## Note

# The difference and ratio of the fractional matching number and the matching number of graphs 

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#### Abstract

Given a graph $G$, the matching number of $G$, written $\alpha^{\prime}(G)$, is the maximum size of a matching in $G$, and the fractional matching number of $G$, written $\alpha_{f}^{\prime}(G)$, is the maximum size of a fractional matching of $G$. In this paper, we prove that if $G$ is an $n$-vertex connected graph that is neither $K_{1}$ nor $K_{3}$, then $\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G) \leq \frac{n-2}{6}$ and $\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)} \leq \frac{3 n}{2 n+2}$. Both inequalities are sharp, and we characterize the infinite family of graphs where equalities hold.


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## 1. Introduction

For undefined terms, see [5]. Throughout this paper, $n$ will always denote the number of vertices of a given graph. A matching in a graph is a set of pairwise disjoint edges. A perfect matching in a graph $G$ is a matching in which each vertex has an incident edge in the matching; its size must be $n / 2$, where $n=|V(G)|$. A fractional matching of $G$ is a function $\phi: E(G) \rightarrow[0,1]$ such that for each vertex $v, \sum_{e \in \Gamma(v)} \phi(e) \leq 1$, where $\Gamma(v)$ is the set of edges incident to $v$, and the size of a fractional matching $\phi$ is $\sum_{e \in E(G)} \phi(e)$. Given a graph $G$, the matching number of $G$, written $\alpha^{\prime}(G)$, is the maximum size of a matching in $G$, and the fractional matching number of $G$, written $\alpha_{f}^{\prime}(G)$, is the maximum size of a fractional matching of $G$.

Given a fractional matching $\phi$, since $\sum_{e \in \Gamma(v)} \phi(e) \leq 1$ for each vertex $v$, we have that $2 \sum_{e \in E(G)} \phi(e) \leq n$, which implies $\alpha_{f}^{\prime}(G) \leq n / 2$. By viewing every matching as a fractional matching it follows that $\alpha_{f}^{\prime}(G) \geq \alpha^{\prime}(G)$ for every graph $G$, but equality need not hold. For example, the fractional matching number of a $k$-regular graph equals $n / 2$ by setting weight $1 / k$ on each edge, but the matching number of a $k$-regular graph can be much smaller than $n / 2$. Thus it is a natural question to find the largest difference between $\alpha_{f}^{\prime}(G)$ and $\alpha^{\prime}(G)$ in a (connected) graph.

In Sections 3 and 4 , we prove tight upper bounds on $\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G)$ and $\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)}$, respectively, for an $n$-vertex connected graph $G$, and we characterize the infinite family of graphs achieving equality for both results. As corollaries of both results, we have upper bounds on both $\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G)$ and $\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)}$ for an $n$-vertex graph $G$, and we characterize the graphs achieving equality for both bounds.

Our proofs use the famous Berge-Tutte Formula [1] for the matching number as well as its fractional analogue. We also use the fact that there is a fractional matching $\phi$ for which $\sum_{e \in E(G)} \phi(e)=\alpha_{f}^{\prime}(G)$ such that $f(e) \in\{0,1 / 2,1\}$ for every edge $e$, and some refinements of the fact. We can prove both Theorems 6 and 8 with two different techniques, and for the sake of the readers we demonstrate each method in the proofs of Theorems 6 and 8 .

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## 2. Tools

In this section, we introduce the tools we used to prove the main results. To prove Theorem 6, we use Theorems 1 and 2. For a graph $H$, let $o(H)$ denote the number of components of $H$ with an odd number of vertices. Given a graph $G$ and $S \subseteq V(G)$, define the deficiency $\operatorname{def}(S)$ by $\operatorname{def}(S)=o(G-S)-|S|$, and let $\operatorname{def}(G)=\max _{S \subseteq V(G)} \operatorname{def}(S)$. Theorem 1 is the famous Berge-Tutte formula, which is a general version of Tutte's 1-factor Theorem [4].

Theorem 1 ([1]). For any n-vertex graph $G, \alpha^{\prime}(G)=\frac{1}{2}(n-\operatorname{def}(G))$.
For the fractional analogue of the Berge-Tutte formula, let $i(H)$ denote the number of isolated vertices in $H$. Given a graph $G$ and $S \subseteq V(G)$, let $\operatorname{def}_{f}(S)=i(G-S)-|S|$ and $\operatorname{def}_{f}(G)=\max _{S \subseteq V(G)} \operatorname{def}_{f}(S)$. Theorem 2 is the fractional version of the Berge-Tutte Formula. This is also the fractional analogue of Tutte's 1 -Factor Theorem saying that $G$ has a fractional perfect matching if and only if $i(G-S) \leq|S|$ for all $S \subseteq V(G)$ (implicit in Pulleyblank [2]), where a fractional perfect matching is a fractional matching $f$ such that $2 \sum_{e \in E(G)} f(e)=n$.

Theorem 2 ([3] See Theorem 2.2.6). For any n-vertex graph $G, \alpha_{f}^{\prime}(G)=\frac{1}{2}\left(n-\operatorname{def}_{f}(G)\right)$.
When we characterize the equalities in the bounds of Theorems 6 and 8, we need the following proposition. Recall that $G[S]$ is the graph induced by a subset of the vertex set $S$.

Proposition 3 ([3] See Proposition 2.2.2). The following are equivalent for a graph $G$.
(a) G has a fractional perfect matching.
(b) There is a partition $\left\{V_{1}, \ldots, V_{n}\right\}$ of the vertex set $V(G)$ such that, for each $i$, the graph $G\left[V_{i}\right]$ is either $K_{2}$ or Hamiltonian.
(c) There is a partition $\left\{V_{1}, \ldots, V_{n}\right\}$ of the vertex set $V(G)$ such that, for each $i$, the graph $G\left[V_{i}\right]$ is either $K_{2}$ or Hamiltonian graph on an odd number of vertices.

Theorem 4 and Observation 5 are used to prove Theorem 8.
Theorem 4 ([3] See Theorem 2.1.5). For any graph G, there is a fractional matching $f$ for which

$$
\sum_{e \in E(G)} f(e)=\alpha_{f}^{\prime}(G)
$$

such that $f(e) \in\{0,1 / 2,1\}$ for every edge $e$.
Given a fractional matching $f$, an unweighted vertex $v$ is a vertex with $\sum_{e \in \Gamma(v)} f(e)=0$, and a full vertex $v$ is a vertex with $f(v w)=1$ for some vertex $w$. Note that $w$ is also a full vertex. An i-edge $e$ is an edge with $f(e)=i$. Note that the existence of an 1-edge guarantees the existence of two full vertices. A vertex subset $S$ of a graph $G$ is independent if $E(G[S])=\emptyset$, where $G[S]$ is the graph induced by $S$.

Observation 5. Among all the fractional matchings of an n-vertex graph $G$ satisfying the conditions of Theorem 4, let $f$ be a fractional matching with the greatest number of edges e with $f(e)=1$. Then we have the following:
(a) The graph induced by the $\frac{1}{2}$-edges is the union of odd cycles. Furthermore, if $C$ and $C^{\prime}$ are two disjoint cycles in the graph induced by $\frac{1}{2}$-edges, then there is no edge $u u^{\prime}$ such that $u \in V(C)$ and $u^{\prime} \in V\left(C^{\prime}\right)$.
(b) The set $S$ of the unweighted vertices is independent. Furthermore, every unweighted vertex is adjacent only to a full vertex.
(c) $\alpha^{\prime}(G) \geq w_{1}+\sum_{i=1}^{\infty} i c_{i}, \alpha_{f}^{\prime}(G)=w_{1}+\sum_{i=1}^{\infty}\left(\frac{2 i+1}{2}\right) c_{i}$, and $n=w_{0}+2 w_{1}+\sum_{i=1}^{\infty}(2 i+1) c_{i}$, where $w_{0}$, $w_{1}$, and $c_{i}$ are the number of unweighted vertices, the number of 1-edges, and the number of odd cycles of length $2 i+1$ in the graph induced by $\frac{1}{2}$-edges in $G$, respectively.
Proof. (a) The graph induced by the $\frac{1}{2}$-edges cannot have a vertex with degree at least 3 since $\sum_{e \in \Gamma(v)} f(e) \leq 1$ for each vertex $v$. Thus the graph must be a disjoint union of paths or cycles. If the graph contains a path or an even cycle, then by replacing weight $1 / 2$ on each edge on the path or the even cycle with weight 1 and 0 alternatively, we can have a fractional matching with the same fractional matching number and more edges with weight 1 , which contradicts the choice of $f$. Thus the graph induced by the $\frac{1}{2}$-edges is the union of odd cycles. If there is an edge $u v$ such that $u \in V(C)$ and $v \in V\left(C^{\prime}\right)$, where $C$ and $C^{\prime}$ are two different odd cycles induced by some $\frac{1}{2}$-edges, then $f(u v)=0$, since $\sum_{e \in \Gamma(x)} f(e) \leq 1$ for each vertex $x$. By replacing weights 0 and $1 / 2$ on the edge $u v$ and the edges on $C$ and $C^{\prime}$ with weight 1 on $u v$, and 0 and 1 on the edges in $E(C)$ and $E\left(C^{\prime}\right)$ alternatively, not violating the definition of a fractional matching, we have a fractional matching with the same fractional matching number with more edges with weight 1 , which is a contradiction. Thus we have the desired result.
(b) If two unweighted vertices $u$ and $v$ are adjacent, then we can put a positive weight on the edge $u v$, which contradicts the choice of $f$. If there exists an unweighted vertex $x$, which is not incident to any full vertex, then $x$ must be adjacent to a vertex $y$ such that $f\left(y y_{1}\right)=1 / 2$ and $f\left(y y_{2}\right)=1 / 2$ for some vertices $y_{1}$ and $y_{2}$. By replacing the weights $0,1 / 2$, and $1 / 2$ on $x y, y y_{1}$, and $y y_{2}$ with $1,0,0$, respectively, we have a fractional matching with the same fractional matching number with more edges with weight 1 , which is a contradiction.
(c) By the definitions of $w_{0}, w_{1}$, and $c_{i}$, we have the desired result.


Fig. 1. All 5-vertex graphs in Theorems 6(i) and 8(i).


Fig. 2. All graphs in Theorems 6(ii) and 8(ii).

## 3. Sharp upper bound for $\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G)$

What are the structures of the graphs having the maximum difference between the fractional matching number and the matching number in an $n$-vertex connected graph? The graphs may have big fractional matching number and small matching number. So, by the Berge-Tutte Formula and its fractional version, they may have a vertex subset $S$ such that almost all of the odd components of $G-S$ have at least three vertices in order to get $S$ to have small fractional deficiency and big deficiency. This is our idea behind the proof of Theorem 6.

Theorem 6. For $n \geq 5$, if $G$ is a connected graph with $n$ vertices, then $\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G) \leq \frac{n-2}{6}$, and equality holds only when either (i) $n=5$ and either $C_{5}$ is subgraph of $G$ or $K_{2}+K_{3}$ is a subgraph of $G$ (see Fig. 1), or
(ii) $G$ has a vertex $v$ such that the components of $G-v$ are all $K_{3}$ except one single vertex (see Fig. 2).

Proof. Among all the vertex subsets with maximum deficiency, let $S$ be the largest set. By the Berge-Tutte Formula, $\alpha^{\prime}(G)=\frac{1}{2}(n-\operatorname{def}(S))$, and by the choice of $S$, all components of $G-S$ have an odd number of vertices. Let $x$ be the number of isolated vertices of $G-S$, and let $y$ be the number of other components of $G-S$. This implies $n \geq|S|+x+3 y$. If $S=\emptyset$, then $\alpha^{\prime}(G) \in\left\{\frac{n}{2}, \frac{n-1}{2}\right\}$, depending on the parity of $n$. In this case, $\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G) \leq \frac{n}{2}-\frac{n-1}{2}=\frac{1}{2} \leq \frac{n-2}{6}$, since $n \geq 5$. Now, assume that $S$ is non-empty.
Case 1: $x=0$. Since $\operatorname{def}_{f}(G) \geq 0,|S| \geq 1$, and $n \geq|S|+3 y$, we have

$$
\begin{aligned}
\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G) & =\frac{1}{2}\left(n-\operatorname{def}_{f}(G)\right)-\frac{1}{2}(n-\operatorname{def}(S))=\frac{1}{2}\left(\operatorname{def}(S)-\operatorname{def}_{f}(G)\right) \\
& \leq \frac{1}{2}(y-|S|-0) \leq \frac{1}{2}\left(\frac{n-|S|}{3}-|S|\right)=\frac{n-4|S|}{6} \leq \frac{n-4}{6}<\frac{n-2}{6}
\end{aligned}
$$

Case 2: $x \geq 1$. Since $n \geq|S|+x+3 y,|S| \geq 1$, and $x \geq 1$, we have

$$
\begin{aligned}
\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G) & =\frac{1}{2}\left(n-\operatorname{def}_{f}(G)\right)-\frac{1}{2}(n-\operatorname{def}(S))=\frac{1}{2}\left(\operatorname{def}(S)-\operatorname{def}_{f}(G)\right) \\
& \leq \frac{1}{2}(x+y-|S|-(x-|S|)) \leq \frac{y}{2}=\frac{n-x-|S|}{6} \leq \frac{n-2}{6}
\end{aligned}
$$

Equality in the bound requires equality in each step of the computation. When $n=5$, we conclude that (i) follows by Proposition 3. In Case 1, we cannot have equality, and in Case 2, we have $|S|=1, x=1$, and $n=|S|+x+3 y=2+3 y$.

Since $G$ is connected, the components of $G-S$ are $P_{3}$ or $K_{3}$ except only one single vertex. If a component of $G-S$ is a copy of $P_{3}$, then by choosing the central vertex $u$ of the path, we have $\operatorname{def}(S \cup\{u\})=o(G-(S \cup\{u\}))-|S \cup\{u\}|=o(G-S)-|S|$, yet $|S \cup\{u\}|>|S|$, which contradict the choice of $S$. Thus we have the desired result.

Corollary 7. For any $n$-vertex graph $G$, we have $\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G) \leq \frac{n}{6}$, and equality holds only when $G$ is the disjoint union of copies of $K_{3}$.
Proof. First, we show that if $n \leq 4$ and $G$ is connected, then $\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G) \leq \frac{n}{6}$, and equality holds only when $G=K_{3}$. If $n \leq 2$, then $G \in\left\{K_{1}, K_{2}\right\}$, which implies that $\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G)=0<n / 6$. If $n=3$, then $G \in\left\{P_{3}, K_{3}\right\}$. Note that $\alpha_{f}^{\prime}\left(P_{3}\right)-\alpha^{\prime}\left(P_{3}\right)=1-1=0<3 / 6$ and $\alpha_{f}^{\prime}\left(K_{3}\right)-\alpha^{\prime}\left(K_{3}\right)=3 / 2-1=1 / 2 \leq 3 / 6$. Furthermore, equality holds only when $G=K_{3}$. If $n=4$, then either $G=K_{1,3}$ or $G$ contains $P_{4}$ as a subgraph. Since $\alpha_{f}^{\prime}\left(K_{1,3}\right)-\alpha^{\prime}\left(K_{1,3}\right)=1-1=0<4 / 6$ and $\alpha_{f}^{\prime}\left(P_{4}\right)-\alpha^{\prime}\left(P_{4}\right)=2-2=0<4 / 6$, we conclude that for any positive integer $n, \alpha_{f}^{\prime}(G)-\alpha^{\prime}(G) \leq \frac{n}{6}$. In fact, if $n \geq 5$, then by Theorem 6 , the difference must be at most $\frac{n-2}{6}$. Thus, for connected graphs, equality holds only when $G=K_{3}$.

Now, if we assume that $G$ is disconnected, then $G$ is the disjoint union of connected graphs $G_{1}, \ldots, G_{k}$. Let $\left|V\left(G_{i}\right)\right|=n_{i}$ for $i \in[k]$. Since

$$
\begin{aligned}
\alpha_{f}^{\prime}(G)-\alpha^{\prime}(G) & =\left[\alpha_{f}^{\prime}\left(G_{1}\right)+\cdots+\alpha_{f}^{\prime}\left(G_{k}\right)\right]-\left[\alpha^{\prime}\left(G_{1}\right)+\cdots+\alpha^{\prime}\left(G_{k}\right)\right] \\
& =\left[\alpha_{f}^{\prime}\left(G_{1}\right)-\alpha^{\prime}\left(G_{1}\right)\right]+\cdots+\left[\alpha_{f}^{\prime}\left(G_{k}\right)-\alpha^{\prime}\left(G_{k}\right)\right] \leq \frac{n_{1}}{6}+\cdots+\frac{n_{k}}{6}=\frac{n}{6}
\end{aligned}
$$

equality holds only when each $G_{i}$ is a copy of $K_{3}$ for $i \in[k]$.

## 4. Sharp upper bound for $\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)}$

To prove the upper bound of Theorem 8, we still can use the Berge-Tutte formula and its fractional analogue. However, we provide an alternative way to prove the theorem.
Theorem 8. For $n \geq 5$, if $G$ is a connected graph with $n$ vertices, then $\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)} \leq \frac{3 n}{2 n+2}$, and equality holds only when either
(i) $n=5$ and either $C_{5}$ is a subgraph of $G$ or $K_{2}+K_{3}$ is a subgraph of $G$ (see Fig. 1), or
(ii) $G$ has $a$ vertex $v$ such that the components of $G-v$ are all $K_{3}$ except one single vertex (see Fig. 2).

Proof. Among all the fractional matchings of an $n$-vertex graph $G$ with the size equal to $\alpha_{f}^{\prime}(G)$, let $f$ be a fractional matching such that the number of edges $e$ with $f(e)=1$ is maximized. We follow the notation in Observation 5.
Case 1: $w_{0}=w_{1}=0$. Since $G$ is connected and $n \geq 5$, there exists only one $i$ such that $i \geq 2$ and $c_{i}$ is not zero, and $\alpha^{\prime}(G)=i c_{i} \neq 0$. Then we have

$$
\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)} \leq \frac{\left(\frac{2 i+1}{2}\right) c_{i}}{i c_{i}}=1+\frac{1}{2 i} \leq \frac{5}{4}
$$

Case 2: $w_{0} \geq 1$ and $w_{1}=0$. By part (b) of Observation 5, this cannot happen.
Case 3: $w_{0}=0$ and $w_{1} \geq 1$. Since $\sum_{i=1}^{\infty} c_{i} \leq \frac{n-2 w_{1}}{3}$, by part (c) of Observation 5 , we have

$$
\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)} \leq \frac{w_{1}+\sum_{i=1}^{\infty}\left(\frac{2 i+1}{2}\right) c_{i}}{w_{1}+\sum_{i=1}^{\infty} i c_{i}}=\frac{\frac{n-w_{0}}{2}}{\frac{n-w_{0}-\sum_{i=1}^{\infty} c_{i}}{2}}=\frac{n}{n-\sum_{i=1}^{\infty} c_{i}} \leq \frac{n}{n-\frac{n-2 w_{1}}{3}}=\frac{3 n}{2 n+2 w_{1}} \leq \frac{3 n}{2 n+2}
$$

Case 4: $w_{0} \geq 1$ and $w_{1} \geq 1$. Since $\sum_{i=1}^{\infty} c_{i} \leq \frac{n-2 w_{1}-w_{0}}{3}$, by part (c) of Observation 5, we have

$$
\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)} \leq \frac{\frac{n-w_{0}}{2}}{\frac{n-w_{0}-\sum_{i=1}^{\infty} c_{i}}{2}} \leq \frac{n-w_{0}}{n-w_{0}-\frac{n-2 w_{1}-w_{0}}{3}}=\frac{3\left(n-w_{0}\right)}{2\left(n+w_{1}-w_{0}\right)}<\frac{3 n}{2\left(n+w_{1}\right)} \leq \frac{3 n}{2(n+1)}
$$

Equality in the bound requires equality in each step of the computation; we only need to check Case 1 and Case 2 . In Case 1 , we have $i=2$, which means that $n=5$ and $G$ contains a copy of $C_{5}$. In Case 3 , we have $w_{1}=1$ and $\sum_{i=1}^{\infty} c_{i}=\frac{n-2}{3}$, which means that the graph induced by the $\frac{1}{2}$-edges is the union of $K_{3}$. Thus $G$ has $K_{2}+k K_{3}$ as a subgraph for some positive integer $k$. Note that there is an edge between the copy of $K_{2}$ and any copy of $K_{3}$ by part (b) of Observation 5 . Also, there are no edges between any pair of two triangles by part (a) of Observation 5. Let $u$ and $v$ be the two vertices corresponding to the copy of $K_{2}$. If there are two different triangles $C$ and $C^{\prime}$ in $G$ such that $u$ and $v$ are incident to $C$ and $C^{\prime}$, respectively, then we have $\alpha^{\prime}(G)>w_{1}+c_{1}$, which implies that we cannot have equality in the first inequality in Case 3 . Thus, we conclude that $G$ contains a copy of either $K_{2}+K_{3}$ as a subgraph or a vertex $v$ such that the components of $G-v$ are all $K_{3}$ except only one single vertex.

Corollary 9. For any n-vertex graph $G$ with at least one edge, we have $\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)} \leq \frac{3}{2}$, and equality holds only when $G$ is the disjoint union of copies of $K_{3}$.
Proof. By the proof of Corollary 7, if $n \leq 4$ and $G$ is connected, then $\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)} \leq \frac{3}{2}$, and equality holds only when $G=K_{3}$. If we assume that $G$ is disconnected, then $G$ is the disjoint union of connected graphs $G_{1}, \ldots, G_{k}$. Let $\left|V\left(G_{i}\right)\right|=n_{i}$ for $i \in[k]$. Without loss of generality, we may assume that $\frac{\alpha_{f}^{\prime}\left(G_{1}\right)}{\alpha^{\prime}\left(G_{1}\right)} \geq \frac{\alpha_{f}^{\prime}\left(G_{i}\right)}{\alpha^{\prime}\left(G_{i}\right)}$ for all $i \in[k]$. Then we have

$$
\frac{\alpha_{f}^{\prime}(G)}{\alpha^{\prime}(G)}=\frac{\alpha_{f}^{\prime}\left(G_{1}\right)+\cdots+\alpha_{f}^{\prime}\left(G_{k}\right)}{\alpha^{\prime}\left(G_{1}\right)+\cdots+\alpha^{\prime}\left(G_{k}\right)} \leq \frac{\alpha_{f}^{\prime}\left(G_{1}\right)}{\alpha^{\prime}\left(G_{1}\right)} \leq \frac{3}{2}
$$

and equality holds only when each $G_{i}$ is a copy of $K_{3}$ for $i \in[k]$.

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