}}$ contains at most $d(H)/2$ Hamilton cycles. Essentially the same construction was given in [1].

3.2. Tools and proof overview. An important concept in our proofs of Theorems 3 and 5 will be the notion of robust expanders. This concept was first introduced by Kühn, Osthus and Treglown [13] for directed graphs. Roughly speaking, a graph is a robust expander if for every set S which is not too small and not too large, its “robust” neighbourhood is at least a little larger than S .

Definition 8. Let G be a graph on n vertices. Given $0 < \nu \leq \tau < 1$ and $S \subseteq V(G)$, we define the ν -robust neighbourhood $RN_{\nu,G}(S)$ of S to be the set of all vertices $v \in V(G)$ with $d_S(v) \geq \nu n$. We say that G is a *robust* (ν, τ) -*expander* if for all $S \subseteq V(G)$ with $\tau n \leq |S| \leq (1 - \tau)n$, we have $|RN_{\nu,G}(S)| \geq |S| + \nu n$.

The main tool for our proofs is the following result of Kühn and Osthus [11] which states that every even-regular robust expander G whose degree is linear in $|G|$ has a Hamilton decomposition.

Theorem 9. *For every $\alpha > 0$, there exists $\tau > 0$ such that for every $\nu > 0$, there exists $n_0(\alpha, \tau, \nu)$ such that the following holds. Suppose that*

- (i) G is an r -regular graph on $n \geq n_0$ vertices, where $r \geq \alpha n$ and r is even;
- (ii) G is a robust (ν, τ) -expander.

Then G has a Hamilton decomposition.

Let G be a graph on n vertices as in Theorem 3. Let $\delta := \delta(G) = (1/2 + \alpha)n$. (So $\alpha \geq \varepsilon$.) As observed in [12], every graph on n vertices whose minimum degree is at least slightly larger than $n/2$ is a robust expander (see Lemma 17). Thus our given graph G is a robust expander. Let G^* be an even factor of largest degree in G . So $d(G^*) \geq \text{reg}_{\text{even}}(n, \delta)$. If G^* is still a robust expander, then we can apply Theorem 9 to obtain a Hamilton decomposition of G^* and thus at least $\text{reg}_{\text{even}}(n, \delta)/2$ edge-disjoint Hamilton cycles in G . The problem is that if α is small, then we could

have $d(G^*) \leq n/2$. So we cannot guarantee that G^* is a robust expander. (However, this approach works if α is at least slightly larger than $3/2 - \sqrt{2}$. Indeed, in this case Theorem 4 implies that $d(G^*)$ will be slightly larger than $n/2$ and so G^* will be a robust expander. This observation was used in [12] to prove Theorem 3 for such values of α .)

So instead of using this simple strategy, in the proof of Theorem 3 we will distinguish two cases depending on whether our graph G contains a subgraph which is close to $G_{n,\delta,\text{ext}}$. Suppose first that G contains such a subgraph, G_1 say. We can choose G_1 in such a way that $\delta(G_1) = \delta$, so G_1 must have an even factor G_2 of degree at least $\text{reg}_{\text{even}}(n, \delta)$. We will then use the fact that G_1 is close to $G_{n,\delta,\text{ext}}$ in order to prove directly that G_2 is a robust expander. As before, this yields a Hamilton decomposition of G_2 by Theorem 9. This part of the argument is contained in Section 4.

If G does not contain a subgraph close to $G_{n,\delta,\text{ext}}$, then we will first find a sparse even factor H of G which is still a robust expander and remove it from G . Call the resulting graph G' . We will then use the fact that G is far from containing $G_{n,\delta,\text{ext}}$ to show that G' still contains an even factor H' of degree at least $\text{reg}_{\text{even}}(n, \delta)$. Since robust expansion is a monotone property, it follows that $H + H'$ is still a robust expander and may therefore be decomposed into Hamilton cycles by Theorem 9. So in this case we even find slightly more than $\text{reg}_{\text{even}}(n, \delta)/2$ edge-disjoint Hamilton cycles. This part of the argument is contained in Section 5. Altogether this will then imply Theorem 3.

In order to prove Theorem 5, we first show that every graph G whose minimum degree is close to $n/2$ either satisfies conditions (i) and (ii) of Theorem 5 or is a robust expander which does not contain a subgraph close to $G_{n,\delta,\text{ext}}$. So suppose G does not satisfy (i) and (ii). We will use the fact that G is a robust expander to find a sparse robustly expanding even factor of G , and then argue similarly to the second part of the proof of Theorem 3 to find slightly more than $\text{reg}_{\text{even}}(n, \delta)/2$ edge-disjoint Hamilton cycles in G . This proof is contained in Section 6.

3.3. η -extremal graphs. The following definition formalises the notion of “containing a subgraph close to $G_{n,\delta,\text{ext}}$ ”. For technical reasons we extend the definition to the case where α is negative – this will be used in Section 6 (with $|\alpha|$ very small). Note that if $\delta = (1/2 + \alpha)n$, then the vertex classes A and B of $G_{n,\delta,\text{ext}}$ have sizes roughly $(1/2 - \sqrt{\alpha/2})n$ and $(1/2 + \sqrt{\alpha/2})n$ respectively, and that $G_{n,\delta,\text{ext}}[B]$ is regular of degree roughly $(\alpha + \sqrt{\alpha/2})n$.

Definition 10. Let $\eta > 0$ and $-1/2 \leq \alpha \leq 1/2$, and let G be a graph on n vertices with $\delta(G) = (1/2 + \alpha)n$. Recall that $\alpha_+ = \max\{\alpha, 0\}$. We say that G is η -extremal if there exist disjoint $A, B \subseteq V(G)$ such that

- (E1) $|A| = (1/2 - \sqrt{\alpha_+/2} \pm \eta)n$;
- (E2) $|B| = (1/2 + \sqrt{\alpha_+/2} \pm \eta)n$;
- (E3) $e(A, B) > (1 - \eta)|A||B|$;
- (E4) $e(B) < (\alpha_+ + \sqrt{\alpha_+/2} + \eta)n|B|/2$.

Note that (E1) and (E2) together imply

- (E5) $n - |A \cup B| \leq 2\eta n$.

The following result states that if G is η -extremal, then $G[B]$ is “almost regular”.

Lemma 11. *Suppose $0 < \eta \ll \alpha, 1/2 - \alpha < 1/2$. Suppose G is an η -extremal graph on n vertices, with $\delta(G) = (1/2 + \alpha)n$, and let $A, B \subseteq V(G)$ be as in Definition 10.*

- (i) *For all vertices $v \in B$, we have $d_B(v) \geq (\alpha + \sqrt{\alpha/2} - 3\eta)n$.*
- (ii) *For all but at most $2\sqrt{\eta}n$ vertices $v \in B$, we have $d_B(v) \leq (\alpha + \sqrt{\alpha/2} + 2\sqrt{\eta})n$.*

Proof. (i) immediately follows from (E1) and (E5). Indeed, for all $v \in B$, we have

$$(2) \quad \begin{aligned} d_B(v) &\geq \delta(G) - d_A(v) - d_{\overline{A \cup B}}(v) \stackrel{(E5)}{\geq} \delta(G) - |A| - 2\eta n \\ &\stackrel{(E1)}{\geq} \left(\alpha + \sqrt{\frac{\alpha}{2}} - 3\eta \right) n, \end{aligned}$$

as desired.

Suppose (ii) fails. Then there exist at least $2\sqrt{\eta}n$ vertices in B with degree greater than $(\alpha + \sqrt{\alpha/2} + 2\sqrt{\eta})n$ in B . We therefore have

$$\begin{aligned} e_G(B) &= \frac{1}{2} \sum_{v \in B} d_B(v) \stackrel{(2)}{>} \frac{1}{2} \left(\left(\alpha + \sqrt{\frac{\alpha}{2}} - 3\eta \right) n|B| + 2\sqrt{\eta}n \cdot 2\sqrt{\eta}n \right) \\ &\geq \frac{1}{2} \left(\alpha + \sqrt{\frac{\alpha}{2}} + \eta \right) n|B|. \end{aligned}$$

But this contradicts (E4), so (ii) must hold. \square

4. THE NEAR-EXTREMAL CASE

Suppose that $0 < 1/n \ll \eta \ll \alpha < 1/2$, and that G is an η -extremal graph on n vertices with $\delta(G) = (1/2 + \alpha)n$. Recall that our aim in this case is to show that G contains a factor of degree $\text{reg}_{\text{even}}(n, \delta)/2$ which is a robust expander. Let $A, B \subseteq V(G)$ be as in Definition 10. We will first show that G contains a spanning subgraph G_1 which is close to $G_{n, \delta, \text{ext}}$ and satisfies $\delta(G_1) = \delta(G)$.

Lemma 12. *Suppose $0 < 1/n \ll \eta \ll 1/C \ll 1/2 - \alpha \leq 1/2$, so that in particular $0 \leq \alpha < 1/2$. Let G be an η -extremal graph on n vertices with $\delta := \delta(G) = (1/2 + \alpha)n$, and let $A, B \subseteq V(G)$ be as in Definition 10. Then there exists a spanning subgraph G_1 of G which satisfies the following properties:*

- (i) *A and B satisfy (E1)–(E4) for the graph G_1 . In particular, G_1 is η -extremal.*
- (ii) *$\delta(G_1) = \delta$.*
- (iii) *$e_{G_1}(A) < C\eta|A|^2$.*

Proof. We will define G_1 using a greedy algorithm. Initially, let $G_1 := G$. Suppose that $G_1[A]$ contains an edge xy such that $d_{G_1}(x), d_{G_1}(y) > \delta$. Then remove xy from G_1 , and continue in this way until G_1 contains no such edge. Note that we have $\delta(G_1) = \delta$, and (E1)–(E4) are not affected by these edge deletions, so G_1 satisfies (i) and (ii).

Suppose $e_{G_1}(A) \geq C\eta|A|^2$, and note that we have

$$\delta = \left(\frac{1}{2} + \alpha\right)n \leq \left(\frac{1}{2} + \sqrt{\frac{\alpha}{2}}\right)n \stackrel{\text{(E2)}}{\leq} |B| + \eta n.$$

(Indeed, $x \leq \sqrt{x/2}$ for all $0 \leq x \leq 1/2$.) If $v \in A$ is a vertex with $d_{G_1}(v) = \delta$, we therefore have

$$d_{G[A,B]}(v) = d_{G_1[A,B]}(v) \leq \delta - d_{G_1[A]}(v) \leq |B| + \eta n - d_{G_1[A]}(v).$$

Each edge in $G_1[A]$ must have at least one endpoint with degree δ in G_1 , so

$$\begin{aligned} e_G(A, B) &= \sum_{v \in A} d_{G[A,B]}(v) \leq |A||B| - \sum_{v \in A, d_{G_1}(v) = \delta} (d_{G_1[A]}(v) - \eta n) \\ &\leq |A||B| + \eta n^2 - e_{G_1}(A) \leq |A| \left(|B| + \eta \frac{n^2}{|A|} - C\eta|A| \right). \end{aligned}$$

Since $1/C \ll 1/2 - \alpha$, we have $C|A| \geq 2|B| + n^2/|A|$ by (E1) and (E2). Hence

$$e_G(A, B) \leq |A|(|B| - 2\eta|B|) = (1 - 2\eta)|A||B|,$$

which contradicts (E3). We therefore have $e_{G_1}(A) < C\eta|A|^2$, and so G_1 satisfies (iii) as desired. \square

Let G_1 be as in Lemma 12, and let G_2 be a degree-maximal even factor of G_1 . (So G_2 is an even-regular spanning subgraph of G_1 whose degree is as large as possible.) By Theorem 4, we have that

$$(3) \quad d(G_2) \geq \text{reg}_{\text{even}}(n, \delta) \geq \frac{n}{4} + \frac{\alpha n}{2} + \sqrt{\frac{\alpha}{2}}n - 2.$$

It can be shown that any degree-maximal even factor of $G_{n, \delta, \text{ext}}$ contains almost all edges inside the larger vertex class B . The following lemma uses a similar argument to prove a similar statement for G_1 .

Lemma 13. *Suppose $0 < 1/n \ll \eta \ll 1/C \ll \alpha, 1/2 - \alpha < 1/2$. Suppose that G is an η -extremal graph on n vertices with $\delta(G) = (1/2 + \alpha)n$. Let G_1 be the graph obtained by applying Lemma 12 to G , and let G_2 be a degree-maximal even factor of G_1 . Let $A, B \subseteq V(G)$ be as in Definition 10. Then for all but at most $3\eta^{1/4}n$ vertices $v \in B$, we have*

$$d_{G_2[B]}(v) \geq \left(\alpha + \sqrt{\frac{\alpha}{2}} - 3\eta^{1/4} \right) n.$$

Proof. Let r be the degree of G_2 . Suppose that $d_{G_2[B]}(v) < (\alpha + \sqrt{\alpha/2} - 3\eta^{1/4})n$ for more than $3\eta^{1/4}n$ vertices. Then by Lemma 11(ii), we have

$$\begin{aligned} r|B| &= \sum_{v \in B} d_{G_2}(v) = e_{G_2}(A, B) + 2e_{G_2}(B) \\ &\leq r|A| + \left(\alpha + \sqrt{\frac{\alpha}{2}} + 2\sqrt{\eta} \right) n|B| + 4\sqrt{\eta}n^2 - 3\eta^{1/4}n \cdot 3\eta^{1/4}n \\ &\leq r|A| + \left(\alpha + \sqrt{\frac{\alpha}{2}} - 3\sqrt{\eta} \right) n|B|. \end{aligned}$$

Since $|B| - |A| \geq (\sqrt{2\alpha} - 2\eta)n$ by (E1) and (E2), it follows that

$$\sqrt{2\alpha}rn \leq \left(\alpha + \sqrt{\frac{\alpha}{2}} - 3\sqrt{\eta} \right) n|B| + 2\eta n^2,$$

and hence

$$\begin{aligned} r &\leq \left(\sqrt{\frac{\alpha}{2}} + \frac{1}{2} - 3\sqrt{\frac{\eta}{2\alpha}} \right) |B| + \eta\sqrt{\frac{2}{\alpha}}n \\ &\stackrel{(E2)}{\leq} \left(\sqrt{\frac{\alpha}{2}} + \frac{1}{2} - 3\sqrt{\frac{\eta}{2\alpha}} \right) \left(\frac{1}{2} + \sqrt{\frac{\alpha}{2}} + \eta \right) n + \eta\sqrt{\frac{2}{\alpha}}n \\ &\leq \left(\frac{1}{4} + \frac{\alpha}{2} + \sqrt{\frac{\alpha}{2}} - \eta^{3/4} \right) n. \end{aligned}$$

(In the last inequality we used that $\eta \ll \alpha$.) It therefore follows from (3) that $r < \text{reg}_{\text{even}}(n, \delta)$. But G_2 was chosen to be degree-maximal, so this is a contradiction. \square

We now collect some robust expansion properties of G_2 . For convenience, if $X \subseteq V(G_2)$, we shall write $X_A := X \cap A$ and $X_B := X \cap B$. In particular, if $S \subseteq V(G)$ then (for example) $RN_\nu(S_A)_B = RN_\nu(S \cap A) \cap B$.

Lemma 14. *Suppose that $0 < 1/n \ll \nu \ll \eta \ll \mu \ll \tau \ll \lambda \ll 1/C \ll \alpha, 1/2 - \alpha < 1/2$. Suppose that G is an η -extremal graph on n vertices with $\delta(G) = (1/2 + \alpha)n$. Let G_1 be the graph obtained by applying Lemma 12 to G , and let G_2 be a degree-maximal even factor of G_1 . Let $A, B \subseteq V(G)$ be as in Definition 10. Then in the graph G_2 , the following properties all hold.*

- (i) *If $S \subseteq A$ with $|S| \geq |A|/2$, then $|RN_\nu(S)_B| \geq (1 - \mu)|B|$.*
- (ii) *If $S \subseteq B$ with $|S| \geq |B|/2$, then $|RN_\nu(S)_A| \geq (1 - \mu)|A|$.*
- (iii) *If $S \subseteq A$ with $|S| \geq \tau n/3$, then $|RN_\nu(S)_B| \geq |B|/2 + \lambda n$.*
- (iv) *If $S \subseteq B$ with $|S| \geq \tau n/3$, then $|RN_\nu(S)_A| \geq |A|/2 + \lambda n$.*
- (v) *If $S \subseteq B$, then $|RN_\nu(S)_B| \geq |S| - \mu n$.*

Proof. Write $d := d(G_2)$. We first prove (i). Suppose $S \subseteq A$ with $|S| \geq |A|/2$. Lemma 11(ii) implies that in G_2 all but at most $2\sqrt{\eta}n \leq \mu|B|$ vertices $v \in B$ satisfy

$$\begin{aligned} d_A(v) &= d - d_{\overline{A \cup B}}(v) - d_B(v) \\ &\stackrel{(3),(E5)}{\geq} \left(\frac{1}{4} + \frac{\alpha}{2} + \sqrt{\frac{\alpha}{2}} \right) n - 2 - 2\eta n - \left(\alpha + \sqrt{\frac{\alpha}{2}} + 2\sqrt{\eta} \right) n \\ &\geq \left(\frac{1}{4} - \frac{\alpha}{2} - 3\sqrt{\eta} \right) n \geq \left(\frac{1}{4} - \frac{1}{2}\sqrt{\frac{\alpha}{2}} + \eta \right) n + \nu n \stackrel{(E1)}{\geq} \frac{|A|}{2} + \nu n, \end{aligned}$$

where the third inequality follows since $x < \sqrt{x}/2$ for all $0 < x < 1/4$. Thus in the graph G_2 we have $|N_A(v) \cap S| \geq \nu n$, and hence $v \in RN_\nu(S)$, for each such v . The result therefore follows.

We now prove (ii). Suppose $S \subseteq B$ with $|S| \geq |B|/2$. Let $A' \subseteq A$ be the set of vertices which in G_2 have less than $|B|/2 + \nu n$ neighbours inside B . Each vertex $v \in A'$ must satisfy

$$\begin{aligned} d_A(v) &= d - d_{\overline{A \cup B}}(v) - d_B(v) \\ &\stackrel{(3),(E5)}{\geq} \left(\frac{1}{4} + \frac{\alpha}{2} + \sqrt{\frac{\alpha}{2}} \right) n - 2 - 2\eta n - \frac{|B|}{2} - \nu n \\ &\stackrel{(E2)}{\geq} \left(\frac{1}{4} + \frac{\alpha}{2} + \sqrt{\frac{\alpha}{2}} - 2\eta - \nu \right) n - 2 - \left(\frac{1}{4} + \frac{1}{2}\sqrt{\frac{\alpha}{2}} + \eta \right) n \\ &\geq \frac{\alpha}{2}n, \end{aligned}$$

and so we have $e_{G_2}(A) \geq \alpha n|A'|/4$. But by Lemma 12(iii) we have $e_{G_2}(A) \leq e_{G_1}(A) < C\eta|A|^2$. Therefore

$$|A'| \leq \frac{4C\eta}{\alpha} \cdot \frac{|A|^2}{n} \leq \sqrt{\eta} \frac{|A|^2}{n} \leq \sqrt{\eta}|A| \leq \mu|A|.$$

However, our assumption that $|S| \geq |B|/2$ and the definition of A' together imply that every vertex $v \in A \setminus A'$ satisfies $|N_B(v) \cap S| \geq \nu n$. Therefore $|RN_\nu(S)| \geq |A \setminus A'| \geq (1 - \mu)|A|$, as required.

We now prove (iii). Suppose $S \subseteq A$ with $|S| \geq \tau n/3$. Then we double-count the edges between S and B in G_2 . The definition of a robust neighbourhood implies that

$$e_{G_2}(S, B) = e_{G_2}(S, RN_\nu(S)_B) + e_{G_2}(S, B \setminus RN_\nu(S)_B) \leq |S||RN_\nu(S)_B| + \nu n^2.$$

On the other hand, Lemma 12(iii) implies that

$$\begin{aligned} e_{G_2}(S, B) &\geq d|S| - 2e_{G_2}(S, A) - e_{G_2}(S, \overline{A \cup B}) \stackrel{(E5)}{\geq} d|S| - 2C\eta|A|^2 - 2\eta n^2 \\ &\geq d|S| - 3C\eta n^2. \end{aligned}$$

Combining the two inequalities yields

$$\begin{aligned}
|RN_\nu(S)_B| &\geq d - 3C\eta \frac{n^2}{|S|} - \nu \frac{n^2}{|S|} \\
&\stackrel{(3)}{\geq} \left(\frac{1}{4} + \frac{\alpha}{2} + \sqrt{\frac{\alpha}{2}} \right) n - 2 - \frac{9C\eta}{\tau} n - \frac{3\nu}{\tau} n \\
&\stackrel{(E2)}{\geq} \frac{|B|}{2} + \left(\frac{\alpha}{2} + \frac{1}{2} \sqrt{\frac{\alpha}{2}} - \eta - \frac{9C\eta}{\tau} - \frac{3\nu}{\tau} \right) n - 2 \geq \frac{|B|}{2} + \frac{\alpha}{2} n,
\end{aligned}$$

and so the result follows.

We now prove (iv). Suppose $S \subseteq B$ with $|S| \geq \tau n/3$. Then we double-count the edges between S and A in G_2 . As before, we have

$$(4) \quad e_{G_2}(S, A) \leq |S| |RN_\nu(S)_A| + \nu n^2.$$

On the other hand,

$$e_{G_2}(S, A) \geq d|S| - \sum_{v \in S} d_B(v) - \sum_{v \in S} d_{A \cup B}(v) \stackrel{(E5)}{\geq} d|S| - \sum_{v \in S} d_B(v) - 2\eta n^2.$$

Lemma 11(ii) implies that

$$\sum_{v \in S} d_B(v) \leq 2\sqrt{\eta} n^2 + \left(\alpha + \sqrt{\frac{\alpha}{2}} + 2\sqrt{\eta} \right) n|S|,$$

and so

$$\begin{aligned}
e_{G_2}(S, A) &\stackrel{(3),(E5)}{\geq} \left(\frac{1}{4} + \frac{\alpha}{2} + \sqrt{\frac{\alpha}{2}} \right) n|S| - 2|S| - \left(\alpha + \sqrt{\frac{\alpha}{2}} + 2\sqrt{\eta} \right) n|S| - (2\eta + 2\sqrt{\eta}) n^2 \\
&\geq \left(\frac{1}{4} - \frac{\alpha}{2} \right) n|S| - 5\sqrt{\eta} n^2.
\end{aligned}$$

Combining this with (4) shows that in G_2 we have

$$\begin{aligned}
|RN_\nu(S)_A| &\geq \left(\frac{1}{4} - \frac{\alpha}{2} \right) n - 6\sqrt{\eta} \cdot \frac{n^2}{|S|} \geq \left(\frac{1}{4} - \frac{\alpha}{2} \right) n - \frac{18\sqrt{\eta}}{\tau} n \\
&= \left(\frac{1}{4} - \frac{1}{2} \sqrt{\frac{\alpha}{2}} \right) n + \left(\frac{1}{2} \sqrt{\frac{\alpha}{2}} - \frac{\alpha}{2} \right) n - \frac{18\sqrt{\eta}}{\tau} n \\
&\stackrel{(E1)}{\geq} \frac{|A|}{2} + \frac{1}{2} \left(\frac{1}{2} \sqrt{\frac{\alpha}{2}} - \frac{\alpha}{2} \right) n \geq \frac{|A|}{2} + \lambda n,
\end{aligned}$$

and so the result follows. (Here we used that $\sqrt{x}/2 > x$ for all $0 < x < 1/4$.)

Finally, we prove (v). Suppose $S \subseteq B$. Recall that $e'_{G_2}(S, RN_\nu(S)_B)$ denotes the number of ordered pairs (u, v) of vertices of G_2 such that $uv \in E(G_2)$, $u \in S$ and $v \in RN_\nu(S)_B$. (Note that this may not equal $e(S, RN_\nu(S)_B)$, as we may have $S \cap RN_\nu(S)_B \neq \emptyset$.) By Lemma 11(ii), counting from $RN_\nu(S)_B$ yields

$$e'_{G_2}(S, RN_\nu(S)_B) \leq \left(\alpha + \sqrt{\frac{\alpha}{2}} + 2\sqrt{\eta} \right) n |RN_\nu(S)_B| + 2\sqrt{\eta} n^2.$$

By Lemma 13, counting from S yields

$$e'_{G_2}(S, RN_\nu(S)_B) \geq \left(\alpha + \sqrt{\frac{\alpha}{2}} - 3\eta^{\frac{1}{4}} \right) n|S| - 3\eta^{\frac{1}{4}}n^2 - \nu n^2.$$

Combining the two inequalities yields $|RN_\nu(S)_B| \geq |S| - \mu n$ as desired. \square

We are now in a position to prove Theorem 3 for η -extremal graphs.

Lemma 15. *Suppose $0 < 1/n \ll \eta \ll \alpha, 1/2 - \alpha < 1/2$. If G is an η -extremal graph on n vertices with $\delta := \delta(G) = (1/2 + \alpha)n$, then G contains at least $\text{reg}_{\text{even}}(n, \delta)/2$ edge-disjoint Hamilton cycles.*

Proof. Let $\tau_0 := \tau(1/4)$ be the constant returned by Theorem 9. Choose additional constants ν, μ, τ, λ and C such that

$$0 < \nu \ll \eta \ll \mu \ll \tau \ll \lambda \ll 1/C \ll \alpha, 1/2 - \alpha, \tau_0.$$

Take $A, B \subseteq V(G)$ as in Definition 10. Apply Lemma 12 to G and C to obtain a spanning subgraph G_1 . Let G_2 be a degree-maximal even factor of G_1 . Note that Lemma 14 may be applied to G_2 .

Claim: G_2 is a robust (ν, τ) -expander.

Note that the claim immediately implies the desired result. Indeed, any robust (ν, τ) -expander is also a robust (ν, τ_0) -expander, and so Theorem 9 implies that G_2 may be decomposed into Hamilton cycles. Moreover, Lemma 12 implies that $\delta(G_1) = \delta$ and so $d(G_2) \geq \text{reg}_{\text{even}}(n, \delta)$. Hence the Hamilton decomposition of G_2 yields the desired collection of $d(G_2)/2 \geq \text{reg}_{\text{even}}(n, \delta)/2$ edge-disjoint Hamilton cycles.

To prove the claim, consider $S \subseteq V(G)$ with $\tau n \leq |S| \leq (1 - \tau)n$. We will use Lemma 14 to show that in G_2 we have $|RN_\nu(S)| \geq |S| + \nu n$. We will split the proof into cases depending on the sizes of $S_A = S \cap A$ and $S_B = S \cap B$. Note that $|S_{\overline{A \cup B}}| \leq 2\eta n$ by (E5).

Case 1: $|S_A| \leq \tau n/3, |S_B| \leq \tau n/3$.

In this case, we have

$$|S| \stackrel{\text{(E5)}}{\leq} \frac{2\tau}{3} + 2\eta n < \tau n,$$

which is a contradiction.

Case 2: $\tau n/3 \leq |S_A| \leq |A|/2, |S_B| \leq \tau n/3$.

In this case, by Lemma 14(iii) we have

$$\begin{aligned} |RN_\nu(S)| &\geq |RN_\nu(S_A)_B| \geq \frac{|B|}{2} + \lambda n \stackrel{\text{(E5)}}{\geq} \frac{|A|}{2} + \sqrt{\frac{\alpha}{2}}n - 2\eta n + \lambda n \\ &\geq \frac{|A|}{2} + \frac{\tau}{3}n + 2\eta n + \nu n \geq |S| + \nu n, \end{aligned}$$

as desired.

Case 3: $|S_A| \geq |A|/2, |S_B| \leq \tau n/3$.

In this case, by Lemma 14(i) we have

$$\begin{aligned} |RN_\nu(S)| &\geq |RN_\nu(S_A)_B| \geq (1 - \mu)|B| \geq |A| + 2\sqrt{\frac{\alpha}{2}}n - 2\eta n - \mu n \\ &\geq |A| + \frac{\tau}{3}n + 2\eta n + \nu n \geq |S| + \nu n, \end{aligned}$$

as desired.

Case 4: $|S_A| \leq |A|/2$, $|S_B| \geq \tau n/3$.

In this case, by Lemma 14(iv) and (v), we have

$$\begin{aligned} |RN_\nu(S)| &\geq |RN_\nu(S_B)_A| + |RN_\nu(S_B)_B| \geq \frac{|A|}{2} + \lambda n + |S_B| - \mu n \\ &\geq |S_A| + |S_B| + 2\eta n + \nu n \geq |S| + \nu n, \end{aligned}$$

as desired.

Case 5: $|S_A| \geq |A|/2$, $\tau n/3 \leq |S_B| \leq |B|/2$.

In this case, by Lemma 14(i) and (iv), we have

$$\begin{aligned} |RN_\nu(S)| &\geq |RN_\nu(S_A)_B| + |RN_\nu(S_B)_A| \geq |B| + \frac{|A|}{2} + (\lambda - \mu)n \\ &\geq \frac{|B|}{2} + |A| + (\lambda - \mu)n \geq |S_B| + |S_A| + 2\eta n + \nu n \geq |S| + \nu n, \end{aligned}$$

as desired, where the third inequality follows since $|B| \geq |A|$ by (E1) and (E2).

Case 6: $|S_A| \geq |A|/2$, $|S_B| \geq |B|/2$.

In this case, by Lemma 14(i) and (ii), we have

$$\begin{aligned} |RN_\nu(S)| &\geq |RN_\nu(S_A)_B| + |RN_\nu(S_B)_A| \geq |B| + |A| - 2\mu n \\ &\stackrel{(E5)}{\geq} n - (2\eta + 2\mu)n \geq (1 - \tau)n + \nu n \geq |S| + \nu n, \end{aligned}$$

as desired.

Thus in all cases we have $|RN_\nu(S)| \geq |S| + \nu n$. Indeed, if $|S_B| \leq \tau n/3$ this follows by Cases 1, 2 and 3; if $\tau n/3 \leq |S_B| \leq |B|/2$ this follows by Cases 4 and 5; and if $|S_B| \geq |B|/2$ this follows by Cases 4 and 6. Hence G_2 is a robust (ν, τ) -expander as desired. This proves the claim and hence the lemma. \square

5. THE NON-EXTREMAL CASE

Suppose now that G is not η -extremal. Our first aim is to find a sparse even factor H of G which is a robust expander. This has essentially already been done in [12], but for digraphs. We first require the following definition, which generalises the notion of robust expanders to digraphs.

Definition 16. Let D be a digraph on n vertices. Given $0 < \nu \leq \tau < 1$, we define the ν -robust outneighbourhood $RN_{\nu, D}^+(S)$ of S to be the set of all vertices $v \in V(D)$ which have at least νn inneighbours in S . We say that D is a *robust (ν, τ) -outexpander* if for all $S \subseteq V(D)$ with $\tau n \leq |S| \leq (1 - \tau)n$, we have $|RN_{\nu, D}^+(S)| \geq |S| + \nu n$.

We will now quote three lemmas from [12]. Lemma 17 implies that our given graph G is a robust expander. We will use Lemmas 18 and 19 to deduce Lemma 20, which together with Lemma 17 implies that G contains a sparse even factor H which is still a robust expander.

Lemma 17. *Suppose $0 < \nu \leq \tau \leq \varepsilon < 1/2$ and $\varepsilon \geq 2\nu/\tau$. Let G be a graph on n vertices with minimum degree $\delta(G) \geq (1/2 + \varepsilon)n$. Then G is a robust (ν, τ) -expander.*

Lemma 18. *Suppose $0 < 1/n \ll \eta \ll \nu, \tau, \alpha < 1$. Suppose that G is a robust (ν, τ) -expander on n vertices with $\delta(G) \geq \alpha n$. Then the edges of G can be oriented in such a way that the resulting oriented graph G' satisfies the following properties:*

- (i) G' is a robust $(\nu/4, \tau)$ -outexpander.
- (ii) $d_{G'}^+(x), d_{G'}^-(x) = (1 \pm \eta)d_G(x)/2$ for every vertex x of G .

An r -factor of a digraph G is a spanning subdigraph of G in which every vertex has in- and outdegree r .

Lemma 19. *Suppose $0 < 1/n \ll \nu' \ll \xi \ll \nu \leq \tau \ll \alpha < 1$. Let G be a robust (ν, τ) -outexpander on n vertices with $\delta^+(G), \delta^-(G) \geq \alpha n$. Then G contains a ξn -factor which is still a robust (ν', τ) -outexpander.*

Lemma 20. *Suppose $0 < 1/n \ll \nu' \ll \varepsilon \ll \nu \ll \tau \ll \alpha < 1$, and suppose in addition that εn is an even integer. If G is a robust (ν, τ) -expander on n vertices with $\delta(G) \geq \alpha n$, then there exists an εn -factor H of G which is a robust (ν', τ) -expander.*

Proof. We apply Lemma 18 to orient the edges of G , forming an oriented graph G' which is a robust $(\nu/4, \tau)$ -outexpander and which satisfies $\delta^+(G'), \delta^-(G') \geq \alpha n/3$. We then apply Lemma 19 to find an $\varepsilon n/2$ -factor H of G' which is a robust (ν', τ) -outexpander. Now remove the orientation on the edges of H to obtain a robust (ν', τ) -expander which is an εn -factor of G , as desired. \square

We will now show that even after removing a sparse factor H , our given graph G still contains an even factor of degree at least $\text{reg}_{\text{even}}(n, \delta)$. To do this, we first show that $G - H$ is still non-extremal.

Lemma 21. *Suppose $0 < 1/n \ll \varepsilon \ll \eta \ll 1/2 - \alpha$, and that $-\varepsilon \leq \alpha < 1/2$. Let G be a graph on n vertices with $\delta(G) = (1/2 + \alpha)n$ which is not 2η -extremal. Suppose H is an εn -factor of G . Then $G - H$ is not η -extremal.*

Proof. Suppose $A, B \subseteq V(G)$ are disjoint with $|A|$ and $|B|$ satisfying (E1) and (E2) of Definition 10. Let $G' := G - H$. Since G is not 2η -extremal, we must have either $e_G(A, B) \leq (1 - 2\eta)|A||B|$ or $e_G(B) \geq (\alpha_+ + \sqrt{\alpha_+/2} + 2\eta)n|B|/2$. In the former case we have

$$e_{G'}(A, B) \leq e_G(A, B) < (1 - \eta)|A||B|,$$

and in the latter case we have

$$\begin{aligned} e_{G'}(B) &\geq e_G(B) - \varepsilon n|B| \geq \frac{1}{2} \left(\alpha_+ + \sqrt{\frac{\alpha_+}{2}} + 2\eta - 2\varepsilon \right) n|B| \\ &\geq \frac{1}{2} \left((\alpha - \varepsilon)_+ + \sqrt{\frac{(\alpha - \varepsilon)_+}{2}} + \frac{3\eta}{2} \right) n|B|. \end{aligned}$$

Since $\delta(G - H) = (1/2 + \alpha - \varepsilon)n$, it follows that $G - H$ is not η -extremal. \square

We now show that $G - H$ contains a large even factor. We will do this using the well-known result of Tutte [17], given below.

Theorem 22. *Let G be a graph. Given disjoint $S, T \subseteq V(G)$ and $r \in \mathbb{N}$, let $Q_r(S, T)$ be the number of connected components C of $G - (S \cup T)$ such that $r|C| + e(C, T)$ is odd, and let*

$$(5) \quad R_r(S, T) := \sum_{v \in T} d(v) - e(S, T) + r(|S| - |T|).$$

Then G contains an r -factor if and only if $Q_r(S, T) \leq R_r(S, T)$ for all disjoint $S, T \subseteq V(G)$.

In proving the following lemma, we follow a similar approach to that used in [1]. We will also make frequent and implicit use of the inequality $\sqrt{x} \leq \sqrt{x+h} \leq \sqrt{x} + \sqrt{h}$ for $x, h \geq 0$.

Lemma 23. *Suppose $0 < 1/n \ll \varepsilon \ll \eta \ll 1/2 - \alpha$ and that $-\varepsilon \leq \alpha < 1/2$. Let G be a graph on n vertices with minimum degree $\delta := \delta(G) = (1/2 + \alpha)n$, and suppose that G is not η -extremal. Let*

$$r := \frac{n}{4} + \frac{(\alpha + \varepsilon)n}{2} + \sqrt{\frac{\alpha + \varepsilon}{2}}n,$$

and suppose that r is an even integer. Then G contains an r -factor.

Proof. Let S, T be two arbitrary disjoint subsets of $V(G)$. We will show that $Q_r(S, T) \leq R_r(S, T)$, from which the result follows by Theorem 22. We first note a useful bound on $Q_r(S, T)$. If $\delta \geq |S| + |T|$ then every vertex outside $S \cup T$ has at least $\delta - |S| - |T|$ neighbours outside $S \cup T$, so every component of $G - (S \cup T)$ contains at least $\delta - |S| - |T| + 1$ vertices. Thus

$$(6) \quad Q_r(S, T) \leq \frac{n - |S| - |T|}{\delta - |S| - |T| + 1} \quad \text{if } \delta \geq |S| + |T|.$$

Also, note that we always have

$$(7) \quad \delta - r = \left(\frac{1}{4} + \frac{\alpha - \varepsilon}{2} - \sqrt{\frac{\alpha + \varepsilon}{2}} \right) n = \left(\frac{1}{4} + \frac{\alpha + \varepsilon}{2} - \sqrt{\frac{\alpha + \varepsilon}{2}} - \varepsilon \right) n \geq \varepsilon n,$$

since $1/4 + x - \sqrt{x} = (1/2 - \sqrt{x})^2 > 0$ for all $0 \leq x < 1/4$ and since $\varepsilon \ll 1/2 - \alpha$. We will now split the proof into cases depending on $|S|$ and $|T|$.

Case 1: $|T| \leq r - 1$, $|S| \leq \delta - r$, and $|S| + |T| \geq 3$.

We have

$$(8) \quad \begin{aligned} R_r(S, T) &\stackrel{(5)}{=} \sum_{v \in T} (d(v) - r) + \sum_{v \in S} (r - d_T(v)) \geq |T|(\delta - r) + \sum_{v \in S} 1 \\ &\stackrel{(7)}{\geq} |S| + |T|. \end{aligned}$$

Let $k := |S| + |T|$. By (6) and (8) it suffices to show that $k \geq (n - k)/(\delta - k + 1)$. This is equivalent to showing that

$$\delta k - k^2 + 2k - n = (k - 2)(\delta - k) + 2\delta - n \geq 0.$$

We have $3 \leq k \leq \delta - 1$ and the function $(k - 2)(\delta - k)$ is concave, so it must be minimised in this range when $k = 3$ or when $k = \delta - 1$. In either case, we have

$$(k - 2)(\delta - k) + 2\delta - n = \delta - 3 + 2\delta - n \geq \delta - 3 - 2\epsilon n \geq 0$$

as desired.

Case 2: $0 \leq |S| + |T| \leq 2$.

If $S = T = \emptyset$, then we have $Q_r(S, T) = R_r(S, T) = 0$ (since r is even). So suppose that $|S| + |T| > 0$. Then it follows from (6) that

$$Q_r(S, T) < \frac{n}{\delta - 1} \leq \frac{3n}{n} = 3.$$

If $T \neq \emptyset$, we have

$$R_r(S, T) \stackrel{(5)}{\geq} \delta|T| - 1 - r|T| \stackrel{(7)}{\geq} 3.$$

If $T = \emptyset$, we have $|S| \geq 1$ and so by (5) we have $R_r(S, T) \geq r \geq 3$. We therefore have $Q_r(S, T) \leq R_r(S, T)$ in all cases.

Case 3: $|T| \geq r$ or $|S| \geq \delta - r$, but not both.

We have

$$\begin{aligned} R_r(S, T) &\stackrel{(5)}{\geq} (\delta - r)|T| - |S||T| + r|S| \\ (9) \quad &= (|T| - r)(\delta - r - |S|) + r(\delta - r) \\ &\geq r(\delta - r) \stackrel{(7)}{\geq} \frac{\epsilon}{4}n^2. \end{aligned}$$

(Note that (9) holds regardless of the values of $|S|$ and $|T|$.) Moreover, we have $Q_r(S, T) \leq n$. Hence $Q_r(S, T) \leq R_r(S, T)$ as desired.

Case 4: $|T| \geq r$, $|S| \geq \delta - r$, and $|T| \neq (n + 2r - \delta)/2 \pm 3\sqrt{\epsilon}n$.

The right hand side of (9) is clearly minimised when $|S| + |T| = n$. It therefore suffices to consider this case alone, yielding

$$\begin{aligned} R_r(S, T) - Q_r(S, T) &\geq (\delta - r)|T| - (n - |T|)|T| + r(n - |T|) - n \\ &= |T|^2 + (\delta - 2r - n)|T| + n(r - 1). \end{aligned}$$

Define a polynomial $f : \mathbb{R} \rightarrow \mathbb{R}$ by

$$f(x) = x^2 + (\delta - 2r - n)x + n(r - 1).$$

Suppose this quadratic has real zeroes at τ_1 and τ_2 , with $\tau_1 < \tau_2$. Then for $|T| \leq \tau_1$ and $|T| \geq \tau_2$, we must have $R_r(S, T) - Q_r(S, T) \geq 0$. The discriminant D of f is given by

$$D = (n + 2r - \delta)^2 - 4n(r - 1) = (n + 2r - \delta)^2 - \left(1 + 2(\alpha + \epsilon) + 2\sqrt{2(\alpha + \epsilon)} - \frac{4}{n}\right)n^2.$$

But

$$(10) \quad n + 2r - \delta = \left(1 + \varepsilon + \sqrt{2(\alpha + \varepsilon)}\right) n,$$

so

$$(n + 2r - \delta)^2 = \left(1 + \varepsilon^2 + 2(\alpha + \varepsilon) + 2\varepsilon + 2(1 + \varepsilon)\sqrt{2(\alpha + \varepsilon)}\right) n^2$$

and

$$D = \varepsilon \left(\varepsilon + 2 + 2\sqrt{2(\alpha + \varepsilon)}\right) n^2 - 4n.$$

Hence $0 < D \leq 5\varepsilon n^2$. In particular, the quadratic does indeed have two real zeroes $\tau_1 < \tau_2$, and from the quadratic formula we have

$$\tau_1 \geq \frac{n + 2r - \delta - 3\sqrt{\varepsilon}n}{2}, \quad \tau_2 \leq \frac{n + 2r - \delta + 3\sqrt{\varepsilon}n}{2}.$$

Since we are in Case 4, we therefore have either $|T| \leq \tau_1$ or $|T| \geq \tau_2$, and the result follows.

Case 5: $|T| = (n + 2r - \delta)/2 \pm 3\sqrt{\varepsilon}n$ and $\delta - r \leq |S| \leq (n - 2r + \delta)/2 - 3\sqrt{\varepsilon}n$.

(Note that our condition on $|T|$ implies that we cannot have $|S| > (n - 2r + \delta)/2 + 3\sqrt{\varepsilon}n$.) Let $x_0 := (n + 2r - \delta)/2 + 3\sqrt{\varepsilon}n \geq |T|$. We then have

$$R_r(S, T) \stackrel{(9)}{\geq} (x_0 - r)(\delta - r - |S|) + r(\delta - r).$$

Since $x_0 + |S| \leq n$, we may now argue exactly as in Case 4 (with x_0 in place of $|T|$) to show that $R_r(S, T) \geq Q_r(S, T)$.

Case 6: $|T| = (n + 2r - \delta)/2 \pm 3\sqrt{\varepsilon}n$ and $|S| = (n - 2r + \delta)/2 \pm 3\sqrt{\varepsilon}n$.

In this case, we will use the fact that G is not η -extremal. From (10), we have

$$\left| \frac{n + 2r - \delta}{2} - \left(\frac{1}{2} + \sqrt{\frac{\alpha_+}{2}}\right) n \right| \leq \left(\frac{\varepsilon}{2} + \sqrt{\frac{\varepsilon}{2}}\right) n.$$

Since $\varepsilon \ll \eta$, we may conclude that

$$\left| |T| - \left(\frac{1}{2} + \sqrt{\frac{\alpha_+}{2}}\right) n \right| < \eta n.$$

A similar argument shows that

$$\left| |S| - \left(\frac{1}{2} - \sqrt{\frac{\alpha_+}{2}}\right) n \right| < \eta n.$$

Since G is not η -extremal, this implies that either $e(S, T) \leq (1 - \eta)|S||T|$ or

$$e(T) \geq \frac{1}{2} \left(\alpha_+ + \sqrt{\frac{\alpha_+}{2}} + \eta\right) n|T|.$$

Case 6a: $e(S, T) \leq (1 - \eta)|S||T|$.

Then we have

$$\begin{aligned}
R_r(S, T) - Q_r(S, T) &\stackrel{(5)}{\geq} (\delta - r)|T| - (1 - \eta)|S||T| + r|S| - n \\
&\geq (\delta - r)|T| - (1 - \eta)(n - |T|)|T| + r(n - |T| - 6\sqrt{\varepsilon}n) - n \\
&= (1 - \eta)|T|^2 + (\delta - 2r - (1 - \eta)n)|T| + (1 - 6\sqrt{\varepsilon})nr - n.
\end{aligned}$$

Write this quadratic as $a|T|^2 + b|T| + c$, and let the discriminant be D . We then have

$$\begin{aligned}
b^2 &= ((1 - \eta)n + 2r - \delta)^2 \stackrel{(10)}{=} \left(1 - \eta + \varepsilon + \sqrt{2(\alpha + \varepsilon)}\right)^2 n^2 \\
&= \left((1 - \eta)^2 + \varepsilon^2 + 2(\alpha + \varepsilon) + 2(1 - \eta)\varepsilon + 2(1 - \eta)\sqrt{2(\alpha + \varepsilon)} + 2\varepsilon\sqrt{2(\alpha + \varepsilon)}\right) n^2 \\
&\leq \left((1 - \eta)^2 + 2\alpha + 2(1 - \eta)\sqrt{2(\alpha + \varepsilon)} + \varepsilon^{\frac{1}{3}}\right) n^2
\end{aligned}$$

and

$$\begin{aligned}
4ac &= 4(1 - \eta)(1 - 6\sqrt{\varepsilon})nr - 4(1 - \eta)n \\
&\geq (1 - \eta)(1 - 6\sqrt{\varepsilon}) \left(1 + 2(\alpha + \varepsilon) + 2\sqrt{2(\alpha + \varepsilon)}\right) n^2 - 4n \\
&\geq (1 - \eta) \left(1 + 2\alpha + 2\sqrt{2(\alpha + \varepsilon)}\right) n^2 - \varepsilon^{\frac{1}{3}}n^2.
\end{aligned}$$

Thus

$$\begin{aligned}
D &= b^2 - 4ac \leq \left((1 - \eta)^2 - (1 - \eta) + 2\eta\alpha + 2\varepsilon^{\frac{1}{3}}\right) n^2 \\
&= \left(-\eta(1 - \eta - 2\alpha) + 2\varepsilon^{\frac{1}{3}}\right) n^2 < 0,
\end{aligned}$$

where the last line follows since $\varepsilon \ll \eta \ll 1/2 - \alpha$ and $\alpha < 1/2$. Hence this quadratic has no real zeroes, and $R_r(S, T) - Q_r(S, T) \geq 0$ as desired.

Case 6b: $e(T) \geq (\alpha_+ + \sqrt{\alpha_+/2} + \eta)n|T|/2$ and $e(S, T) \geq (1 - \eta)|S||T|$.

Then we have

$$\begin{aligned}
\sum_{v \in T} d(v) &\geq e(S, T) + 2e(T) \\
&\geq \left((1 - \eta)|S| + \left(\alpha_+ + \sqrt{\frac{\alpha_+}{2}} + \eta\right)n\right) |T| \\
&\geq \left((1 - \eta) \left(\frac{n - 2r + \delta}{2} - 3\sqrt{\varepsilon}n\right) + \left(\alpha + \sqrt{\frac{\alpha_+}{2}} + \eta\right)n\right) |T| \\
&\stackrel{(10)}{\geq} \left((1 - \eta) \left(\frac{1}{2} - \frac{\varepsilon}{2} - \sqrt{\frac{\alpha + \varepsilon}{2}} - 3\sqrt{\varepsilon}\right) + \alpha + \sqrt{\frac{\alpha_+}{2}} + \eta\right) n|T| \\
&\geq \left(\frac{1}{2} - \sqrt{\frac{\alpha + \varepsilon}{2}} - 4\sqrt{\varepsilon} - \frac{\eta}{2} + \alpha + \sqrt{\frac{\alpha_+}{2}} + \eta\right) n|T| \\
&\geq \left(\frac{1}{2} + \frac{\eta}{3} + \alpha\right) n|T|.
\end{aligned}$$

Hence

$$\begin{aligned}
R_r(S, T) - Q_r(S, T) &\stackrel{(5)}{\geq} \sum_{v \in T} d(v) - (n - |T|)|T| + r(|S| - |T|) - n \\
&\geq \sum_{v \in T} d(v) + |T|^2 - n|T| + r(n - |T| - 6\sqrt{\varepsilon}n) - r|T| - n \\
&= \sum_{v \in T} d(v) + |T|^2 - (n + 2r)|T| + (1 - 6\sqrt{\varepsilon})nr - n \\
&\geq |T|^2 - \left(\left(\frac{1}{2} - \frac{\eta}{3} - \alpha \right) n + 2r \right) |T| + (1 - 6\sqrt{\varepsilon})nr - n \\
&\geq |T|^2 - \left(1 + \varepsilon + \sqrt{2(\alpha + \varepsilon)} - \frac{\eta}{3} \right) n|T| + (1 - 6\sqrt{\varepsilon})nr - n \\
&\geq |T|^2 - \left(1 + \sqrt{2(\alpha + \varepsilon)} - \frac{\eta}{4} \right) n|T| + (1 - 7\sqrt{\varepsilon})nr.
\end{aligned}$$

Write this quadratic as $|T|^2 + b|T| + c$, and let the discriminant be D . We then have

$$b^2 \leq \left(1 + 2\alpha + 2\varepsilon + \frac{\eta^2}{16} + 2\sqrt{2(\alpha + \varepsilon)} - \frac{\eta}{2} \right) n^2 \leq \left(1 + 2\alpha + 2\sqrt{2(\alpha + \varepsilon)} - \frac{\eta}{3} \right) n^2$$

and

$$\begin{aligned}
4c &= 4(1 - 7\sqrt{\varepsilon})nr = (1 - 7\sqrt{\varepsilon}) \left(1 + 2(\alpha + \varepsilon) + 2\sqrt{2(\alpha + \varepsilon)} \right) n^2 \\
&\geq \left(1 + 2\alpha + 2\sqrt{2(\alpha + \varepsilon)} \right) n^2 - \varepsilon^{\frac{1}{3}} n^2.
\end{aligned}$$

Thus

$$D = b^2 - 4c \leq \left(\varepsilon^{\frac{1}{3}} - \frac{\eta}{3} \right) n^2 < 0$$

since $\varepsilon \ll \eta$. Hence this quadratic has no real zeroes, and $R_r(S, T) - Q_r(S, T) \geq 0$ as desired. This completes the proof. \square

It is now simple to prove that every non-extremal graph G whose minimum degree δ is slightly larger than $n/2$ contains significantly more than $\text{reg}_{\text{even}}(n, \delta)/2$ edge-disjoint Hamilton cycles.

Lemma 24. *Suppose $0 < 1/n \ll c \ll \eta \ll \alpha, 1/2 - \alpha < 1/2$. Let G be a graph on n vertices with $\delta := \delta(G) = (1/2 + \alpha)n$ such that G is not η -extremal. Then G contains at least $\text{reg}_{\text{even}}(n, \delta)/2 + cn$ edge-disjoint Hamilton cycles.*

Proof. Let $\tau_0 := \tau(1/4)$ be as defined in Theorem 9. Choose new constants $\varepsilon, \varepsilon', \nu, \nu', \tau$ such that

$$0 < 1/n \ll \nu', c \ll \varepsilon, \varepsilon' \ll \eta \ll \nu \ll \tau \ll \alpha, 1/2 - \alpha, \tau_0.$$

Let

$$r := \left(\frac{1}{4} + \frac{\alpha + \varepsilon'}{2} + \sqrt{\frac{\alpha + \varepsilon'}{2}} \right) n.$$

By reducing ε' and ε slightly if necessary we may assume that both r and εn are even integers. By Lemmas 17 and 20, G contains an εn -factor H which is a robust

(ν', τ) -expander. Let $G' := G - H$. By Lemma 21, G' is not $(\eta/2)$ -extremal. Since also $\delta(G') = (1/2 + \alpha - \varepsilon)n$, we can apply Lemma 23 with $\varepsilon + \varepsilon'$ and $\alpha - \varepsilon$ playing the roles of ε and α to find an r -factor H' of G' .

Since H is a robust (ν', τ) -expander (and thus also a robust (ν', τ_0) -expander), the same holds for $H + H'$. Hence by Theorem 9, $H + H'$ can be decomposed into $d(H + H')/2$ edge-disjoint Hamilton cycles. By Theorem 4 we have $r \geq \text{reg}_{\text{even}}(n, \delta)$, and so

$$\frac{1}{2}d(H + H') \geq \frac{1}{2}(\text{reg}_{\text{even}}(n, \delta) + \varepsilon n) \geq \frac{1}{2}\text{reg}_{\text{even}}(n, \delta) + cn$$

as desired. \square

6. PROOF OF THEOREMS 3 AND 5

We first combine Lemmas 15 and 24 to prove Theorem 3.

Proof of Theorem 3. Choose $n_0 \in \mathbb{N}$ and an additional constant η such that $1/n_0 \ll \eta \ll \varepsilon$. Define α by $\delta(G) = (1/2 + \alpha)n$. Recall from Section 1 that Theorem 3 was already proved in [12] for the case when $\delta(G) \geq (2 - \sqrt{2} + \varepsilon)n$. So we may assume that $\alpha \leq 3/2 - \sqrt{2} + \varepsilon$ and so $\eta \ll \alpha, 1/2 - \alpha$. Thus we can apply Lemma 15 (if G is η -extremal) or Lemma 24 (if G is not η -extremal) to find $\text{reg}_{\text{even}}(n, \delta(G))/2$ edge-disjoint Hamilton cycles in G . \square

Let G be a graph on n vertices whose minimum degree is not much smaller than $n/2$. Before we can prove Theorem 5, we must first show that either G is a robust expander or it is close to either the complete bipartite graph $K_{n/2, n/2}$ or the disjoint union $K_{n/2} \dot{\cup} K_{n/2}$ of two cliques. The former case corresponds to (i) of Theorem 5, and the latter case corresponds to (ii).

Definition 25. We say that a graph G is ε -close to $K_{n/2, n/2}$ if there exists $A \subseteq V(G)$ with $|A| = \lfloor n/2 \rfloor$ and such that $e(A) \leq \varepsilon n^2$. We say that G is ε -close to $K_{n/2} \dot{\cup} K_{n/2}$ if there exists $A \subseteq V(G)$ with $|A| = \lfloor n/2 \rfloor$ and such that $e(A, \bar{A}) \leq \varepsilon n^2$.

Suppose that G is a graph of minimum degree roughly $n/2$. If G is ε -close to $K_{n/2, n/2}$ then the bipartite subgraph of G induced by A and \bar{A} is almost complete. However, \bar{A} may also contain many edges. If G is ε -close to $K_{n/2} \dot{\cup} K_{n/2}$ then both $G[A]$ and $G[\bar{A}]$ are almost complete.

Lemma 26. *Suppose $0 < 1/n \ll \kappa \ll \nu \ll \tau, \varepsilon < 1$. Let G be a graph on n vertices of minimum degree $\delta := \delta(G) \geq (1/2 - \kappa)n$. Then G satisfies one of the following properties:*

- (i) G is ε -close to $K_{n/2, n/2}$;
- (ii) G is ε -close to $K_{n/2} \dot{\cup} K_{n/2}$;
- (iii) G is a robust (ν, τ) -expander.

Proof. Suppose $S \subseteq V(G)$ with $\tau n \leq |S| \leq (1 - \tau)n$. Our aim is to show that either $RN := RN_\nu(S)$ has size at least $|S| + \nu n$ or that G is close to either $K_{n/2, n/2}$ or $K_{n/2} \dot{\cup} K_{n/2}$. We will split the proof into cases depending on $|S|$.

Case 1: $\tau n \leq |S| \leq (1/2 - \sqrt{\nu})n$.

In this case, we have

$$\delta|S| \leq e'(S, V(G)) = e'(S, RN) + e'(S, \overline{RN}) \leq |S||RN| + \nu n^2 \leq |S||RN| + \nu n \frac{|S|}{\tau},$$

and so $|RN| \geq (1/2 - \kappa - \nu/\tau)n \geq |S| + \nu n$ as desired. (Recall that $e'(A, B)$ denotes the number of ordered pairs (a, b) with $ab \in E(G)$, $a \in A$ and $b \in B$.)

Case 2: $(1/2 + 2\nu)n \leq |S| \leq (1 - \tau)n$.

In this case, we have $RN = V(G)$ and so the result is immediate. Indeed, for all $v \in V(G)$, we have $d(v) \geq (1/2 - \kappa)n$ and so $|N(v) \cap S| \geq (2\nu - \kappa)n \geq \nu n$.

Case 3: $(1/2 - \sqrt{\nu})n \leq |S| \leq (1/2 + 2\nu)n$.

Suppose that $|RN| < |S| + \nu n$. We will first show that either $|S \setminus RN| < \sqrt{\nu}n$ or G is ε -close to $K_{n/2, n/2}$. Suppose $|S \setminus RN| \geq \sqrt{\nu}n$. Then

$$\begin{aligned} |S \setminus RN|(\delta - \nu n) &\leq e(S \setminus RN, \overline{S}) = e(S \setminus RN, \overline{S} \cap RN) + e(S \setminus RN, \overline{S} \setminus RN) \\ &\leq |S \setminus RN||\overline{S} \cap RN| + \nu n^2 \leq |S \setminus RN||\overline{S} \cap RN| + \sqrt{\nu}n|S \setminus RN|, \end{aligned}$$

and so $|\overline{S} \cap RN| \geq \delta - 2\sqrt{\nu}n$. But then together with our assumption that $|RN| < |S| + \nu n$, this implies $|S \cap RN| < 3\sqrt{\nu}n$. Hence $e(S) \leq 3\sqrt{\nu}n^2 + |S|\nu n < 4\sqrt{\nu}n^2$. By adding or removing at most $\sqrt{\nu}n$ arbitrary vertices to or from S , we can form a set A of $\lfloor n/2 \rfloor$ vertices with

$$e(A) < 4\sqrt{\nu}n^2 + \sqrt{\nu}n^2 = 5\sqrt{\nu}n^2 \leq \varepsilon n^2.$$

Thus G is ε -close to $K_{n/2, n/2}$.

We may therefore assume that $|S \setminus RN| < \sqrt{\nu}n$, from which it follows that $|\overline{S} \cap RN| < 2\sqrt{\nu}n$ (by our initial assumption that $|RN| < |S| + \nu n$). We will now show that G is ε -close to $K_{n/2} \dot{\cup} K_{n/2}$. We have $e(S, \overline{S} \cap RN) \leq |S||\overline{S} \cap RN| \leq 2\sqrt{\nu}n^2$, and hence $e(S, \overline{S}) \leq 3\sqrt{\nu}n^2$. As before, by adding or removing at most $\sqrt{\nu}n$ arbitrary vertices to or from S , we can therefore form a set A of $\lfloor n/2 \rfloor$ vertices with $e(A, \overline{A}) \leq e(S, \overline{S}) + \sqrt{\nu}n^2 \leq \varepsilon n^2$. Hence G is ε -close to $K_{n/2} \dot{\cup} K_{n/2}$.

If G is not ε -close to either $K_{n/2, n/2}$ or $K_{n/2} \dot{\cup} K_{n/2}$, we must therefore have $|RN| \geq |S| + \nu n$ for all $S \subseteq V(G)$ with $\tau n \leq |S| \leq (1 - \tau)n$, so that G is a robust (ν, τ) -expander as required. \square

We now have all the tools we need to prove Theorem 5.

Proof of Theorem 5. Let $\tau := \tau(1/4)$ be as defined in Theorem 9. Choose $n_0 \in \mathbb{N}$ and new constants $\varepsilon', \varepsilon'', \nu, \nu'$ such that

$$0 < 1/n_0 \ll \nu' \ll \varepsilon \ll \varepsilon', \varepsilon'' \ll \nu \ll \tau, \eta.$$

Consider any graph G on $n \geq n_0$ vertices as in Theorem 5. Let $\delta := \delta(G)$ and define α by $\delta = (1/2 + \alpha)n$. So $-\varepsilon \leq \alpha \leq \varepsilon$. Let

$$r := \left(\frac{1}{4} + \frac{\alpha + \varepsilon'}{2} + \sqrt{\frac{\alpha + \varepsilon'}{2}} \right) n.$$

By reducing ε' and ε'' slightly if necessary we may assume that both r and $\varepsilon''n$ are even integers.

Suppose that G does not satisfy (i), i.e. $e(X) \geq \eta n^2$ for all $X \subseteq V(G)$ with $|X| = \lfloor n/2 \rfloor$. We claim that G is not $(\eta/4)$ -extremal. To show this, consider any set $B \subseteq V(G)$ with

$$|B| = \left(\frac{1}{2} + \sqrt{\frac{\alpha_+}{2}} \pm \frac{\eta}{4} \right) n.$$

By adding or removing at most $\eta n/2$ arbitrary vertices to and from B , we obtain a set B' with $|B'| = \lfloor n/2 \rfloor$ and such that $e(B') \leq e(B) + \eta n^2/2$. Together with our assumption that (i) does not hold, this implies that

$$e(B) \geq \frac{\eta n^2}{2} \geq \frac{1}{2} \left(\alpha_+ + \sqrt{\frac{\alpha_+}{2}} + \frac{\eta}{4} \right) n|B|.$$

Hence G is not $(\eta/4)$ -extremal.

Suppose moreover that (ii) does not hold, so that G fails to be η -close to $K_{n/2} \dot{\cup} K_{n/2}$. By Lemma 26, it follows that G is a robust (ν, τ) -expander. By Lemma 20, G therefore contains an $\varepsilon''n$ -factor H which is a robust (ν', τ) -expander. Let $G' := G - H$. By Lemma 21, G' is not $(\eta/8)$ -extremal. Since also $\delta(G') = (1/2 + \alpha - \varepsilon'')n$, we can apply Lemma 23 with $\varepsilon' + \varepsilon''$ and $\alpha - \varepsilon''$ playing the roles of ε and α to find an r -factor H' of G' .

Since H is a robust (ν', τ) -expander, the same holds for $H + H'$. Hence by Theorem 9, $H + H'$ can be decomposed into $d(H + H')/2$ edge-disjoint Hamilton cycles. By Theorem 4 (and the fact that $\text{reg}_{\text{even}}(n, \delta) = 0$ if $\delta < n/2$) we have $r \geq \max\{\text{reg}_{\text{even}}(n, \delta), n/8\}$, and so

$$\frac{1}{2}d(H + H') \geq \frac{1}{2} \left(\max\{\text{reg}_{\text{even}}(n, \delta), n/8\} + \varepsilon''n \right) \geq \frac{1}{2} \max\{\text{reg}_{\text{even}}(n, \delta), n/8\} + \varepsilon n,$$

as desired. \square

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