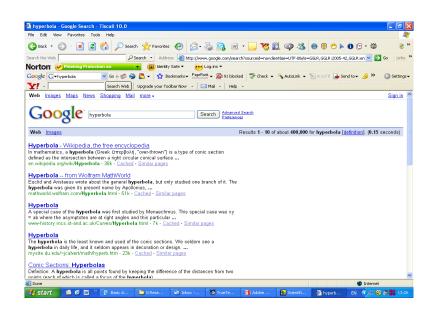
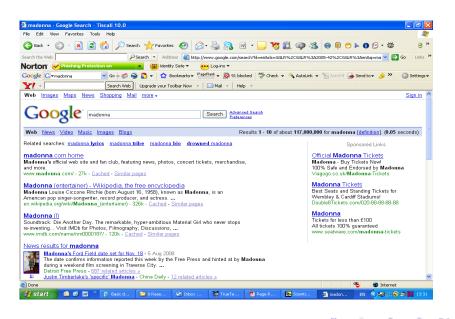
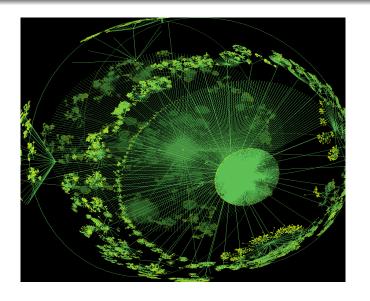
Google: An Application of Linear Algebra (The mathematics of a great success)

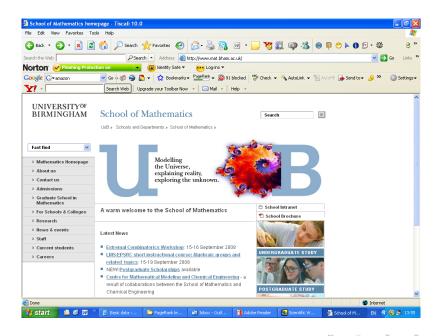
Peter Butkovic

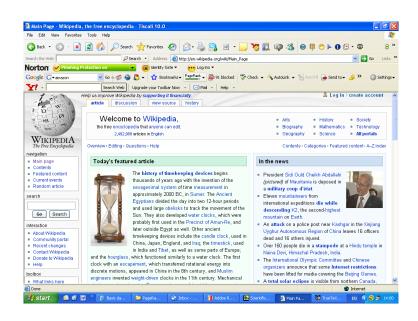






Hyun's Map of the Web





• Google ($^{\sim}$ googol = 10^{100})

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- Founders of Google: Larry Page and Sergey Brin

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- 1995: Research students at Stanford
- 1996: Started a student project on search engines
- 1998: Google incorporates as a company (initial investment: \$1.1million) and files patent for PageRank

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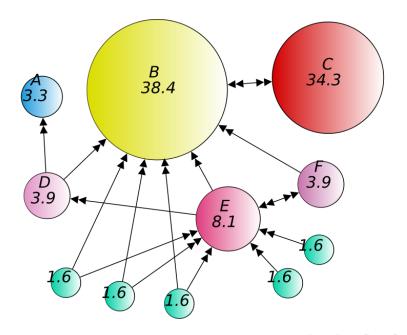
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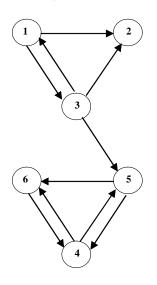
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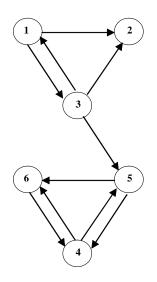
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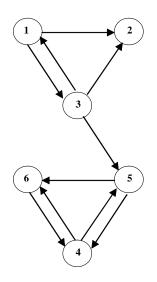


•
$$r(P_1) = \frac{1}{3}r(P_3)$$



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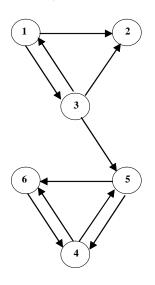
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$$r(P_2) = \frac{1}{2}r(P_1) + \frac{1}{3}r(P_3)$$



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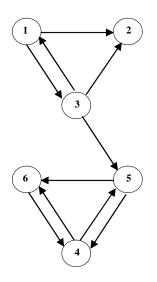


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$$r(P_1) = \frac{1}{3}r(P_3)$$

•
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$$r(P_3) = \frac{1}{2}r(P_1)$$

•
$$r(P_4) = \frac{1}{2}r(P_5) + r(P_6)$$



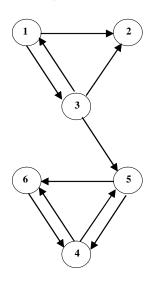
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$$r(P_1) = \frac{1}{3}r(P_3)$$

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$$r(P_4) = \frac{1}{2}r(P_5) + r(P_6)$$

•
$$r(P_5) = \frac{1}{3}r(P_3) + \frac{1}{2}r(P_4)$$



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$$r(P_1) = \frac{1}{3}r(P_3)$$

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$$r(P_3) = \frac{1}{2}r(P_1)$$

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$$r(P_4) = \frac{1}{2}r(P_5) + r(P_6)$$

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$$r(P_5) = \frac{1}{3}r(P_3) + \frac{1}{2}r(P_4)$$

•
$$r(P_6) = \frac{1}{2}r(P_4) + \frac{1}{2}r(P_5)$$

•
$$r_0 = \left(\frac{1}{n}, \frac{1}{n}, \ldots\right)$$

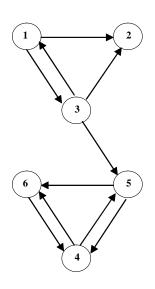
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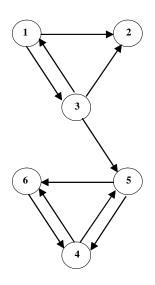
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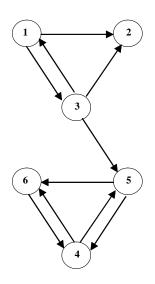


• $r_0 = (\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6})$



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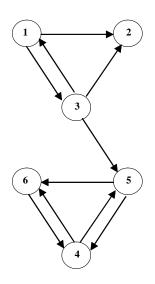
•
$$r_1(P_1) = \frac{1}{3}r_0(P_3) = \frac{1}{18}$$



•
$$r_0 = (\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6})$$

•
$$r_1(P_1) = \frac{1}{3}r_0(P_3) = \frac{1}{18}$$

•
$$r_1(P_2) = \frac{1}{2}r_0(P_1) + \frac{1}{3}r_0(P_3) = \frac{5}{36}$$

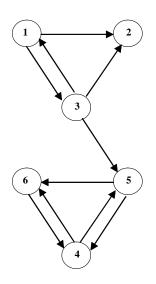


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$$r_0 = (\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6})$$

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$$r_1(P_1) = \frac{1}{3}r_0(P_3) = \frac{1}{18}$$

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• $r_1(P_3) = \frac{1}{2}r_0(P_1) = \frac{1}{12}$



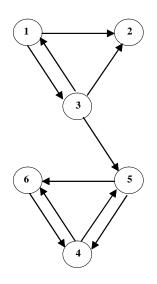
•
$$r_0 = (\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6})$$

•
$$r_1(P_1) = \frac{1}{3}r_0(P_3) = \frac{1}{18}$$

•
$$r_1(P_2) = \frac{1}{2}r_0(P_1) + \frac{1}{3}r_0(P_3) = \frac{5}{36}$$

•
$$r_1(P_3) = \frac{1}{2}r_0(P_1) = \frac{1}{12}$$

•
$$r_1(P_4) = \frac{1}{2}r_0(P_5) + r_0(P_6) = \frac{1}{4}$$



•
$$r_0 = (\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6})$$

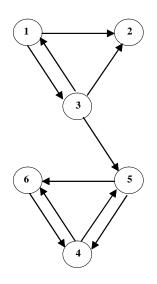
•
$$r_1(P_1) = \frac{1}{3}r_0(P_3) = \frac{1}{18}$$

•
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•
$$r_1(P_3) = \frac{1}{2}r_0(P_1) = \frac{1}{12}$$

•
$$r_1(P_4) = \frac{1}{2}r_0(P_5) + r_0(P_6) = \frac{1}{4}$$

•
$$r_1(P_5) = \frac{1}{3}r_0(P_3) + \frac{1}{2}r_0(P_4) = \frac{5}{36}$$



•
$$r_0 = (\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6})$$

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$$r_1(P_1) = \frac{1}{3}r_0(P_3) = \frac{1}{18}$$

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•
$$r_1(P_4) = \frac{1}{2}r_0(P_5) + r_0(P_6) = \frac{1}{4}$$

•
$$r_1(P_5) = \frac{1}{3}r_0(P_3) + \frac{1}{2}r_0(P_4) = \frac{5}{36}$$

•
$$r_1(P_6) = \frac{1}{2}r_0(P_4) + \frac{1}{2}r_0(P_5) = \frac{1}{6}$$

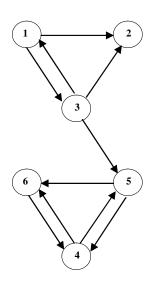
	r_0	r_1
$\overline{P_1}$	1/6	1/18
P_2	1/6	5/36
P_3	1/6	1/12
P_4	1/6	1/4
P_5	1/6	5/36
P_6	1/6	1/6

	r ₀	r_1	rank
$\overline{P_1}$	1/6	1/18	6
P_2	1/6	5/36	3 - 4
P_3	1/6	1/12	5
P_4	1/6	1/4	1
P_5	1/6	5/36	3 - 4
P_6	1/6	1/6	2

	<i>r</i> ₀	r_1	rank	<i>r</i> ₂
P_1	1/6	1/18	6	1/36
P_2	1/6	5/36	3 - 4	1/18
P_3	1/6	1/12	5	1/36
P_4	1/6	1/4	1	17/72
P_5	1/6	5/36	3 - 4	11/72
P_6	1/6	1/6	2	14/72

	r_0	r_1	rank	<i>r</i> ₂	rank
P_1	1/6	1/18	6	1/36	5 – 6
P_2	1/6	5/36	3 - 4	1/18	4
P_3	1/6	1/12	5	1/36	5 - 6
P_4	1/6	1/4	1	17/72	1
P_5	1/6	5/36	3 - 4	11/72	3
P_6	1/6	1/6	2	1/36 1/18 1/36 17/72 11/72 14/72	2

	r_0	r_1	rank	<i>r</i> ₂	rank
P_1	1/6	1/18	6	1/36	5 – 6
P_2	1/6	5/36	3 - 4	1/18	4
P_3	1/6	1/12	5	1/36	5 - 6
P_4	1/6	1/4	1	17/72	1
P_5	1/6	5/36	3 - 4	11/72	3
P_6	1/6	1/6	2	1/36 1/18 1/36 17/72 11/72 14/72	2



Hyperlink matrix H:

	P_1	P_2	P_3	P_4	P_5	P_6
P_1	0	$\frac{1}{2}$	$\frac{1}{2}$	0	0	0
P_2	0	0	Ō	0	0	0
P_3	$\frac{1}{3}$	$\frac{1}{3}$	0	0	$\frac{1}{3}$	0
P_4	Ö	Ö	0	0	$\frac{1}{2}$	$\frac{1}{2}$
P_5	0	0	0	$\frac{1}{2}$	Ō	$\frac{1}{2}$ $\frac{1}{2}$
P_6	0	0	0	$\bar{1}$	0	Ō

Stochastic matrix: Every row is ≥ 0 and sums to 1.

H could be a stochastic matrix if it was not for the rows corresponding to the dangling nodes

$$\underbrace{\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)}_{r_0} \cdot \underbrace{\left(\begin{array}{cccc} 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0\\ 0 & 0 & 0 & 0 & 0 & 0\\ \frac{1}{3} & \frac{1}{3} & 0 & 0 & \frac{1}{3} & 0\\ 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2}\\ 0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2}\\ 0 & 0 & 0 & 1 & 0 & 0 \end{array}\right)}_{r_1}$$

$$= \underbrace{\left(\frac{1}{18}, \frac{5}{36}, \frac{1}{12}, \frac{1}{4}, \frac{5}{36}, \frac{1}{6}\right)}_{r_1}$$

$$\underbrace{\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)}_{r_0} \cdot \underbrace{\left(\begin{array}{cccc} 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 \\ \frac{1}{3} & \frac{1}{3} & 0 & 0 & \frac{1}{3} & 0 \\ 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\ 0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\ 0 & 0 & 0 & 1 & 0 & 0 \end{array}\right)}_{r_1}$$

$$= \underbrace{\left(\frac{1}{18}, \frac{5}{36}, \frac{1}{12}, \frac{1}{4}, \frac{5}{36}, \frac{1}{6}\right)}_{r_1}$$

$$r_1 = r_0 \cdot H$$

$$\underbrace{\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)}_{r_0} \cdot \underbrace{\left(\begin{array}{cccc} 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0\\ 0 & 0 & 0 & 0 & 0 & 0\\ \frac{1}{3} & \frac{1}{3} & 0 & 0 & \frac{1}{3} & 0\\ 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2}\\ 0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2}\\ 0 & 0 & 0 & 1 & 0 & 0 \end{array}\right)}_{r_1}$$

$$= \underbrace{\left(\frac{1}{18}, \frac{5}{36}, \frac{1}{12}, \frac{1}{4}, \frac{5}{36}, \frac{1}{6}\right)}_{r_1}$$

$$r_1 = r_0 \cdot H$$

Similarly:

$$r_2 = r_1 \cdot H = (r_0 \cdot H) \cdot H = r_0 \cdot H^2$$

$$r_3 = r_2 \cdot H = ... = r_0 \cdot H^3$$

$$r_3 = r_2 \cdot H = ... = r_0 \cdot H^3$$

• In general:

•

$$r_k = r_{k-1} \cdot H = \dots = r_0 \cdot H^k$$

$$r_3 = r_2 \cdot H = ... = r_0 \cdot H^3$$

• In general:

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$$r_k = r_{k-1} \cdot H = \dots = r_0 \cdot H^k$$

• ... "Power Method"

• Will this power method converge? If not what conditions must be satisfied?

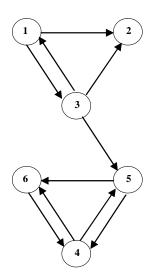
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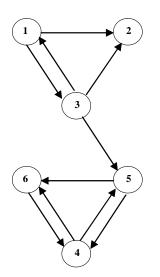
- Will this power method converge? If not what conditions must be satisfied?
- If it converges, will it do so to a vector meaningful for page ranking?
- Does the convergence depend on the starting vector?
- Does the limit depend on the starting vector?
- If it converges, how long does it take?

Page and Brin first had to deal with a number of problems: Rank sink pages



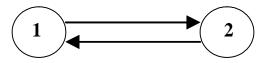
	1		
	<i>r</i> ₀	r_1	 r_{13}
$\overline{P_1}$	1/6	1/18	 0
P_2	1/6	5/36	 0
P_3	1/6	1/12	 0
P_4	1/6	1/4	 2/3
P_5	1/6	5/36	 1/3
P_6	1/6	1/6	 1/5

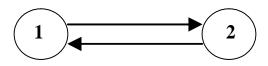
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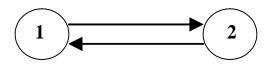
	1		
	<i>r</i> ₀	r_1	 r_{13}
P_1	1/6	1/18	 0
P_2	1/6	5/36	 0
P_3	1/6	1/12	 0
P_4	1/6	1/4	 2/3
P_5	1/6	5/36	 1/3
P_6	1/6	1/6	 1/5

• Nodes 4,5,6 are a *link* farm.



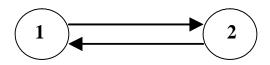


$$\bullet \ H = \left(\begin{array}{cc} 0 & 1 \\ 1 & 0 \end{array}\right)$$



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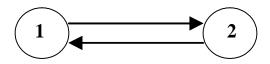
•
$$(1,0)\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = (0,1)$$



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$$(1,0)\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = (0,1)$$

•
$$(0,1)\begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} = (1,0)$$

 \bullet Flip-flop \rightarrow no convergence

Elements of the Markov chains theory:

• If H is a stochastic matrix and x_0 a stochastic vector then the sequence

$$\{x_0, x_1 = x_0H, x_2 = x_1H, x_3 = x_2H, ...\}$$

is called a Markov chain.

Elements of the Markov chains theory:

 If H is a stochastic matrix and x₀ a stochastic vector then the sequence

$$\{x_0, x_1 = x_0 H, x_2 = x_1 H, x_3 = x_2 H, ...\}$$

is called a Markov chain.

• *H* is called the *transition probability matrix*.

Theorem (Markov, 1906)

If H is a **positive** transition probability matrix of a Markov chain then this chain converges to a unique positive vector (called stationary vector) independently of the starting vector. • Brin and Page: Adjustments to the basic model using the concept of a *random surfer*.

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- **Adjustment 1** (*stochasticity*): Zero rows in *H* (corresponding to the dangling nodes) are replaced by

$$\left(\frac{1}{n},\frac{1}{n},...,\frac{1}{n}\right)$$

The new matrix is S.

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The new matrix is S.

• Our example:

$$H o S = \left(egin{array}{cccccc} 0 & rac{1}{2} & rac{1}{2} & 0 & 0 & 0 \\ rac{1}{16} & rac{1}{16} & rac{1}{16} & rac{1}{16} & rac{1}{16} & rac{1}{16} \\ rac{1}{3} & rac{1}{3} & 0 & 0 & rac{1}{3} & 0 \\ 0 & 0 & 0 & 0 & rac{1}{2} & rac{1}{2} \\ 0 & 0 & 0 & rac{1}{2} & 0 & rac{1}{2} \\ 0 & 0 & 0 & 1 & 0 & 0 \end{array}
ight)$$

- Brin and Page: Adjustments to the basic model using the concept of a *random surfer*.
- Adjustment 1 (stochasticity): Zero rows in H (corresponding to the dangling nodes) are replaced by

$$\left(\frac{1}{n},\frac{1}{n},...,\frac{1}{n}\right)$$

The new matrix is S.

• Our example:

$$H \to S = \left(\begin{array}{ccccc} 0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0\\ \frac{1}{6} & \frac{1}{6} & \frac{1}{6} & \frac{1}{6} & \frac{1}{6} & \frac{1}{6}\\ \frac{1}{3} & \frac{1}{3} & 0 & 0 & \frac{1}{3} & 0\\ 0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2}\\ 0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2}\\ 0 & 0 & 0 & 1 & 0 & 0 \end{array}\right)$$

• *S* is stochastic! (But not yet positive)



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$$E = \left(\begin{array}{ccc} \frac{1}{n} & \cdots & \frac{1}{n} \\ \vdots & \ddots & \vdots \\ \frac{1}{n} & \cdots & \frac{1}{n} \end{array}\right)$$

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• and $\alpha \in (0,1)$ controls the proportion of the time RS follows the *hyperlinks* as opposed to *teleportation*

• $G = \alpha S + (1 - \alpha) E$ is called the *Google matrix*: it is stochastic and positive!

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- Hence any Markov chain with *G* is guaranteed to converge to a unique positive vector, independently of the starting vector.

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- Hence any Markov chain with *G* is guaranteed to converge to a unique positive vector, independently of the starting vector.
- Actually used $\alpha \approx .85$

• Our example for $\alpha = 0.9$:

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Google's PageRank vector is the stationary vector of G which is

 $\left(.03721,.05396,.04151,.3751,.206,.2802\right)$

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- Random surfer spends 3.721% of their time on P_1 , etc...
- The importance ranking therefore is: P_4 , P_6 , P_5 , P_2 , P_3 , P_1 .

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- Does the limit depend on the starting vector? NO
- If it converges, how long does it take?

$$\bullet \ r_k \cdot G = r_{k+1}$$

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- In general: $x \cdot M = \lambda x$
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- The stationary point of any Markov chain with transition matrix G is an eigenvector of G corresponding to the eigenvalue 1

Perron-Frobenius theory of Linear Algebra solves the eigenvector-eigenvalue problem for non-negative matrices.

Theorem (Perron, 1912)

If G is a positive, stochastic matrix then $\lambda=1$ is an eigenvalue of G and G has a unique positive eigenvector (up to multiples).

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- Total: $\approx 50n$ operations (recall n > 8 billion)
- CONCLUSION: The power method with Google matrix is very fast!

THANK YOU

Next session in this room at 12.00: "Careers, degrees and mathematics"