# Google: An Application of Linear Algebra (The mathematics of a great success) 

Peter Butkovic





## Hyun's Map of the Web




- Google ( ${ }^{\sim}$ googol $=10^{100}$ )
- Google ( ${ }^{\text {googol }}=10^{100}$ )
- Founders of Google: Larry Page and Sergey Brin
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- Founders of Google: Larry Page and Sergey Brin
- 1995: Research students at Stanford
- 1996: Started a student project on search engines
- 1998: Google incorporates as a company (initial investment: \$1.1million) and files patent for PageRank
- 2000: Selling advertisements began
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- 2001: Patent granted
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- 2004: Total capital reached \$23billion
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- 2008: Google uses 450,000 servers in 25 locations around the world to index $>8$ billion websites


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- Each webpage is assigned a measure of importance ... PageRank


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- Each webpage is assigned a measure of importance ... PageRank
- A hyperlink is a "recommendation": A webpage is important if other webpages point to it (via hyperlinks)
- The PageRank of the recommender matters (the higher the better)
- The generosity of the recommender matters (the higher the worse)

A.N.Langville and C.D.Meyer: Google's PageRank and Beyond (PUP 2006)


$$
\text { - } r\left(P_{1}\right)=\frac{1}{3} r\left(P_{3}\right)
$$

Node 2 - dangling
A.N.Langville and C.D.Meyer: Google's PageRank and Beyond (PUP 2006)


- $r\left(P_{1}\right)=\frac{1}{3} r\left(P_{3}\right)$
- $r\left(P_{2}\right)=$ $\frac{1}{2} r\left(P_{1}\right)+\frac{1}{3} r\left(P_{3}\right)$

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- $r\left(P_{3}\right)=\frac{1}{2} r\left(P_{1}\right)$

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- $r\left(P_{3}\right)=\frac{1}{2} r\left(P_{1}\right)$
- $r\left(P_{4}\right)=\frac{1}{2} r\left(P_{5}\right)+r\left(P_{6}\right)$

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- $r\left(P_{1}\right)=\frac{1}{3} r\left(P_{3}\right)$
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- $r\left(P_{5}\right)=$ $\frac{1}{3} r\left(P_{3}\right)+\frac{1}{2} r\left(P_{4}\right)$

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- $r\left(P_{3}\right)=\frac{1}{2} r\left(P_{1}\right)$
- $r\left(P_{4}\right)=\frac{1}{2} r\left(P_{5}\right)+r\left(P_{6}\right)$
- $r\left(P_{5}\right)=$ $\frac{1}{3} r\left(P_{3}\right)+\frac{1}{2} r\left(P_{4}\right)$
- $r\left(P_{6}\right)=$ $\frac{1}{2} r\left(P_{4}\right)+\frac{1}{2} r\left(P_{5}\right)$

Node 2 - dangling

- $r_{0}=\left(\frac{1}{n}, \frac{1}{n}, \ldots\right)$
- $r_{0}=\left(\frac{1}{n}, \frac{1}{n}, \ldots\right)$
- $r_{0} \longrightarrow r_{1}$
- $r_{0}=\left(\frac{1}{n}, \frac{1}{n}, \ldots\right)$
- $r_{0} \longrightarrow r_{1}$
- $r_{1} \longrightarrow r_{2}$
- $r_{0}=\left(\frac{1}{n}, \frac{1}{n}, \ldots\right)$
- $r_{0} \longrightarrow r_{1}$
- $r_{1} \longrightarrow r_{2}$
- ...
- $r_{0}=\left(\frac{1}{n}, \frac{1}{n}, \ldots\right)$
- $r_{0} \longrightarrow r_{1}$
- $r_{1} \longrightarrow r_{2}$
- ...
- $r_{k} \longrightarrow r_{k+1}$
- $r_{0}=\left(\frac{1}{n}, \frac{1}{n}, \ldots\right)$
- $r_{0} \longrightarrow r_{1}$
- $r_{1} \longrightarrow r_{2}$
- ...
- $r_{k} \longrightarrow r_{k+1}$
- ...

- $r_{0}=\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)$


$$
\begin{aligned}
& \text { - } r_{0}=\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right) \\
& \text { - } r_{1}\left(P_{1}\right)=\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{1}{18}
\end{aligned}
$$



$$
\begin{aligned}
& \text { - } r_{0}=\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right) \\
& \text { - } r_{1}\left(P_{1}\right)=\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{1}{18}
\end{aligned}
$$

- $r_{1}\left(P_{2}\right)=$ $\frac{1}{2} r_{0}\left(P_{1}\right)+\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{5}{36}$

- $r_{0}=\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)$
- $r_{1}\left(P_{1}\right)=\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{1}{18}$
- $r_{1}\left(P_{2}\right)=$ $\frac{1}{2} r_{0}\left(P_{1}\right)+\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{5}{36}$
- $r_{1}\left(P_{3}\right)=\frac{1}{2} r_{0}\left(P_{1}\right)=\frac{1}{12}$


$$
\begin{aligned}
& \text { - } r_{0}=\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right) \\
& \text { - } r_{1}\left(P_{1}\right)=\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{1}{18} \\
& \text { - } r_{1}\left(P_{2}\right)= \\
& \frac{1}{2} r_{0}\left(P_{1}\right)+\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{5}{36} \\
& \text { - } r_{1}\left(P_{3}\right)=\frac{1}{2} r_{0}\left(P_{1}\right)=\frac{1}{12} \\
& \text { - } r_{1}\left(P_{4}\right)= \\
& \frac{1}{2} r_{0}\left(P_{5}\right)+r_{0}\left(P_{6}\right)=\frac{1}{4}
\end{aligned}
$$



- $r_{0}=\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)$
- $r_{1}\left(P_{1}\right)=\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{1}{18}$
- $r_{1}\left(P_{2}\right)=$ $\frac{1}{2} r_{0}\left(P_{1}\right)+\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{5}{36}$
- $r_{1}\left(P_{3}\right)=\frac{1}{2} r_{0}\left(P_{1}\right)=\frac{1}{12}$
- $r_{1}\left(P_{4}\right)=$ $\frac{1}{2} r_{0}\left(P_{5}\right)+r_{0}\left(P_{6}\right)=\frac{1}{4}$
- $r_{1}\left(P_{5}\right)=$ $\frac{1}{3} r_{0}\left(P_{3}\right)+\frac{1}{2} r_{0}\left(P_{4}\right)=\frac{5}{36}$

- $r_{0}=\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)$
- $r_{1}\left(P_{1}\right)=\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{1}{18}$
- $r_{1}\left(P_{2}\right)=$ $\frac{1}{2} r_{0}\left(P_{1}\right)+\frac{1}{3} r_{0}\left(P_{3}\right)=\frac{5}{36}$
- $r_{1}\left(P_{3}\right)=\frac{1}{2} r_{0}\left(P_{1}\right)=\frac{1}{12}$
- $r_{1}\left(P_{4}\right)=$ $\frac{1}{2} r_{0}\left(P_{5}\right)+r_{0}\left(P_{6}\right)=\frac{1}{4}$
- $r_{1}\left(P_{5}\right)=$ $\frac{1}{3} r_{0}\left(P_{3}\right)+\frac{1}{2} r_{0}\left(P_{4}\right)=\frac{5}{36}$
- $r_{1}\left(P_{6}\right)=$ $\frac{1}{2} r_{0}\left(P_{4}\right)+\frac{1}{2} r_{0}\left(P_{5}\right)=\frac{1}{6}$

|  | $r_{0}$ | $r_{1}$ |
| :---: | :---: | :---: |
| $P_{1}$ | $1 / 6$ | $1 / 18$ |
| $P_{2}$ | $1 / 6$ | $5 / 36$ |
| $P_{3}$ | $1 / 6$ | $1 / 12$ |
| $P_{4}$ | $1 / 6$ | $1 / 4$ |
| $P_{5}$ | $1 / 6$ | $5 / 36$ |
| $P_{6}$ | $1 / 6$ | $1 / 6$ |


|  | $r_{0}$ | $r_{1}$ | rank |
| :---: | :---: | :---: | :---: |
| $P_{1}$ | $1 / 6$ | $1 / 18$ | 6 |
| $P_{2}$ | $1 / 6$ | $5 / 36$ | $3-4$ |
| $P_{3}$ | $1 / 6$ | $1 / 12$ | 5 |
| $P_{4}$ | $1 / 6$ | $1 / 4$ | 1 |
| $P_{5}$ | $1 / 6$ | $5 / 36$ | $3-4$ |
| $P_{6}$ | $1 / 6$ | $1 / 6$ | 2 |


|  | $r_{0}$ | $r_{1}$ | rank | $r_{2}$ |
| :---: | :---: | :---: | :---: | :---: |
| $P_{1}$ | $1 / 6$ | $1 / 18$ | 6 | $1 / 36$ |
| $P_{2}$ | $1 / 6$ | $5 / 36$ | $3-4$ | $1 / 18$ |
| $P_{3}$ | $1 / 6$ | $1 / 12$ | 5 | $1 / 36$ |
| $P_{4}$ | $1 / 6$ | $1 / 4$ | 1 | $17 / 72$ |
| $P_{5}$ | $1 / 6$ | $5 / 36$ | $3-4$ | $11 / 72$ |
| $P_{6}$ | $1 / 6$ | $1 / 6$ | 2 | $14 / 72$ |


|  | $r_{0}$ | $r_{1}$ | rank | $r_{2}$ | rank |
| :---: | :---: | :---: | :---: | :---: | :---: |
| $P_{1}$ | $1 / 6$ | $1 / 18$ | 6 | $1 / 36$ | $5-6$ |
| $P_{2}$ | $1 / 6$ | $5 / 36$ | $3-4$ | $1 / 18$ | 4 |
| $P_{3}$ | $1 / 6$ | $1 / 12$ | 5 | $1 / 36$ | $5-6$ |
| $P_{4}$ | $1 / 6$ | $1 / 4$ | 1 | $17 / 72$ | 1 |
| $P_{5}$ | $1 / 6$ | $5 / 36$ | $3-4$ | $11 / 72$ | 3 |
| $P_{6}$ | $1 / 6$ | $1 / 6$ | 2 | $14 / 72$ | 2 |


|  | $r_{0}$ | $r_{1}$ | rank | $r_{2}$ | rank |
| :---: | :---: | :---: | :---: | :---: | :---: |
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| $P_{2}$ | $1 / 6$ | $5 / 36$ | $3-4$ | $1 / 18$ | 4 |
| $P_{3}$ | $1 / 6$ | $1 / 12$ | 5 | $1 / 36$ | $5-6$ |
| $P_{4}$ | $1 / 6$ | $1 / 4$ | 1 | $17 / 72$ | 1 |
| $P_{5}$ | $1 / 6$ | $5 / 36$ | $3-4$ | $11 / 72$ | 3 |
| $P_{6}$ | $1 / 6$ | $1 / 6$ | 2 | $14 / 72$ | 2 |



Hyperlink matrix $H$ :

|  | $P_{1}$ | $P_{2}$ | $P_{3}$ | $P_{4}$ | $P_{5}$ | $P_{6}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $P_{1}$ | 0 | $\frac{1}{2}$ | $\frac{1}{2}$ | 0 | 0 | 0 |
| $P_{2}$ | 0 | 0 | 0 | 0 | 0 | 0 |
| $P_{3}$ | $\frac{1}{3}$ | $\frac{1}{3}$ | 0 | 0 | $\frac{1}{3}$ | 0 |
| $P_{4}$ | 0 | 0 | 0 | 0 | $\frac{1}{2}$ | $\frac{1}{2}$ |
| $P_{5}$ | 0 | 0 | 0 | $\frac{1}{2}$ | 0 | $\frac{1}{2}$ |
| $P_{6}$ | 0 | 0 | 0 | 1 | 0 | 0 |

Stochastic matrix: Every row is $\geq 0$ and sums to 1 .
$H$ could be a stochastic matrix if it was not for the rows corresponding to the dangling nodes

$$
\begin{aligned}
-\underbrace{\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)}_{r_{0}} & \cdot\left(\begin{array}{cccccc}
0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 \\
0 & \frac{1}{3} & 0 & 0 & 0 & 0 \\
\frac{1}{3} & \frac{1}{3} & 0 & 0 & 0 & \frac{1}{3} \\
0 & 0 & 0 \\
0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\
0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\
0 & 0 & 0 & 1 & 0 & 0
\end{array}\right) \\
& =\underbrace{\left(\frac{1}{18}, \frac{5}{36}, \frac{1}{12}, \frac{1}{4}, \frac{5}{36}, \frac{1}{6}\right)}_{n_{1}}
\end{aligned}
$$

$$
\begin{aligned}
& \text { - } \underbrace{\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)}_{0} \cdot\left(\begin{array}{llllll}
0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 \\
\frac{1}{3} & \frac{1}{3} & 0 & 0 & \frac{1}{3} & 0 \\
0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\
0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\
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\end{array}\right) \\
& =\underbrace{\left(\frac{1}{18}, \frac{5}{36}, \frac{1}{12}, \frac{1}{4}, \frac{5}{36}, \frac{1}{6}\right)}_{1}
\end{aligned}
$$

$$
r_{1}=r_{0} \cdot H
$$

$$
\begin{aligned}
& \underbrace{\left(\frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}, \frac{1}{6}\right)}_{r_{0}} \cdot
\end{aligned} \cdot\left(\begin{array}{cccccc}
0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 \\
0 & 0 & 0 & 0 & 0 & 0 \\
\frac{1}{3} & \frac{1}{3} & 0 & 0 & \frac{1}{3} & 0 \\
0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\
0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\
0 & 0 & 0 & 1 & 0 & 0
\end{array}\right)
$$

$$
r_{1}=r_{0} \cdot H
$$

- Similarly:

$$
r_{2}=r_{1} \cdot H=\left(r_{0} \cdot H\right) \cdot H=r_{0} \cdot H^{2}
$$

- 

$$
r_{3}=r_{2} \cdot H=\ldots=r_{0} \cdot H^{3}
$$

- 

$$
r_{3}=r_{2} \cdot H=\ldots=r_{0} \cdot H^{3}
$$

- In general:

$$
r_{k}=r_{k-1} \cdot H=\ldots=r_{0} \cdot H^{k}
$$

- 

$$
r_{3}=r_{2} \cdot H=\ldots=r_{0} \cdot H^{3}
$$

- In general:

$$
r_{k}=r_{k-1} \cdot H=\ldots=r_{0} \cdot H^{k}
$$

- ... "Power Method"


## QUESTIONS

- Will this power method converge? If not what conditions must be satisfied?


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- If it converges, will it do so to a vector meaningful for page ranking?
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- Does the limit depend on the starting vector?


## QUESTIONS

- Will this power method converge? If not what conditions must be satisfied?
- If it converges, will it do so to a vector meaningful for page ranking?
- Does the convergence depend on the starting vector?
- Does the limit depend on the starting vector?
- If it converges, how long does it take?

Page and Brin first had to deal with a number of problems:
Rank sink pages


|  | $r_{0}$ | $r_{1}$ | $\ldots$ | $r_{13}$ |
| :---: | :---: | :---: | :---: | :---: |
| $P_{1}$ | $1 / 6$ | $1 / 18$ | $\ldots$ | 0 |
| $P_{2}$ | $1 / 6$ | $5 / 36$ | $\ldots$ | 0 |
| $P_{3}$ | $1 / 6$ | $1 / 12$ | $\ldots$ | 0 |
| $P_{4}$ | $1 / 6$ | $1 / 4$ | $\ldots$ | $2 / 3$ |
| $P_{5}$ | $1 / 6$ | $5 / 36$ | $\ldots$ | $1 / 3$ |
| $P_{6}$ | $1 / 6$ | $1 / 6$ | $\ldots$ | $1 / 5$ |

Page and Brin first had to deal with a number of problems:
Rank sink pages


|  | $r_{0}$ | $r_{1}$ | $\ldots$ | $r_{13}$ |
| :---: | :---: | :---: | :---: | :---: |
| $P_{1}$ | $1 / 6$ | $1 / 18$ | $\ldots$ | 0 |
| $P_{2}$ | $1 / 6$ | $5 / 36$ | $\ldots$ | 0 |
| $P_{3}$ | $1 / 6$ | $1 / 12$ | $\ldots$ | 0 |
| $P_{4}$ | $1 / 6$ | $1 / 4$ | $\ldots$ | $2 / 3$ |
| $P_{5}$ | $1 / 6$ | $5 / 36$ | $\ldots$ | $1 / 3$ |
| $P_{6}$ | $1 / 6$ | $1 / 6$ | $\ldots$ | $1 / 5$ |

- Nodes 4,5,6 are a link farm.
- Another problem: Cycles

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- $H=\left(\begin{array}{ll}0 & 1 \\ 1 & 0\end{array}\right)$
- Another problem: Cycles

- $H=\left(\begin{array}{ll}0 & 1 \\ 1 & 0\end{array}\right)$
- $(1,0)\left(\begin{array}{ll}0 & 1 \\ 1 & 0\end{array}\right)=(0,1)$
- Another problem: Cycles

- $H=\left(\begin{array}{ll}0 & 1 \\ 1 & 0\end{array}\right)$
- $(1,0)\left(\begin{array}{ll}0 & 1 \\ 1 & 0\end{array}\right)=(0,1)$
- $(0,1)\left(\begin{array}{ll}0 & 1 \\ 1 & 0\end{array}\right)=(1,0)$
- Another problem: Cycles

- $H=\left(\begin{array}{ll}0 & 1 \\ 1 & 0\end{array}\right)$
- $(1,0)\left(\begin{array}{ll}0 & 1 \\ 1 & 0\end{array}\right)=(0,1)$
- $(0,1)\left(\begin{array}{ll}0 & 1 \\ 1 & 0\end{array}\right)=(1,0)$
- Flip-flop $\rightarrow$ no convergence

Elements of the Markov chains theory:

- If $H$ is a stochastic matrix and $x_{0}$ a stochastic vector then the sequence

$$
\left\{x_{0}, x_{1}=x_{0} H, x_{2}=x_{1} H, x_{3}=x_{2} H, \ldots\right\}
$$

is called a Markov chain.

Elements of the Markov chains theory:

- If $H$ is a stochastic matrix and $x_{0}$ a stochastic vector then the sequence

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$$

is called a Markov chain.

- $H$ is called the transition probability matrix.


## Theorem (Markov, 1906)

If $H$ is a positive transition probability matrix of a Markov chain then this chain converges to a unique positive vector (called stationary vector) independently of the starting vector.

- Brin and Page: Adjustments to the basic model using the concept of a random surfer.
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- Adjustment 1 (stochasticity): Zero rows in $H$ (corresponding to the dangling nodes) are replaced by

$$
\left(\frac{1}{n}, \frac{1}{n}, \ldots, \frac{1}{n}\right)
$$

The new matrix is $S$.

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$$

The new matrix is $S$.

- Our example:

$$
H \rightarrow S=\left(\begin{array}{cccccc}
0 & \frac{1}{2} & \frac{1}{2} & 0 & 0 & 0 \\
\frac{1}{6} & \frac{1}{6} & \frac{1}{6} & \frac{1}{6} & \frac{1}{6} & \frac{1}{6} \\
\frac{1}{3} & \frac{1}{3} & 0 & 0 & \frac{1}{3} & 0 \\
0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\
0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\
0 & 0 & 0 & 1 & 0 & 0
\end{array}\right)
$$

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\frac{1}{3} & \frac{1}{3} & 0 & 0 & \frac{1}{3} & 0 \\
0 & 0 & 0 & 0 & \frac{1}{2} & \frac{1}{2} \\
0 & 0 & 0 & \frac{1}{2} & 0 & \frac{1}{2} \\
0 & 0 & 0 & 1 & 0 & 0
\end{array}\right)
$$

- $S$ is stochastic! (But not yet positive)
- Adjustment 2 (primitivity): The random surfer (RS) from time to time moves to a new webpage without using a hyperlink (he "teleports") and then he resumes hyperlink surfing again.
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- 

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S \rightarrow G=\alpha S+(1-\alpha) E
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- where

$$
E=\left(\begin{array}{ccc}
\frac{1}{n} & \cdots & \frac{1}{n} \\
\vdots & \ddots & \vdots \\
\frac{1}{n} & \cdots & \frac{1}{n}
\end{array}\right)
$$

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- and $\alpha \in(0,1)$ controls the proportion of the time RS follows the hyperlinks as opposed to teleportation
- $G=\alpha S+(1-\alpha) E$ is called the Google matrix: it is stochastic and positive!
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- Hence any Markov chain with $G$ is guaranteed to converge to a unique positive vector, independently of the starting vector.
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- Hence any Markov chain with $G$ is guaranteed to converge to a unique positive vector, independently of the starting vector.
- Actually used $\alpha \approx .85$
- Our example for $\alpha=0.9$ :
- Our example for $\alpha=0.9$ :

$$
-G=0.9 S+0.1 E=\frac{1}{60}\left(\begin{array}{cccccc}
1 & 28 & 28 & 1 & 1 & 1 \\
10 & 10 & 10 & 10 & 10 & 10 \\
19 & 19 & 1 & 1 & 19 & 1 \\
1 & 1 & 1 & 1 & 28 & 28 \\
1 & 1 & 1 & 28 & 1 & 28 \\
1 & 1 & 1 & 55 & 1 & 1
\end{array}\right)
$$

- Google's PageRank vector is the stationary vector of $G$ which is
(.03721, . $05396, .04151, .3751, .206, .2802$ )
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- Random surfer spends $3.721 \%$ of their time on $P_{1}$, etc...
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- Random surfer spends $3.721 \%$ of their time on $P_{1}$, etc...
- The importance ranking therefore is: $P_{4}, P_{6}, P_{5}, P_{2}, P_{3}, P_{1}$.
- Will this iterative process converge? YES
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- If it converges, will it do so to a vector meaningful for page ranking? YES
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- If it converges, will it do so to a vector meaningful for page ranking? YES
- Does the convergence depend on the starting vector? NO
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- Does the limit depend on the starting vector? NO
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- If it converges, will it do so to a vector meaningful for page ranking? YES
- Does the convergence depend on the starting vector? NO
- Does the limit depend on the starting vector? NO
- If it converges, how long does it take?
- $r_{k} \cdot G=r_{k+1}$
- $r_{k} \cdot G=r_{k+1}$
- If $r_{k} \rightarrow r$ then $r \cdot G=r$
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- In general: $x \cdot M=\lambda x$
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- $r_{k} \cdot G=r_{k+1}$
- If $r_{k} \rightarrow r$ then $r \cdot G=r$
- In general: $x \cdot M=\lambda x$
- $\lambda \ldots$ eigenvalue of $M$
- $x$... eigenvector of $M$ (if $x \neq 0$ ) corresponding to the eigenvalue $\lambda$
- The stationary point of any Markov chain with transition matrix $G$ is an eigenvector of $G$ corresponding to the eigenvalue 1

Perron-Frobenius theory of Linear Algebra solves the eigenvector-eigenvalue problem for non-negative matrices.

## Theorem (Perron, 1912)

If $G$ is a positive, stochastic matrix then $\lambda=1$ is an eigenvalue of $G$ and $G$ has a unique positive eigenvector (up to multiples).

- If $G$ is a Google matrix - the power method needs roughly 50 iterations
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- In each iteration $\approx n$ arithmetic operations
- Total: $\approx 50 n$ operations (recall $n>8$ billion)
- CONCLUSION: The power method with Google matrix is very fast!


## THANK YOU

Next session in this room at 12.00: "Careers, degrees and mathematics"

