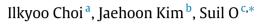
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# The difference and ratio of the fractional matching number and the matching number of graphs

ABSTRACT



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#### 1. Introduction

For undefined terms, see [5]. Throughout this paper, *n* will always denote the number of vertices of a given graph. A *matching* in a graph is a set of pairwise disjoint edges. A *perfect matching* in a graph *G* is a matching in which each vertex has an incident edge in the matching; its size must be n/2, where n = |V(G)|. A *fractional matching* of *G* is a function  $\phi : E(G) \rightarrow [0, 1]$  such that for each vertex v,  $\sum_{e \in \Gamma(v)} \phi(e) \le 1$ , where  $\Gamma(v)$  is the set of edges incident to v, and the *size* of a fractional matching  $\phi$  is  $\sum_{e \in E(G)} \phi(e)$ . Given a graph *G*, the *matching number* of *G*, written  $\alpha'(G)$ , is the maximum size of a matching in *G*, and the *fractional matching number* of *G*, written  $\alpha'_f(G)$ , is the maximum size of a fractional matching of *G*.

Given a fractional matching  $\phi$ , since  $\sum_{e \in \Gamma(v)} \phi(e) \leq 1$  for each vertex v, we have that  $2 \sum_{e \in E(G)} \phi(e) \leq n$ , which implies  $\alpha'_f(G) \leq n/2$ . By viewing every matching as a fractional matching it follows that  $\alpha'_f(G) \geq \alpha'(G)$  for every graph G, but equality need not hold. For example, the fractional matching number of a k-regular graph equals n/2 by setting weight 1/k on each edge, but the matching number of a k-regular graph can be much smaller than n/2. Thus it is a natural question to find the largest difference between  $\alpha'_f(G)$  and  $\alpha'(G)$  in a (connected) graph.

In Sections 3 and 4, we prove tight upper bounds on  $\alpha'_f(G) - \alpha'(G)$  and  $\frac{\alpha'_f(G)}{\alpha'(G)}$ , respectively, for an *n*-vertex connected graph *G*, and we characterize the infinite family of graphs achieving equality for both results. As corollaries of both results, we have upper bounds on both  $\alpha'_f(G) - \alpha'(G)$  and  $\frac{\alpha'_f(G)}{\alpha'(G)}$  for an *n*-vertex graph *G*, and we characterize the graphs achieving equality for both bounds.

Our proofs use the famous Berge–Tutte Formula [1] for the matching number as well as its fractional analogue. We also use the fact that there is a fractional matching  $\phi$  for which  $\sum_{e \in E(G)} \phi(e) = \alpha'_f(G)$  such that  $f(e) \in \{0, 1/2, 1\}$  for every edge e, and some refinements of the fact. We can prove both Theorems 6 and 8 with two different techniques, and for the sake of the readers we demonstrate each method in the proofs of Theorems 6 and 8.

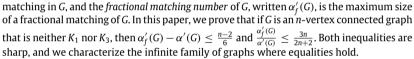
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Note

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Given a graph G, the matching number of G, written  $\alpha'(G)$ , is the maximum size of a

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### 2. Tools

In this section, we introduce the tools we used to prove the main results. To prove Theorem 6, we use Theorems 1 and 2. For a graph *H*, let o(H) denote the number of components of *H* with an odd number of vertices. Given a graph *G* and  $S \subseteq V(G)$ , define the *deficiency* def(*S*) by def(*S*) = o(G - S) - |S|, and let def(*G*) =  $\max_{S \subseteq V(G)} def(S)$ . Theorem 1 is the famous Berge–Tutte formula, which is a general version of Tutte's 1-factor Theorem [4].

## **Theorem 1** ([1]). For any *n*-vertex graph G, $\alpha'(G) = \frac{1}{2}(n - \text{def}(G))$ .

For the fractional analogue of the Berge–Tutte formula, let i(H) denote the number of isolated vertices in H. Given a graph G and  $S \subseteq V(G)$ , let def<sub>f</sub>(S) = i(G - S) - |S| and def<sub>f</sub>(G) = max<sub>S \subseteq V(G)</sub> def<sub>f</sub>(S). Theorem 2 is the fractional version of the Berge–Tutte Formula. This is also the fractional analogue of Tutte's 1-Factor Theorem saying that G has a fractional perfect matching if and only if  $i(G - S) \leq |S|$  for all  $S \subseteq V(G)$  (implicit in Pulleyblank [2]), where a fractional perfect matching is a fractional matching f such that  $2 \sum_{e \in E(G)} f(e) = n$ .

**Theorem 2** ([3] See Theorem 2.2.6). For any n-vertex graph G,  $\alpha'_f(G) = \frac{1}{2}(n - \text{def}_f(G))$ .

When we characterize the equalities in the bounds of Theorems 6 and 8, we need the following proposition. Recall that G[S] is the graph induced by a subset of the vertex set *S*.

Proposition 3 ([3] See Proposition 2.2.2). The following are equivalent for a graph G.

(a) *G* has a fractional perfect matching.

(b) There is a partition  $\{V_1, \ldots, V_n\}$  of the vertex set V(G) such that, for each i, the graph  $G[V_i]$  is either  $K_2$  or Hamiltonian.

(c) There is a partition  $\{V_1, \ldots, V_n\}$  of the vertex set V(G) such that, for each *i*, the graph  $G[V_i]$  is either  $K_2$  or Hamiltonian graph on an odd number of vertices.

Theorem 4 and Observation 5 are used to prove Theorem 8.

**Theorem 4** ([3] See Theorem 2.1.5). For any graph G, there is a fractional matching f for which

$$\sum_{e \in E(G)} f(e) = \alpha'_f(G)$$

such that  $f(e) \in \{0, 1/2, 1\}$  for every edge e.

Given a fractional matching f, an *unweighted* vertex v is a vertex with  $\sum_{e \in \Gamma(v)} f(e) = 0$ , and a *full* vertex v is a vertex with f(vw) = 1 for some vertex w. Note that w is also a full vertex. An *i-edge* e is an edge with f(e) = i. Note that the existence of an 1-edge guarantees the existence of two full vertices. A vertex subset S of a graph G is *independent* if  $E(G[S]) = \emptyset$ , where G[S] is the graph induced by S.

**Observation 5.** Among all the fractional matchings of an n-vertex graph *G* satisfying the conditions of *Theorem* 4, let *f* be a fractional matching with the greatest number of edges e with f(e) = 1. Then we have the following:

(a) The graph induced by the  $\frac{1}{2}$ -edges is the union of odd cycles. Furthermore, if C and C' are two disjoint cycles in the graph induced by  $\frac{1}{2}$ -edges, then there is no edge uu' such that  $u \in V(C)$  and  $u' \in V(C')$ .

(b) The set S of the unweighted vertices is independent. Furthermore, every unweighted vertex is adjacent only to a full vertex.

(c)  $\alpha'(G) \ge w_1 + \sum_{i=1}^{\infty} ic_i, \alpha'_f(G) = w_1 + \sum_{i=1}^{\infty} (\frac{2i+1}{2})c_i$ , and  $n = w_0 + 2w_1 + \sum_{i=1}^{\infty} (2i+1)c_i$ , where  $w_0, w_1$ , and  $c_i$  are the number of unweighted vertices, the number of 1-edges, and the number of odd cycles of length 2i + 1 in the graph induced by  $\frac{1}{2}$ -edges in *G*, respectively.

**Proof.** (a) The graph induced by the  $\frac{1}{2}$ -edges cannot have a vertex with degree at least 3 since  $\sum_{e \in \Gamma(v)} f(e) \leq 1$  for each vertex v. Thus the graph must be a disjoint union of paths or cycles. If the graph contains a path or an even cycle, then by replacing weight 1/2 on each edge on the path or the even cycle with weight 1 and 0 alternatively, we can have a fractional matching with the same fractional matching number and more edges with weight 1, which contradicts the choice of f. Thus the graph induced by the  $\frac{1}{2}$ -edges is the union of odd cycles. If there is an edge uv such that  $u \in V(C)$  and  $v \in V(C')$ , where C and C' are two different odd cycles induced by some  $\frac{1}{2}$ -edges, then f(uv) = 0, since  $\sum_{e \in \Gamma(x)} f(e) \leq 1$  for each vertex x. By replacing weights 0 and 1/2 on the edge uv and the edges on C and C' with weight 1 on uv, and 0 and 1 on the edges in E(C) and E(C') alternatively, not violating the definition of a fractional matching, we have a fractional matching with the same fractional matching number with more edges with weight 1, which is a contradiction. Thus we have the desired result.

(b) If two unweighted vertices u and v are adjacent, then we can put a positive weight on the edge uv, which contradicts the choice of f. If there exists an unweighted vertex x, which is not incident to any full vertex, then x must be adjacent to a vertex y such that  $f(yy_1) = 1/2$  and  $f(yy_2) = 1/2$  for some vertices  $y_1$  and  $y_2$ . By replacing the weights 0, 1/2, and 1/2 on xy,  $yy_1$ , and  $yy_2$  with 1, 0, 0, respectively, we have a fractional matching with the same fractional matching number with more edges with weight 1, which is a contradiction.

(c) By the definitions of  $w_0$ ,  $w_1$ , and  $c_i$ , we have the desired result.  $\Box$ 

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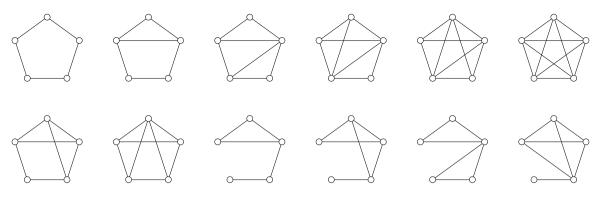


Fig. 1. All 5-vertex graphs in Theorems 6(i) and 8(i).

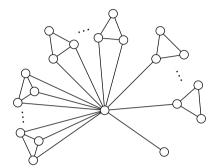


Fig. 2. All graphs in Theorems 6(ii) and 8(ii).

## 3. Sharp upper bound for $\alpha'_f(G) - \alpha'(G)$

What are the structures of the graphs having the maximum difference between the fractional matching number and the matching number in an *n*-vertex connected graph? The graphs may have big fractional matching number and small matching number. So, by the Berge–Tutte Formula and its fractional version, they may have a vertex subset *S* such that almost all of the odd components of G - S have at least three vertices in order to get *S* to have small fractional deficiency and big deficiency. This is our idea behind the proof of Theorem 6.

**Theorem 6.** For  $n \ge 5$ , if *G* is a connected graph with *n* vertices, then  $\alpha'_f(G) - \alpha'(G) \le \frac{n-2}{6}$ , and equality holds only when either (i) n = 5 and either  $C_5$  is subgraph of *G* or  $K_2 + K_3$  is a subgraph of *G* (see Fig. 1), or (ii) *C* has a current *n* with that the common state of *G* and *G* with *n* and *n* and *G* with *n* and *n* and *G* with *n* and *n* and *G* with *n* and *G* with

(ii) *G* has a vertex v such that the components of G - v are all  $K_3$  except one single vertex (see Fig. 2).

**Proof.** Among all the vertex subsets with maximum deficiency, let *S* be the largest set. By the Berge–Tutte Formula,  $\alpha'(G) = \frac{1}{2}(n - \text{def}(S))$ , and by the choice of *S*, all components of *G* – *S* have an odd number of vertices. Let *x* be the number of isolated vertices of *G* – *S*, and let *y* be the number of other components of *G* – *S*. This implies  $n \ge |S| + x + 3y$ . If  $S = \emptyset$ , then  $\alpha'(G) \in \{\frac{n}{2}, \frac{n-1}{2}\}$ , depending on the parity of *n*. In this case,  $\alpha'_f(G) - \alpha'(G) \le \frac{n}{2} - \frac{n-1}{2} = \frac{1}{2} \le \frac{n-2}{6}$ , since  $n \ge 5$ . Now, assume that *S* is non-empty.

*Case* 1: x = 0. Since def<sub>f</sub>(G)  $\geq 0$ ,  $|S| \geq 1$ , and  $n \geq |S| + 3y$ , we have

$$\begin{aligned} \alpha_f'(G) - \alpha'(G) &= \frac{1}{2}(n - \operatorname{def}_f(G)) - \frac{1}{2}(n - \operatorname{def}(S)) = \frac{1}{2}(\operatorname{def}(S) - \operatorname{def}_f(G)) \\ &\leq \frac{1}{2}(y - |S| - 0) \leq \frac{1}{2}\left(\frac{n - |S|}{3} - |S|\right) = \frac{n - 4|S|}{6} \leq \frac{n - 4}{6} < \frac{n - 2}{6}. \end{aligned}$$

*Case* 2:  $x \ge 1$ . Since  $n \ge |S| + x + 3y$ ,  $|S| \ge 1$ , and  $x \ge 1$ , we have

$$\begin{aligned} \alpha_f'(G) - \alpha'(G) &= \frac{1}{2}(n - \operatorname{def}_f(G)) - \frac{1}{2}(n - \operatorname{def}(S)) = \frac{1}{2}(\operatorname{def}(S) - \operatorname{def}_f(G)) \\ &\leq \frac{1}{2}\left(x + y - |S| - (x - |S|)\right) \le \frac{y}{2} = \frac{n - x - |S|}{6} \le \frac{n - 2}{6}. \end{aligned}$$

Equality in the bound requires equality in each step of the computation. When n = 5, we conclude that (i) follows by Proposition 3. In Case 1, we cannot have equality, and in Case 2, we have |S| = 1, x = 1, and n = |S| + x + 3y = 2 + 3y.

Since *G* is connected, the components of G - S are  $P_3$  or  $K_3$  except only one single vertex. If a component of G - S is a copy of  $P_3$ , then by choosing the central vertex *u* of the path, we have def $(S \cup \{u\}) = o(G - (S \cup \{u\})) - |S \cup \{u\}| = o(G - S) - |S|$ , yet  $|S \cup \{u\}| > |S|$ , which contradict the choice of *S*. Thus we have the desired result.  $\Box$ 

**Corollary 7.** For any *n*-vertex graph *G*, we have  $\alpha'_f(G) - \alpha'(G) \le \frac{n}{6}$ , and equality holds only when *G* is the disjoint union of copies of *K*<sub>3</sub>.

**Proof.** First, we show that if  $n \le 4$  and *G* is connected, then  $\alpha'_f(G) - \alpha'(G) \le \frac{n}{6}$ , and equality holds only when  $G = K_3$ . If  $n \le 2$ , then  $G \in \{K_1, K_2\}$ , which implies that  $\alpha'_f(G) - \alpha'(G) = 0 < n/6$ . If n = 3, then  $G \in \{P_3, K_3\}$ . Note that  $\alpha'_f(P_3) - \alpha'(P_3) = 1 - 1 = 0 < 3/6$  and  $\alpha'_f(K_3) - \alpha'(K_3) = 3/2 - 1 = 1/2 \le 3/6$ . Furthermore, equality holds only when  $G = K_3$ . If n = 4, then either  $G = K_{1,3}$  or *G* contains  $P_4$  as a subgraph. Since  $\alpha'_f(K_{1,3}) - \alpha'(K_{1,3}) = 1 - 1 = 0 < 4/6$  and  $\alpha'_f(P_4) - \alpha'(P_4) = 2 - 2 = 0 < 4/6$ , we conclude that for any positive integer n,  $\alpha'_f(G) - \alpha'(G) \le \frac{n}{6}$ . In fact, if  $n \ge 5$ , then by Theorem 6, the difference must be at most  $\frac{n-2}{2}$ . Thus, for connected graphs, equality holds only when  $C = K_3$ .

by Theorem 6, the difference must be at most  $\frac{n-2}{6}$ . Thus, for connected graphs, equality holds only when  $G = K_3$ . Now, if we assume that G is disconnected, then G is the disjoint union of connected graphs  $G_1, \ldots, G_k$ . Let  $|V(G_i)| = n_i$  for  $i \in [k]$ . Since

$$\alpha'_{f}(G) - \alpha'(G) = \left[\alpha'_{f}(G_{1}) + \dots + \alpha'_{f}(G_{k})\right] - \left[\alpha'(G_{1}) + \dots + \alpha'(G_{k})\right]$$
$$= \left[\alpha'_{f}(G_{1}) - \alpha'(G_{1})\right] + \dots + \left[\alpha'_{f}(G_{k}) - \alpha'(G_{k})\right] \le \frac{n_{1}}{6} + \dots + \frac{n_{k}}{6} = \frac{n}{6},$$

equality holds only when each  $G_i$  is a copy of  $K_3$  for  $i \in [k]$ .  $\Box$ 

## 4. Sharp upper bound for $\frac{\alpha'_f(G)}{\alpha'(G)}$

To prove the upper bound of Theorem 8, we still can use the Berge–Tutte formula and its fractional analogue. However, we provide an alternative way to prove the theorem.

**Theorem 8.** For  $n \ge 5$ , if *G* is a connected graph with *n* vertices, then  $\frac{\alpha'_f(G)}{\alpha'(G)} \le \frac{3n}{2n+2}$ , and equality holds only when either (i) n = 5 and either  $C_5$  is a subgraph of *G* or  $K_2 + K_3$  is a subgraph of *G* (see Fig. 1), or

(ii) *G* has a vertex v such that the components of G - v are all  $K_3$  except one single vertex (see Fig. 2).

**Proof.** Among all the fractional matchings of an *n*-vertex graph *G* with the size equal to  $\alpha'_f(G)$ , let *f* be a fractional matching such that the number of edges *e* with f(e) = 1 is maximized. We follow the notation in Observation 5.

*Case* 1:  $w_0 = w_1 = 0$ . Since *G* is connected and  $n \ge 5$ , there exists only one *i* such that  $i \ge 2$  and  $c_i$  is not zero, and  $\alpha'(G) = ic_i \ne 0$ . Then we have

$$\frac{\alpha'_f(G)}{\alpha'(G)} \le \frac{\left(\frac{2i+1}{2}\right)c_i}{ic_i} = 1 + \frac{1}{2i} \le \frac{5}{4}.$$

*Case* 2:  $w_0 \ge 1$  and  $w_1 = 0$ . By part (b) of Observation 5, this cannot happen. *Case* 3:  $w_0 = 0$  and  $w_1 \ge 1$ . Since  $\sum_{i=1}^{\infty} c_i \le \frac{n-2w_1}{3}$ , by part (c) of Observation 5, we have

 $\frac{\alpha'_f(G)}{\alpha'(G)} \leq \frac{w_1 + \sum_{i=1}^{\infty} \left(\frac{2i+1}{2}\right)c_i}{w_1 + \sum_{i=1}^{\infty} ic_i} = \frac{\frac{n-w_0}{2}}{\frac{n-w_0 - \sum_{i=1}^{\infty} c_i}{2}} = \frac{n}{n - \sum_{i=1}^{\infty} c_i} \leq \frac{n}{n - \frac{n-2w_1}{3}} = \frac{3n}{2n + 2w_1} \leq \frac{3n}{2n + 2}.$ 

Case 4:  $w_0 \ge 1$  and  $w_1 \ge 1$ . Since  $\sum_{i=1}^{\infty} c_i \le \frac{n-2w_1-w_0}{3}$ , by part (c) of Observation 5, we have

$$\frac{\alpha'_f(G)}{\alpha'(G)} \leq \frac{\frac{n-w_0}{2}}{\frac{n-w_0-\sum\limits_{i=1}^{\infty}c_i}{\frac{1}{2}}} \leq \frac{n-w_0}{n-w_0-\frac{n-2w_1-w_0}{3}} = \frac{3(n-w_0)}{2(n+w_1-w_0)} < \frac{3n}{2(n+w_1)} \leq \frac{3n}{2(n+1)}.$$

Equality in the bound requires equality in each step of the computation; we only need to check Case 1 and Case 2. In Case 1, we have i = 2, which means that n = 5 and G contains a copy of  $C_5$ . In Case 3, we have  $w_1 = 1$  and  $\sum_{i=1}^{\infty} c_i = \frac{n-2}{3}$ , which means that the graph induced by the  $\frac{1}{2}$ -edges is the union of  $K_3$ . Thus G has  $K_2 + kK_3$  as a subgraph for some positive integer k. Note that there is an edge between the copy of  $K_2$  and any copy of  $K_3$  by part (b) of Observation 5. Also, there are no edges between any pair of two triangles by part (a) of Observation 5. Let u and v be the two vertices corresponding to the copy of  $K_2$ . If there are two different triangles C and C' in G such that u and v are incident to C and C', respectively, then we have  $\alpha'(G) > w_1 + c_1$ , which implies that we cannot have equality in the first inequality in Case 3. Thus, we conclude that G contains a copy of either  $K_2 + K_3$  as a subgraph or a vertex v such that the components of G - v are all  $K_3$  except only one single vertex.  $\Box$ 

**Corollary 9.** For any *n*-vertex graph *G* with at least one edge, we have  $\frac{\alpha'_f(G)}{\alpha'(G)} \leq \frac{3}{2}$ , and equality holds only when *G* is the disjoint union of copies of  $K_3$ .

**Proof.** By the proof of Corollary 7, if  $n \le 4$  and *G* is connected, then  $\frac{\alpha'_f(G)}{\alpha'(G)} \le \frac{3}{2}$ , and equality holds only when  $G = K_3$ . If we assume that *G* is disconnected, then *G* is the disjoint union of connected graphs  $G_1, \ldots, G_k$ . Let  $|V(G_i)| = n_i$  for  $i \in [k]$ . Without loss of generality, we may assume that  $\frac{\alpha'_f(G_1)}{\alpha'(G_1)} \ge \frac{\alpha'_f(G_i)}{\alpha'(G_i)}$  for all  $i \in [k]$ . Then we have

$$\frac{\alpha'_f(G)}{\alpha'(G)} = \frac{\alpha'_f(G_1) + \dots + \alpha'_f(G_k)}{\alpha'(G_1) + \dots + \alpha'(G_k)} \le \frac{\alpha'_f(G_1)}{\alpha'(G_1)} \le \frac{3}{2},$$

and equality holds only when each  $G_i$  is a copy of  $K_3$  for  $i \in [k]$ .  $\Box$ 

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