Typical distances in a geometric model for complex networks [*]

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Abstract

The theme of this paper is the study of *typical* distances in a random graph model that was introduced by Krioukov et al. [25] and envisages basic properties of complex networks as the expression of an underlying hyperbolic geometry. This model gives rise to sparse random graphs on the hyperbolic plane which exhibit power law degree distribution as well as local clustering (i.e., they are locally dense). In this paper, we show that in fact these random graphs are *ultra-small* worlds. When the parameters of the model yield a power law degree distribution with exponent between 2 and 3, we show that the distance between two given vertices conditional on belonging to the same component is proportional to $\log \log N$ a.a.s., where N is the number of vertices of the random graph. To be more precise, we show that the distance rescaled by $\log \log N$ converges *in probability* to a certain constant that depends on the exponent of the power law. This constant actually emerges in the same setting in the *Chung-Lu* model.

We also consider the regime where the exponent of the power law is larger than 3. There we show that most pairs of vertices that belong to the same component are within distance $poly(\log \log N)$ with high probability asymptotically as $N \to \infty$.

1 Introduction

The *small-world problem* was first stated by Stanley Milgram in his 1967 paper [26] through which he gave strong evidence of the so-called *small-world effect*. The simplest formulation of the small-world problem [26] is: "Starting with any two people in the world, what is the probability that they will know each other?". A more sophisticated formulation of the problem asks whether any two people, if they do not directly know of each other, have common acquaintances. Milgram's experiment indicated that this is indeed the case within a relatively small random sample of the population of the United States. In particular, it turned out the at least half of the sample was within *six degrees of separation* from the "target" individual. In graph theoretic terms, in the graph of acquaintances the nodes that represent these individuals are within distance 6 from the node that was representing the target individual.

The small-world phenomenon is ubiquitous in natural and technological networks such as neural networks, the Internet, the World-Wide-Web or the power grid – see the book of Chung and Lu [14] as well as the book of Dorogovtsev [10] for experimental evidence regarding such networks. For example, it was announced relatively recently that between any two active users of Facebook there are 3.74 degrees of separation on average [31].

There have been numerous attempts to explain this phenomenon through the theory of complex networks. Among the initial attempts was the "small-world" model of Watts and Strogatz which is defined through random re-wiring of the edges of a cyclic lattice. This model exhibits small average distance, but lacks a basic feature of such large self-organizing networks which is the *scale freeness*. Experimental evidence [1] suggests that these networks have a distribution of degrees whose tail decays like a power law with exponent usually between 2 and 3.

Of course the term "small-world" itself is somewhat vague. Loosely speaking, the term refers to average distances that are slowly growing functions of the number of vertices of the network. A possible candidate is the logarithmic function. Thus, the classical Erdős-Rényi random graph may be thought of as a small-world graph as it has logarithmic diameter – see [4]. However, it lacks the scale freeness as well and, furthermore, it represents a very homogeneous network. This is a very unrealistic feature as most large scale networks contain vertices that have very different properties

from each other. Sub-logarithmic bounds on the diameter were established for the preferential attachment model [2] by Bollobás and Riordan [6]. As it was shown by Bollobás et al. [7], this is scale-free with exponent equal to 3.

Recent research that focused on models for complex networks that are scale free with power law exponent between 2 and 3 identified cases of such networks that are *ultrasmall*. This term is associated with models in which the distance between two randomly chosen connected vertices grows *doubly logarithmically* in the number of vertices of the random graph. With N denoting the number of vertices, the function $\log \log N$ is a very slowly growing function. Presumably this is closer to empirical evidence which comes from networks that have millions of vertices but whose average distance between two randomly chosen vertices is very small.

An analytical relation between the two was first established by Cohen and Havlin [8] and by Dorogovtsev, Mendes and Samukhin [11]. It was shown rigorously for a variety of random graph models which exhibit power law degree distribution such as the Chung-Lu model [12], the Noros-Reittu model [27], the configuration model [20] as well as variations of the preferential attachment model [15] [9].

1.1   A geometric framework for complex networks

Recently, Krioukov et al. [25] introduced a geometric framework in order to represent the inherent inhomogeneity of a complex network. Their basic assumption is that the intrinsic hierarchies that are present in a complex network induce a tree-like structure. This suggests that the geometry of a complex network is hyperbolic.

The most common representations of the hyperbolic plane are the upper-half plane representation $\{z = x + iy \ : \ y > 0\}$ as well as the Poincaré unit disk which is simply the open disk of radius one, that is, $\{(u,v) \in \mathbb{R}^2 \ : \ 1 - u^2 - v^2 > 0\}$. Both spaces are equipped with the hyperbolic metric; in the former case this is $\frac{dx^2+dy^2}{y^2}$ whereas in the latter this is $4 \frac{du^2+dv^2}{(1-u^2-v^2)^2}$. It is well-known that the (Gaussian) curvature in both cases is equal to -1 and that the two spaces are isometric. In fact, there are more representations of the hyperbolic plane of curvature -1, which are isometrically equivalent to the above two. We will denote by $\mathbb{H}^2$ the class of these spaces.

We are now ready to give the definition of the basic model introduced

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in [25]. Consider the Poincaré disk representation of the hyperbolic plane. Let N be the number of vertices of the random graph, of which we assume that it tends to infinity. Consider also some fixed constant $\nu > 0$ and let $R > 0$ satisfy $N = \nu e^{R/2}$. It turns out that this parameter determines the average degree of the random graph. Let $V_N := \{v_1, \ldots, v_N\}$ be the set of vertices, where each v_i is a point selected randomly and independently from the disk of radius R centered at the origin of the Poincaré disk O; we denote this disk by $\mathcal{D}_R$.

Each of these points is distributed as follows. Assume that a random point u has *polar* coordinates (r, θ). The angle θ is uniformly distributed in $(0, 2\pi]$ and the probability density function of r, which we denote by $\rho_N(r)$, is determined by a parameter $\alpha > 0$ and is equal to

$$\rho(r) = \rho_N(r) = \begin{cases} \alpha \frac{\sinh \alpha r}{\cosh \alpha R - 1}, & \text{if } 0 \leq r \leq R \\ 0, & \text{otherwise} \end{cases}. \tag{1.1}$$

The above distribution is simply the uniform distribution on $\mathcal{D}_R$, but *on the hyperbolic plane of curvature $-\alpha^2$*. With elementary but tedious calculations, it can be shown that the length of a circle of radius r (centered at the origin) on the hyperbolic plane of curvature $-\alpha^2$ is $\frac{2\pi}{\alpha} \sinh(\alpha r)$, whereas the area of the circle of radius R (centered at the origin) is $\text{Area}_\alpha(\mathcal{D}_R) = \frac{2\pi}{\alpha^2}(\cosh(\alpha R) - 1)$. We will be using the function $\text{Area}_\alpha(\cdot)$ to denote the area of a measurable set in $\mathcal{D}_R$ on the hyperbolic plane of curvature $-\alpha^2$. Thus, if we set $\alpha = 1$, then the above becomes the density of the uniform distribution.

Alternatively, consider the disk $\mathcal{D}'_R$ of radius R around the origin O' of (the Poincaré disk representing) the hyperbolic plane of curvature $-\alpha^2$. We select N points within $\mathcal{D}'_R$ independently of each other, uniformly at random. These points are projected onto $\mathcal{D}_R$ preserving their polar coordinates. The projections of these points will be the vertex set V_N of the random graph.

The curvature of the hyperbolic plane determines the rate of growth of the space. A tedious calculation shows that the curvature of the hyperbolic plane is $-\alpha^2$, if we multiply the metric by $1/\alpha^2$. Hence, when $\alpha < 1$, the N points are distributed on a disk (namely $\mathcal{D}'_R$) which has smaller area compared to $\mathcal{D}_R$. This naturally increases the density of those points that are located closer to the origin. Similarly, when $\alpha > 1$ the area of the disk

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$\mathcal{D}'_R$ is larger than that of $\mathcal{D}_R$, and most of the N points are significantly more likely to be located near the boundary of $\mathcal{D}'_R$, due to the exponential growth of the volume.

Given the set V_N on $\mathcal{D}_R$ we define the random graph $\mathcal{G}(N; \alpha, \nu)$ on V_N, where two distinct vertices are joined if and only if they are within (hyperbolic) distance R from each other.

Notation

We now introduce some notation which we use throughout out proofs. Let a_N, b_N be two sequences of positive real numbers. We write $a_N \approx b_N$ to indicate that $a_N = \Theta(b_N)$, that is, there are real numbers $c, C > 0$ such that $cb_N \leq a_N \leq Cb_N$, for all natural numbers N. We also write $a_N \sim b_N$ to denote that $a_N / b_N \to \infty$, as $N \to \infty$.

If $\mathcal{E}_N$ is an event on the probability space $(\Omega_N, \mathbb{P}_N, \mathcal{F}_N)$, for each $N \in \mathbb{N}$, we say that $\mathcal{E}_N$ *occurs asymptotically almost surely (a.a.s.)* if $\mathbb{P}(\mathcal{E}_N) \to 1$ as $N \to \infty$. In our context, we mainly use the sequence of probability spaces that is induced by $\mathcal{G}(N; \alpha, \nu)$. However, later we introduce a variant of this model which is its *Poissonisation*. We will be using the term *a.a.s.* for that model as well.

1.1.1 Some facts about $\mathcal{G}(N; \alpha, \nu)$

We argue that the above model can be thought of as a geometrization of the random graph model that was introduced by F. Chung and L. Lu [12] [13] and is a special case of an inhomogeneous random graph. The notion of *inhomogeneous random graphs* was introduced Söderberg [28], but was defined more generally and studied in great detail by Bollobás, Janson and Riordan in [5]. In its most general setting, there is an underlying compact metric space $\mathcal{S}$ equipped with a measure μ on its Borel σ-algebra. This is the space of *types* of the vertices (defined below). A *kernel* κ is a bounded real-valued, non-negative function on $\mathcal{S} \times \mathcal{S}$, which is symmetric and measurable. The vertices of the random graph can be understood as points in $\mathcal{S}$. If $x, y \in \mathcal{S}$, then the corresponding vertices are joined with probability $\frac{\kappa(x,y)}{N} \wedge 1$, independently of every other pair (N is the total number of vertices). The points that are the vertices of the graph are approximately distributed according to μ. More specifically, the empirical distribution function on the N points converges weakly to μ as $N \to \infty$.

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Of particular interest is the case where the kernel function can be factorized and can be written $\kappa(x, y) = t(x)t(y)$; this is called a *kernel of rank 1*. Intuitively, the function $t(x)$ can be thought of as the weight or the type of vertex x. It is approximately its expected degree. In the special case where $t(x)$ follows a distribution that has a power law tail, the model becomes the so-called Chung-Lu model that was introduced in a series of papers [12] [13] (see also [19]).

We argue that in the random graph $\mathcal{G}(N; \alpha, \nu)$, the probability that two vertices are adjacent has this form. The proof of this fact relies on Lemma 2.5, which we will state later and is proved in [3]. It provides an approximate characterization of what it means for two points u, v to have hyperbolic distance at most R in terms of their *relative angle*, which we denote by $\theta_{u,v}$. For this lemma, we need the notion of the *type* of a vertex. For a vertex $v \in V_N$, if r_v is the distance of v from the origin, that is, the radius of v, then we set $t_v = R - r_v$ – we call this quantity the *type* of vertex v. As we shall shortly see, the type of a vertex is approximately exponentially distributed. If we substitute $R - t$ for r in (1.1), then assuming that t is fixed that expression becomes asymptotically equal to $\alpha e^{-\alpha t}$. Roughly speaking, Lemma 2.5 states that two vertices u and v of types t_u and t_v are within distance R (essentially) if and only if $\theta_{u,v} < 2\nu e^{t_u/2} e^{t_v/2}/N$. Hence, conditional on their types the probability that u and v are adjacent is proportional to $e^{t_u/2} e^{t_v/2}/N$. If we set $t(u) = e^{t_u/2}$, then $\mathbb{P}(t(u) \geq x) = \mathbb{P}(t_u \geq 2 \ln x) \approx e^{-2\alpha \ln x} = 1/x^{2\alpha}$. In other words, the distribution of $t(u)$ has a power-law tail with parameter 2α. Thus, the random graph $\mathcal{G}(N; \alpha, \nu)$ is a *dependent* version of the Chung-Lu model that emerges naturally from the hyperbolic geometry of the underlying space. The fact that this is a random geometric graph gives rise to local clustering, which is missing in the Chung-Lu model. There, most vertices have tree-like neighborhoods.

In fact, it can be shown that the degree of a vertex u in $\mathcal{G}(N; \alpha, \nu)$ that has type t_u is approximately distributed as a Poisson random variable with parameter proportional to $e^{t_u/2}$.

Gugelmann, Panagiotou and Peter [18] showed that the degree of a vertex has a power law with exponent $2\alpha + 1$. If $\alpha > 1/2$, then the exponent of the power law may take any value greater than 2. When $1 > \alpha > 1/2$, this exponent is between 2 and 3. They also showed that the average degree is a constant that depends on α and ν, and that the clustering coefficient (the probability of two vertices with a common neighbor to be joined by an

edge) of $\mathcal{G}(N; \alpha, \nu)$ is asymptotically bounded away from 0 with probability $1 - o(1)$ as $N \to \infty$.

Furthermore, the last two authors together with Müller [3] showed that $\mathcal{G}(N; \alpha, \nu)$ with high probability has a *giant component*, that is, a connected component containing a linear number of vertices if $1/2 < \alpha < 1$. When $\alpha > 1$, the size of the largest component is bounded by a function that is *sublinear* in N. When $\alpha = 1$, the existence of a giant component depends on the value of ν.

1.2   Results

In this contribution, we give an almost sure bound on the (graph) distance between two randomly chosen vertices that belong to the same connected component. We show that $\mathcal{G}(N; \alpha, \nu)$ is *ultrasmall* when $\frac{1}{2} < \alpha < 1$, that is, when the degree distribution has a power law tail with exponent between 2 and 3. More specifically, we show that a.a.s. the graph distance between two randomly chosen vertices that belong to the same component is of order $\log \log N$. However, the diameter of $\mathcal{G}(N; \alpha, \nu)$ grows at least logarithmically in N. This is a recent result of Kiwi and Mitsche [24], where they show that there is a connected component of diameter proportional to $\log N$. They also derive an upper bound on the diameter showing that the diameter is at most proportional to R^{1+C} a.a.s., for some positive constant C that depends on the parameters of the model. More recently, Friedrich and Krohmer [17] improved the constant showing that the exponent is at most $1/(2(1 - \alpha))$. They also show that if ν is small enough, then the exponent is equal to 1. Note that the Chung-Lu model exhibits logarithmic diameter [12].

For $\alpha > 1$, we show that a.a.s. $\mathcal{G}(N; \alpha, \nu)$ is *almost* ultrasmall: the graph distance between two randomly chosen vertices that belong to the same component is a.a.s. bounded by some polynomial of $\log \log N$. This range of α yields a power law degree distribution with exponent greater than 3. For this range, Chung and Lu [12] proved that the Chung-Lu model exhibits average distances of order $\log N$ asymptotically with high probability.

Let $d_G(u, v)$ denote the graph distance between two vertices u and v.

Theorem 1.1. *For $\zeta > 0$, assume that $1/2 < \alpha < 1$ and $u, v \in V_N$. Let τ be such that $\tau^{-1} = \log \left(\frac{1}{2\alpha - 1} \right)$. A.a.s. if $d_G(u, v) < \infty$, then $\left| \frac{d_G(u,v)}{\log R} - 2\tau \right| < \zeta$.*

In this regime, $\mathcal{G}(N; \alpha, \nu)$ does have a giant component and therefore for any two distinct vertices u, v we have $d_G(u, v) < \infty$ with probability that is asymptotically bounded away from 0. The upper bound (which is probably the most important) in the above result was also derived by Chung and Lu [12] for the Chung-Lu model with power law exponent between 2 and 3. That was under the assumption that the average degree is greater than 1. However, in our case a giant component is formed *independently* of what the average degree is, as long as $1/2 < \alpha < 1$. The full result for the Chung-Lu model can be found in [19].

The above result was also derived in the case of random graphs with given degree distribution that follows a power law with parameter between 2 and 3 by van der Hofstad et al. [20] in a stronger form which involves convergence in distribution.

Our second result provides an upper bound on the typical distance between two connected vertices when $\alpha > 1$. In this case there is no giant component a.a.s. However, the largest component contains *polynomially many* vertices as there is a number of a vertices of degree that scales polynomially in N. However, these components form also (almost) *ultrasmall* worlds.

Theorem 1.2. *Let* $\alpha > 1$, $\varepsilon > 0$. *There is a subset V' of vertices of $\mathcal{G}(N; \alpha, \nu)$ of size $(1 - o(1))N$ so that if $u, v \in V'$ and $d_G(u, v) < \infty$, then $d_G(u, v) \leq \log^{1+\varepsilon} \log N$.*

In the next section, we introduce the *Poissonisation* of $\mathcal{G}(N; \alpha, \nu)$ which is convenient for our calculations. Thereafter, we will state and prove some basic geometric facts regarding the hyperbolic plane, which allow us to express distances on the hyperbolic plane in terms of polar coordinates on $\mathcal{D}_R$. Subsequently, we proceed with the proof of Theorems 1.1 and 1.2.

The main idea behind the proof of Theorem 1.1 makes use of the existence of a very dense core that is formed by those vertices that have type at least $R/2$. We show that if two vertices are connected, then most likely they have short paths to the core which itself is a complete graph. These paths, which we call *exploding*, emerge also in the Chung-Lu model [12, 19].

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2 Preliminary results

2.1 Poissonisation

It will be significantly easier to work in a setting where, instead of having exactly N random points, our vertex set consists of $\mathrm{Po}(N)$ points on $\mathcal{D}_R$, in the hyperbolic plane of curvature $-\alpha^2$. Two vertices/points are declared adjacent exactly as in $\mathcal{G}(N; \alpha, \nu)$. We denote the resulting graph by $\mathcal{P}(N; \alpha, \nu)$. More specifically, the vertex set consists of the points of a Poisson point process in $\mathcal{D}_R$ (see [23]). In every measurable set $U \subseteq \mathcal{D}_R$, the number of points in U follows the Poisson distribution with parameter equal to $N \frac{\mathrm{Area}_\alpha(U)}{\mathrm{Area}_\alpha(\mathcal{D}_R)}$. Moreover, the numbers of points in any finite collection of pairwise disjoint measurable subsets of $\mathcal{D}_R$ are independent Poisson-distributed random variables.

We prove the following lemma that allows us to transfer results from the Poisson model into the $\mathcal{G}(N; \alpha, \nu)$ model. Let $\mathcal{A}_n$ denote a set of graphs on $V_n := \{1, \ldots, n\}$ that is closed under automorphisms. We call a family $\mathcal{A} = \{\mathcal{A}_n\}_{n \in \mathbb{N}}$ of graphs *(vertex-) non-decreasing*, if $G - v \in \mathcal{A}_{n-1}$ for any[1] $v \in V(G)$ implies $G \in \mathcal{A}_n$. Similarly, we call the family *(vertex-) non-increasing*, if $G - v \notin \mathcal{A}_{n-1}$ for any $v \in V(G)$ implies $G \notin \mathcal{A}_n$.

Lemma 2.1. *Assume that $\alpha > 0$ is fixed. Let $\mathcal{A}$ be a (vertex-) non-increasing family of graphs. For N large enough we have $\mathbb{P}(\mathcal{G}(N; \alpha, \nu) \notin \mathcal{A}) < 4\mathbb{P}(\mathcal{P}(N; \alpha, \nu) \notin \mathcal{A})$. The same holds if $\mathcal{A}$ is (vertex-) non-decreasing.*

Proof. Denote by E_{Po} and E the events that $\mathcal{P}(N; \alpha, \nu) \notin \mathcal{A}$ and $\mathcal{G}(N; \alpha, \nu) \notin \mathcal{A}$, respectively. We write

$$
\mathbb{P}(E_{Po}) = \sum_{N'=0}^{\infty} \mathbb{P}(E_{Po}|Po(N) = N') \cdot \mathbb{P}(Po(N) = N')
$$

$$
\geq \sum_{N'=N}^{\infty} \mathbb{P}(E_{Po}|Po(N) = N') \cdot \mathbb{P}(Po(N) = N')
$$

$$
\geq \sum_{N'=N}^{\infty} \mathbb{P}(E_{Po}|Po(N) = N) \cdot \mathbb{P}(Po(N) = N'),
$$

where we have used in the last line that, since $\mathcal{A}$ is non-increasing, we have $\mathbb{P}(E_{\mathrm{Po}}|\mathrm{Po}(N) = N') \geq \mathbb{P}(E_{\mathrm{Po}}|\mathrm{Po}(N) = N)$ for $N' \geq N$. Let us also note

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[1] $G - v \in \mathcal{A}_{n-1}$ means that $G - v$ is isomorphic to a member of $\mathcal{A}_{n-1}$

that $\mathbb{P}(E_{\mathrm{Po}}|\mathrm{Po}(N) = N) = \mathbb{P}(E)$. Thus,

$$
\begin{aligned}
\mathbb{P}(E_{\mathrm{Po}}) &\geq \sum_{N'=N}^{\infty} \mathbb{P}(E) \cdot \mathbb{P}(\mathrm{Po}(N) = N') \\
&= \mathbb{P}(E) \cdot \mathbb{P}(\mathrm{Po}(N) \geq N) \\
&> \frac{1}{4} \cdot \mathbb{P}(E),
\end{aligned}
$$

where the last line holds N large enough (by an application of, say, the central limit theorem). The second part of the lemma follows similarly, bounding the sum by taking only the terms where $N' \leq N$. $\qquad\square$

This implies that if $\mathbb{P}(\mathcal{P}(N; \alpha, \nu) \notin \mathcal{A}) = o(1)$, then $\mathbb{P}(\mathcal{G}(N; \alpha, \nu) \notin \mathcal{A}) = o(1)$.

During some of our proofs, we will need to bound probabilities of events that are associated with a certain subset of vertices X, whose positions in $\mathcal{D}_R$ have been realised. For a certain measurable subset $U \subset \mathcal{D}_R$ which does not contain any vertex in X so that $\mathcal{D}_R \setminus U$ has positive Lebesgue measure, the vertices of the random graph $\mathcal{P}_{X,U}(N; \alpha, \nu)$ consist of X together with set of points of a Poisson process on $\mathcal{D}_R \setminus U$ with curvature $-\alpha^2$ with parameter $N - |X|$. Hence, this process "produces" $N - |X|$ vertices on average, thus giving N vertices in total on average. If we condition on the number of vertices of this Poisson process being N', then the resulting random graph is distributed as $\mathcal{G}(N'; \alpha, \nu)$ conditional on U *being empty* and X being located at its particular positions.

Let $\mathcal{A}_X$ be a graph property associated with the set X. We call this *non-decreasing* if

$$
\mathbb{P}_{\mathcal{P}_{X,U}(N;\alpha,\nu)}(\mathcal{A}_X \mid \mathrm{Po}(N-|X|) = N_1) \leq \mathbb{P}_{\mathcal{P}_{X,U}(N;\alpha,\nu)}(\mathcal{A}_X \mid \mathrm{Po}(N-|X|) = N_2),
$$

whenever $N_1 \leq N_2$. If the opposite inequality holds, we call the property *non-increasing*. Note that $\mathbb{P}_{\mathcal{P}_{X,U}(N;\alpha,\nu)}(\mathcal{A}_X \mid \mathrm{Po}(N-|X|) = N')$ is the probability of $\mathcal{A}_X$ in the space $\mathcal{G}(N' + |X|; \alpha, \nu)$ conditional on X being at certain positions in $\mathcal{D}_R$ and U being empty – we denote this by $\mathcal{G}_{X,U}(N; \alpha, \nu)$. Hence, arguing as in the proof of the previous lemma we have

Lemma 2.2. *If $\mathcal{A}_X$ is either a non-decreasing or a non-increasing property*

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that is associated with a certain set of vertices X, then

$$\mathbb{P}_{\mathcal{P}_{X,U}(N;\alpha,\nu)}(\mathcal{A}_X) \geq \frac{1}{4}\mathbb{P}_{\mathcal{G}_{X,U}(N;\alpha,\nu)}(\mathcal{A}_X),$$

for any measurable $U \subset \mathcal{D}_R$ such that $X \cap U = \emptyset$ and $\mathcal{D}_R \setminus U$ has positive Lebesgue measure.

The following useful fact follows directly from the definition of the process, using the measure defined for the distribution of the points.

Fact 2.3. *Let A be a subset of $\mathcal{D}_R \setminus U$, for some measurable subset $U \subset \mathcal{D}_R$, and X be a set of vertices located in $\mathcal{D}_R$, such that $X \cap U, A \cap U = \emptyset$. Let N_A be the expected number of vertices in A, in $\mathcal{G}_{X,U}(N;\alpha,\nu)$, and denote by $\mathcal{E}_A$ the event that A is empty. We have*

$$\mathbb{P}_{\mathcal{P}_{X,U}(N;\alpha,\nu)}(\mathcal{E}_A) = \exp(-N_A).$$

2.2 Geometric properties of $\mathcal{D}_R$

We state a simple geometric fact, which we will use in the following sections. With O being the origin, we say that a vertex v lies *above* some edge uw when v is inside the (hyperbolic) triangle Ouw, where uw is the geodesic path in $\mathcal{D}_R$ that joins u with w. Similarly, v lies *below* the edge uw, if v does not lie above uw but some radial projection of v towards O lies above uw.

Fact 2.4. *If the vertex w lies above the edge $u'u''$, then w is adjacent to u' and to w'. Moreover, the geodesic segments connecting w to u' and u'' lie entirely in the triangle $Ou'u''$.*

Proof. The hyperbolic triangle $Ou'u''$ has only sides of length at most R. The vertex w lies inside this triangle, so it has distance at most R from O, u' and u''. This is the case for any point v in the triangle $Ou'u''$. The geodesic from v to u' is entirely in the triangle, since otherwise it would have to cross one of the sides. A crossing point would therefore have two paths of minimum length to u', which is a contradiction. The same argument also works for the geodesic segment between v and u''. □

The following lemma provides a useful (almost) characterization of the fact that two vertices are within hyperbolic distance R, given their types.

Recall first that for any two points/vertices u, v on $\mathcal{D}_R$, their *relative angle* $\theta_{u,v}$ is defined as $\min\{u\hat{O}v, 2\pi - u\hat{O}v\}$. Note that $\theta_{u,v} \leq \pi$ always. The lemma reduces a statement about hyperbolic distances to a statement about the relative angle between two points. Its proof can be found in [16] and [3]. For two points p, v, let

$$\hat{\theta}_{p,v} := \min\left\{2(1+\varepsilon)e^{\frac{t_p+t_v-R}{2}}, \pi\right\} = \min\left\{2(1+\varepsilon)\frac{\nu}{N}e^{\frac{t_p+t_v}{2}}, \pi\right\}, \text{ and}$$

$$\check{\theta}_{p,v} := \min\left\{2(1-\varepsilon)e^{\frac{t_p+t_v-R}{2}}, \pi\right\} = \min\left\{2(1-\varepsilon)\frac{\nu}{N}e^{\frac{t_p+t_v}{2}}, \pi\right\}.$$

For $c_0 = c_0(\varepsilon)$, that depends on ε as in the following lemma, we call the set

$$T_\varepsilon^+(v) := \left\{p \in D_R \ : \ t_p + t_v - R < -c_0, \ \theta_{p,v} \leq \hat{\theta}_{p,v}\right\}$$

the *outer tube* of v. Similarly, we call the set

$$T_\varepsilon^-(v) := \left\{p \in D_R \ : \ t_p + t_v - R < -c_0, \ \theta_{p,v} \leq \check{\theta}_{p,v}\right\}$$

the *inner tube* of v.

Lemma 2.5. *For any $\varepsilon > 0$ there exists an $N_0 > 0$ and a $c_0 > 0$ such that for any $N > N_0$ and $u, v \in \mathcal{D}_R$ with $t_u + t_v < R - c_0$ the following hold.*

- *If $u \in T_\varepsilon^-(v)$, then $d(u,v) < R$.*

- *If $u \notin T_\varepsilon^+(v)$, then $d(u,v) > R$.*

2.3 Properties of $\mathcal{G}(N; \alpha, \nu)$

We state some general results about the graphs, the proofs of which can be found in [3].

Lemma 2.6. *Let $\bar{\rho}(t)$ be the distribution of the types. For any $\varepsilon \in (0,1)$, uniformly for $0 \leq t < (1-\varepsilon)R$ we have*

$$\bar{\rho}(t) = \rho(R - t) = (1 + o(1))\alpha e^{-\alpha t}. \tag{2.1}$$

The following fact is an immediate consequence of the above.

Corollary 2.7. *Let $\omega : \mathbb{N} \to \mathbb{N}$ be an increasing function such that $\omega(N) \to \infty$ as $N \to \infty$. The expected number of vertices of type at least $R/(2\alpha) +$*

$\omega(N)$ in $\mathcal{G}(N; \alpha, \nu)$ is $o(1)$. Hence, with probability $1 - o(1)$ all vertices in V_N have type at most $\frac{1}{2\alpha}R + \omega(N)$.

3 Proof of Theorem 1.1: upper bound

In this section we assume $1/2 < \alpha < 1$.

Definition 3.1. *For $G \in \mathcal{P}(N; \alpha, \nu)$ or $G \in \mathcal{G}(N; \alpha, \nu)$, let $\operatorname{Core}(G) = \{v \in V(G) \ : \ t_v \geq R/2\}$ be the core of G.*

Note that for every pair of vertices $u, v \in \operatorname{Core}(G)$, by the triangle inequality the distance between u and v is at most R, so $uv \in E(G)$. In other words, the subgraph that is induced by the vertices in $\operatorname{Core}(G)$ is complete.

Lemma 3.2. *Let $\omega(N)$ be such that $\omega(N) \to \infty$ as $N \to \infty$ but $\omega(N) = o(R)$. Let x be a vertex such that $t_x < \log \log R$ and $U \subset \mathcal{D}_R$ an open subset of $\mathcal{D}_R$ which does not contain any points of type at least $\log \log R$ and has $\operatorname{Area}_\alpha(U) = o(\operatorname{Area}_\alpha(\mathcal{D}_R))$. Let $G \in \mathcal{P}_{x,U}(N; \alpha, \nu)$. A.a.s. there is a vertex $u \in \operatorname{Core}(G)$ such that $uv \in E(G)$ for every vertex v with $t_v \geq \frac{2\alpha - 1}{2\alpha}R + \omega(N)$.*

Proof. By the triangle inequality, any such vertex v is adjacent to any vertex of radius at most $R(2\alpha - 1)/(2\alpha) + \omega(N)$, so it is sufficient to show that a.a.s. the disc D_r of radius $r := \frac{2\alpha - 1}{2\alpha}R + \omega(N)$ is non-empty. Note that $r < R/2$, for any N large enough, as $\omega(N) = o(R)$ and $\alpha < 1$, so any vertex in D_r belongs to the core. Let N_r be the number of vertices in D_r.

Note first that $\frac{2\alpha - 1}{2\alpha} - 1 = -\frac{1}{2\alpha}$. Thus $r - R = -\frac{R}{2\alpha} + \omega(N)$, whereby $\alpha(r - R) = -R/2 + \alpha\omega(N)$. As $D_r \cap U = \emptyset$, these identities imply that

$$\mathbb{E}\,[\,N_r\,] = (N - 2)\frac{\cosh(\alpha r) - 1}{\operatorname{Area}_\alpha(\mathcal{D}_R) - \operatorname{Area}_\alpha(U)} = (N - 2)\frac{\cosh(\alpha r) - 1}{\cosh(\alpha R)(1 - o(1))}$$

$$\sim Ne^{\alpha(r-R)} = Ne^{-R/2 + \omega(N)} \stackrel{N = \nu e^{R/2}}{=\!=\!=} \nu e^{\alpha\omega(N)}.$$

Using this and Fact 2.3 we get

$$\mathbb{P}(N_r \neq 0) = 1 - e^{-(1 + o(1))\nu e^{\alpha\omega(N)}} = 1 - o(1).$$

$\square$

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In fact, the only component we consider is the one containing the vertices in the core. We show that most pairs of vertices that are connected have a short path into the core. These paths naturally give short paths connecting all the vertices in the component. We are interested in the following paths in which the type of the vertices increases exponentially along the path.

Definition 3.3. *For $\delta > 0$, we call a path $P = v_1, v_2, \ldots, v_m$ in G a δ-exploding path if $v_m \in \mathrm{Core}(G)$ and $t_{v_{i+1}} \geq (1+\delta)t_{v_i}$ for $1 \leq i \leq m-2$.*

Not every vertex in the giant component has an exploding path into the core. However, the vertices that do not have such a path are more likely to have a very low type. In particular, we prove that any vertex of type at least $\log \log R$ has an exploding path into the core with probability $1 - o(1)$. We actually show this lemma for the Poisson model. The result does transfer to $\mathcal{G}(N; \alpha, \nu)$, due to its monotonicity, but we are going to use it later in this form.

Lemma 3.4. *Let $\delta = 2\frac{1-\alpha}{2\alpha-1}$ and $\zeta < \delta$ be a positive real number. Assume that v and x are vertices such that $t_v \geq \log \log R \geq t_x$ and $U \subset \mathcal{D}_R$ an open subset which does not contain any points of type at least t_v so that U is contained in a sector of $\mathcal{D}_R$ that spans a $o(1)$ angle. With probability (in the space $\mathcal{P}_{\{v,x\},U}(N; \alpha, \nu)$) $1 - e^{-\Theta(\log^{(\alpha-\frac{1}{2})\zeta} R)}$, there is a $(\delta - \zeta)$-exploding path starting at v.*

Proof. Take any $\varepsilon < \frac{1}{4}$ and assume that $N > N_0$, where N_0 is as in Lemma 2.5.

By Lemma 3.2, if v satisfies $t_v \geq \frac{2\alpha-1}{2\alpha}R + \omega(N)$, then a.a.s. there is a vertex $u \in G$ with $t_u \geq R/2$ and $vu \in E(G)$. In other words, if $t_v \geq \frac{2\alpha-1}{2\alpha}R + \omega(N)$, then we are done.

Assume now that $t_v < \frac{2\alpha-1}{2\alpha}R + \omega(N)$. As $1 + \delta = \frac{1}{2\alpha-1}$, it follows that $(1+\delta)t_v < \frac{1}{2\alpha}R + \frac{\omega(N)}{2\alpha-1}$. Note that by Corollary 2.7, it suffices to consider only points of type no larger than $\frac{1}{2\alpha}R + \frac{\omega(N)}{2\alpha-1}$.

Let $v_1 = v$. We will construct inductively a series of (random) sets $T_i \subset \mathcal{D}_R$, for $i \geq 2$, in each of which we find a vertex v_i, which will be the ith vertex in the exploding path.

For two points p, p', let $\vartheta_{p,p'} = \theta_{p,p'}$ if p' is in the anti-clockwise direction from p, but $\vartheta_{p,p'} = -\theta_{p,p'}$, otherwise.

Assume that we have exposed v_i. For any point $p \in \mathcal{D}_R$ we let

$$\hat{T}_\varepsilon^-(p) := \left\{ p' \in \mathcal{D}_R \ : \ |t_{p'} - (1+\delta)t_p| < \zeta t_p, \ \frac{\varepsilon\nu}{N} e^{\frac{t_p + t_{p'}}{2}} \leq \vartheta_{p',p} \leq \frac{2(1-\varepsilon)\nu}{N} e^{\frac{t_{p'} + t_p}{2}} \right\}.$$

We take $T_i := \hat{T}_\varepsilon^-(v_i)$. Let A be the set of vertices that are located in $\hat{T}_\varepsilon^-(v_i)$. Note that, as the angle covered by U is $o(1)$, we have that $\mathrm{Area}_\alpha(U) = o(\mathrm{Area}_\alpha(\mathcal{D}_R))$. Hence, the area of a set in $\mathcal{D}_R \setminus U$ is within a $1 - o(1)$ factor from the area in $\mathcal{D}_R$ (both on the hyperbolic plane of curvature $-\alpha^2$).

So, for any $\varepsilon \in (0, 1/4)$ and for N large enough we have

$$\mathbb{E}\left[|A|\right] \geq 2(1 - \frac{3}{2}\varepsilon)\frac{N-2}{2\pi} \int_{(1+\delta-\zeta)t_{v_i}}^{(1+\delta+\zeta)t_{v_i}} e^{\frac{1}{2}(t_{v_i}+t-R)}(1 - o(1))e^{-\alpha t}dt$$

$$\geq 2(1 - \frac{3}{2}\varepsilon)(1 - o(1))\frac{N}{2\pi}\frac{\nu}{N}e^{\frac{t_{v_i}}{2}} \int_{(1+\delta-\zeta)t_{v_i}}^{(1+\delta+\zeta)t_{v_i}} e^{(\frac{1}{2}-\alpha)t}dt$$

$$\overset{\varepsilon<1/4}{\geq} \frac{\nu}{2\pi}e^{\frac{1}{2}t_{v_i}}\frac{1}{2\alpha-1}\left(e^{(\frac{1}{2}-\alpha)(1+\delta-\zeta)t_{v_i}} - e^{(\frac{1}{2}-\alpha)(1+\delta+\zeta)t_{v_i}}\right).$$

But $(1 + \delta + \zeta)t_{v_i} - (1 + \delta)t_{v_i} + \zeta t_{v_i} > 2\zeta t_{v_i} \to \infty$, whereby the above becomes:

$$\mathbb{E}\left[|A|\right] \geq \frac{\nu}{2\pi}\frac{1}{2\alpha-1}e^{\frac{1}{2}t_{v_i}-(\alpha-\frac{1}{2})(1+\delta-\zeta)t_{v_i}}(1 - o(1)).$$

Furthermore, $(\alpha - \frac{1}{2})(1 + \delta) = \frac{2\alpha-1}{2}\frac{1}{2\alpha-1} = \frac{1}{2}$ and finally, we have

$$\mathbb{E}\left[|A|\right] \geq \frac{\nu}{2\pi}\frac{1}{2\alpha-1}e^{(\alpha-\frac{1}{2})\zeta t_{v_i}}(1 - o(1)) \overset{2\alpha-1<1}{\geq} \frac{\nu}{2\pi}e^{(\alpha-\frac{1}{2})\zeta t_{v_i}},$$

for N large enough. Hence, by Fact 2.3 we have

$$\mathbb{P}(|A| > 0) = 1 - \mathbb{P}(|A| = 0)$$
$$\geq 1 - \exp\left(-\frac{\nu}{2\pi}e^{(\alpha-\frac{1}{2})\zeta t_{v_i}}\right).$$

As $t_{v_i} \geq \log\log R$, we have $\mathbb{P}(|A| = 0) \leq \exp\left(-\frac{\nu}{\pi}(\log R)^{(\alpha-\frac{1}{2})\zeta}\right)$. If $|A| > 0$, then there are vertices that are located inside T_i and we let v_{i+1} be one of them – the choice is arbitrary. The following claim guarantees that $T_{i+1} = \hat{T}_\varepsilon^-(v_{i+1})$ is disjoint from T_i and when we repeat the argument there is no danger to expose again area which we have already exposed.

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Claim 3.5. *For all N large enough and for all $i \geq 1$ the following holds. For all $p \in \hat{T}_\varepsilon^-(v_i)$ we have $T_\varepsilon^+(v_i) \cap \hat{T}_\varepsilon^-(p) = \emptyset$.*

Proof of Claim 3.5. Consider a point $p \in \hat{T}_\varepsilon^-(v_i)$ and let $p' \in \hat{T}_\varepsilon^-(p)$. We will show that

$$\vartheta_{p',v_i} \gg 2(1+\varepsilon)\frac{\nu}{N}e^{\frac{t_v+t_{p'}}{2}}.$$

We write $\vartheta_{p',v_i} = \vartheta_{p',p} + \vartheta_{p,v_i}$. Since $p' \in \hat{T}_\varepsilon^-(p)$ and $p \in \hat{T}_\varepsilon^-(v_i)$ we have

$$\vartheta_{p',p} \geq \varepsilon\frac{\nu}{N}e^{\frac{t_{p'}+t_p}{2}} \quad \text{and} \quad \vartheta_{p,v_i} \geq \varepsilon\frac{\nu}{N}e^{\frac{t_p+t_{v_i}}{2}}.$$

Hence

$$
\begin{aligned}
\vartheta_{p',p} + \vartheta_{p,v_i} &\geq \varepsilon\frac{\nu}{N}\left(e^{\frac{t_{p'}+t_p}{2}} + e^{\frac{t_p+t_{v_i}}{2}}\right) \\
&= \varepsilon\frac{\nu}{N}e^{\frac{t_{p'}+t_{v_i}}{2}}\left(e^{\frac{t_p-t_{v_i}}{2}} + e^{\frac{t_p-t_{p'}}{2}}\right) > \varepsilon\frac{\nu}{N}e^{\frac{t_{p'}+t_{v_i}}{2}}e^{\frac{t_p-t_{v_i}}{2}} \\
&\geq \varepsilon\frac{\nu}{N}e^{\frac{t_{p'}+t_{v_i}}{2}}e^{(\delta-\zeta)t_{v_i}} \stackrel{(\delta-\zeta)t_{v_i}\to\infty}{\gg} 2(1+\varepsilon)\frac{\nu}{N}e^{\frac{t_{p'}+t_{v_i}}{2}}.
\end{aligned}
$$

In fact, $(\delta - \zeta)t_{v_i} \geq (\delta - \zeta)\log\log R$, and therefore the inequality holds uniformly for all N that are large enough. $\qquad\square$

If we start at type at least $\log\log R$, it takes $O(\log R)$ steps to reach type $\frac{2\alpha-1}{2\alpha}R + \omega(N)$; at that point we can complete the exploding path using the vertex whose existence is guaranteed by Lemma 3.2. Thus for any given vertex v with $t_v > \log\log R$ we have

$$
\begin{aligned}
\mathbb{P}(\exists \text{ sequence of vertices } v_2, \ldots) &= \left(1 - \exp\left(-\frac{\nu}{\pi}(\log R)^{(\alpha-\frac{1}{2})\zeta}\right)\right)^{O(\log R)} \\
&= 1 - O(\log R)\exp\left(-\frac{\nu}{\pi}(\log R)^{(\alpha-\frac{1}{2})\zeta}\right) \\
&= 1 - \exp\left(-\Theta\left(\log^{(\alpha-\frac{1}{2})\zeta} R\right)\right),
\end{aligned}
$$

as $xe^{-ax^b} = o(1)$ for $0 < a, b$ and $x \to \infty$.

$\qquad\square$

Remark 3.6. *In fact, if the type of v is $O(1)$, that is, v is a typical vertex, then the probability that there is a $(\delta - \zeta)$-exploding path starting at v is bounded away from 0. With slightly more work, one can show that*

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two vertices u and v have both an exploding path with probability that is asymptotically bounded away from 0. Thus, $d_G(u,v) < \infty$ with probability that is asymptotically bounded away from 0. Alternatively, this follows from the main theorem in [3], according to which $\mathcal{G}(N;\alpha,\nu)$ has giant component a.a.s. if $1/2 < \alpha < 1$.

We are now ready to proceed with the upper bound in Theorem 1.1

Proof of Theorem 1.1: upper bound. Let u,v be two vertices. We will show that the event $d_G(u,v) < \infty$ but $d_G(u,v) \geq (2\tau + \zeta^{1/2})\log R$ occurs with probability $o(1)$. Note that this is in the $\mathcal{G}(N;\alpha,\nu)$ space. Also, for convenience, we have taken the ζ that appears in the statement of Theorem 1.1 as $\zeta^{1/2}$. We denote this event by $\mathcal{E}_N(\tau,\zeta)$. Also, let $\mathcal{A}_N$ denote the event that the relative angle between u and v is greater than $\nu \frac{2\zeta_\varepsilon \log R}{N}$, where $\zeta_\varepsilon := \zeta(1-\varepsilon)$, for some $\varepsilon \in (0,1)$. Of course, the probability of $\overline{\mathcal{A}_N}$ is $o(1)$ and therefore it suffices to prove that $\mathbb{P}\left[\mathcal{E}_N(\tau,\zeta) \cap \mathcal{A}_N\right] = o(1)$.

If $\mathcal{E}_N(\tau,\zeta)$ is realised, then there must be a *minimal* path between vertices u and v. In this context, a minimal path is meant to be an induced path. Let $P_{\min}$ denote such a path. Assume, in addition, that $\mathcal{A}_N$ is simultaneously realised, that is, $\theta_{u,v} > \nu \frac{2\zeta_\varepsilon \log R}{N}$. With this assumption, let $P_{\min}(u)$ denote the sub-path of $P_{\min}$ starting at u and ending at the first vertex whose relative angle with u exceeds $\nu \frac{\zeta_\varepsilon \log R}{N}$. Similarly, let $P_{\min}(v)$ denote the sub-path of $P_{\min}$ starting at v and ending at the first vertex whose relative angle with v exceeds $\nu \frac{\zeta_\varepsilon \log R}{N}$. Clearly, as $\mathcal{A}_N$ is realized, the two paths may overlap, but they have at most one edge in common.

Assume without loss of generality that v is at angle $\theta_{u,v} \leq \pi$ in the anti-clockwise direction from u. Consider the sectors consisting of points of relative angle at most $\nu \frac{\zeta_\varepsilon \log R}{N}$ from a point x:

$$S_h^+(x) := \left\{ p \in \mathcal{D}_R \; : \; t_p > \log \log R, \; 0 < \vartheta_{x,p} < \nu \frac{\zeta_\varepsilon \log R}{N} \right\}$$

and

$$S_h^-(x) := \left\{ p \in \mathcal{D}_R \; : \; t_p > \log \log R, \; -\nu \frac{\zeta_\varepsilon \log R}{N} < \vartheta_{x,p} < 0 \right\}.$$

There are two cases:

1. either each one of $S_h^+(u), S_h^-(u), S_h^+(v), S_h^-(v)$ contains a vertex that

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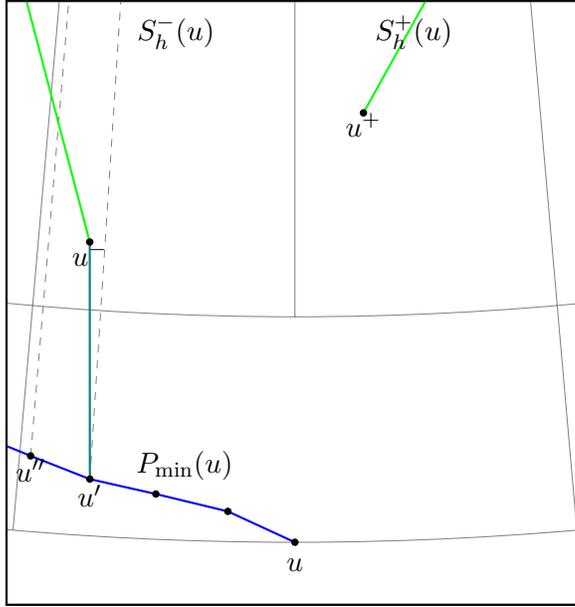
Figure 1: Creating a short path into the core under $\mathcal{S}$.

    is the starting vertex of a $(\delta - \zeta)$-exploding path,

2. or at least one of them is either empty or all of its vertices are not the endpoints of a $(\delta - \zeta)$-exploding path.

Let $\mathcal{S}$ denote the former and let $\overline{\mathcal{S}}$ denote the latter. We will show that $\mathbb{P}(\overline{\mathcal{S}}) = o(1)$. First consider, without loss of generality, the set $S_h^+(u)$. The probability that this set is empty is $o(1)$. Indeed, let $N_{S_h^+(u)}$ be the number of vertices that appear into this sector. Then

$$\mathbb{E}\left[N_{S_h^+(u)}\right] = N\frac{\cosh(\alpha(R - \log\log R)) - 1}{\cosh(\alpha R) - 1}\frac{1}{2\pi}\nu\frac{\zeta_\varepsilon \log R}{N} \approx \log^{1-\alpha} R \to \infty.$$

The distribution of $N_{S_h^+(u)}$ is binomial and the application of a standard Chernoff bound implies that $\mathbb{P}\left[N_{S_h^+(u)} = 0\right] = o(1)$.

    If $S_h^+(u)$ is not empty and all of its vertices are not the beginnings of a $(\delta - \zeta)$- exploding path, then the vertex with lowest type in $S_h^+(u)$ does not have a $(\delta - \zeta)$-exploding path starting at it as well. We call this vertex the *first* vertex in $S_h^+(u)$.

Claim 3.7. *The probability that the first vertex in $S_h^+(u)$ does not have a $(\delta - \zeta)$-exploding path starting at it is $o(1)$.*

Proof of Claim 3.7. Conditional on having at least one vertex in $S_h^+(u)$, let u' be the first vertex (with probability 1 there will be exactly one such vertex) which we expose and assume that the area in $S_h^+(u)$ that consists of points with type greater than $t_{u'}$ has not been exposed. Let us switch temporarily to $\mathcal{P}_{X,U}(N; \alpha, \nu)$, where $X = u, u'$ and U the subset of $S_h^+(u)$ below u'. Then by Lemma 3.4, there is a $(\delta - \zeta)$-exploding path starting at u' with probability $1 - o(1)$ uniformly over $t_{u'} \geq \log\log R$. This lemma can be applied as the area above u' has not been exposed in the corresponding Poisson process and the proof of Lemma 3.4 deals only only with that area. The result transfers to $\mathcal{G}(N; \alpha, \nu)$ (conditional on U being empty and on the realisations of u and u'), through Lemma 2.2, due to the fact that this property is non-decreasing. $\qquad\square$

Then, since the probability that $S_h^+(u)$ is empty is $o(1)$, the union bound implies that $\mathbb{P}\left[\overline{\mathcal{S}}\right] = o(1)$.

We will show that $\mathbb{P}\left[\mathcal{E}_N(\tau, \zeta) \cap \mathcal{A}_N \cap \mathcal{S}\right] = 0$. Observe that any vertex which belongs to $S_h^+(u) \cup S_h^-(u)$ (or to $S_h^+(v) \cup S_h^-(v)$, respectively) will be adjacent to a vertex in $P_{\min}(u)$ ($P_{\min}(v)$, resp.). Indeed, if $P_{\min}(u)$ contains a vertex in $S_h^+(u) \cup S_h^-(u)$, then this must be adjacent to any other vertex in $S_h^+(u) \cup S_h^-(u)$. This is the case as $S_h^+(u) \cup S_h^-(u) \subseteq T_\varepsilon^-(u')$ for any $u' \in S_h^+(u) \cup S_h^-(u)$, provided that $\zeta < 1$. To see this, note that any two points in $S_h^+(u) \cup S_h^-(u)$ have relative angle at most $2\zeta_\varepsilon \frac{\nu}{N}$. However, for any point in $S_h^+(u) \cup S_h^-(u)$, its inner tube consists of all points of relative angle at most $2(1 - \varepsilon)\frac{\nu e^{\log\log R}}{N}$ from it. Thus, if $\zeta_\varepsilon < 1 - \varepsilon$ (that is, $\zeta < 1$), then the containment follows. In this case, some vertex of $P_{\min}(u)$ will be connected to the first vertex in $S_h^+(u) \cup S_h^-(u)$.

Suppose now that all vertices of $P_{\min}(u)$ do not belong to $S_h^+(u) \cup S_h^-(u)$. Let u^+, u^- be vertices in $S_h^+(u)$ and $S_h^-(u)$ respectively, which are the starting vertices of $(\delta - \zeta)$-exploding paths P_{u^+} and P_{u^-}. There are two *consecutive* vertices in $P_{\min}(u)$ say u', u'' such that either $\vartheta_{u'',u^+} > 0 > \vartheta_{u',u^+}$ or $\vartheta_{u'',u^-} > 0 > \vartheta_{u',u^-}$. Thus, either u^+ or u^- is "above" the edge $u'u''$ and therefore by Fact 2.4 either u^+ or u^- is adjacent to both $u'u''$. The length of any exploding path is at most $\log R / \log(1 + \delta - \zeta)$. Thus, $|P_{u^+}|, |P_{u^-}| \leq \log R / \log(1 + \delta - \zeta)$. The following bounds the length of $P_{\min}(u), P_{\min}(v)$:

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Claim 3.8. *Both $P_{\min}(u)$ and $P_{\min}(v)$ have length at most $\zeta \log R$.*

Proof of Claim 3.8. Consider $P_{\min}(u)$ (the proof for $P_{\min}(v)$ is identical). Since $P_{\min}(u)$ is part of a minimal path, it follows that if we take the set of vertices of $P_{\min}$ that are at even distance from u, then there cannot be an edge between any two of them, for this would contradict the minimality of $P_{\min}$. Let $P^e_{\min}(u)$ be this set of vertices. For any vertex $u' \in P^e_{\min}(u)$ consider the sector $T(u') := \{p \in \mathcal{D}_R \ : \ \theta_{u',p} < (1-\varepsilon)\frac{\nu}{N}\}$. There cannot be distinct $u', u'' \in P^e_{\min}(u)$ such that $T(u') \cap T(u'') \neq \emptyset$. If this were the case, then their relative angle would be at most $2(1-\varepsilon)\frac{\nu}{N}$ and by Lemma 2.5 they would be adjacent. But there are at most $\nu\frac{\zeta_\varepsilon \log R}{N} / \left(2(1-\varepsilon)\frac{\nu}{N}\right) = \frac{\zeta}{2}\log R$ such sectors inside the sector of angle $\nu\frac{\zeta_\varepsilon \log R}{N}$ in the anti-clockwise direction from u. Thus $|P^e_{\min}(u)| \leq \frac{\zeta}{2}\log R$, whereby the length of $P_{\min}$ is at most $\zeta \log R$. $\qquad \square$

Thus

$$d_G(u,v) \leq |P_{\min}(u)| + |P_{u^+}| + 1 + |P_{u^-}| + |P_{\min}(v)|$$

$$\leq 2\left(\frac{1}{\log(1+\delta-\zeta)} + \zeta + o(1)\right)\log R$$

Hence, there exists a ζ such that for all N large enough $\frac{1}{\log(1+\delta-\zeta)} + \zeta + o(1) < \tau + \zeta^{1/2}$. This implies that $\mathcal{E}_N(\tau, \zeta)$ is not realised. $\qquad \square$

Remark 3.9. *If we replace the angles that determine the domains S^+_h and S^-_h by a quantity that is proportional to $R^{\frac{1}{1-\alpha}}/N$ and the lower bound on the type by $\frac{1}{2(1-\alpha)}\log R$, then the probabilities that appear above become $o(N^{-2})$. Thus, the analogous of the above bound on $d_G(u,v)$ holds for all pairs of vertices, and implies that the diameter is proportional to $R^{\frac{1}{1-\alpha}}$ a.a.s. This upper bound is worse than the one obtained in [17].*

4 Proof of Theorem 1.1: lower bound

For given vertices $u, v \in V_N$, let $\mathcal{L}_{\zeta,N}(u,v)$ be the event that $d_G(u,v) < (2\tau - \zeta)\log R =: L$, for some $\zeta > 0$. Assume that $t_u, t_v < \log\log R$ - by Lemma 2.6 this event occurs with probability $1 - o(1)$. Let $\mathcal{T}_{u,v}$ denote this event. By Lemma 2.5, for any $T \leq R/2 - 2\log\log R$, if u and v are connected

through a path of length at most ℓ_u where the intermediate vertices have type at most T, then

$$\theta_{u,v} \leq 4\nu \frac{e^T}{N} L \leq 4\nu \frac{e^{R/2}}{N} \frac{L}{\log^2 R} = 4 \frac{L}{\log^2 R}.$$

Conditional on $\mathcal{T}_{u,v}$, the probability of this event is $O(L/\log^2 R) = o(1)$. Now, if there is a path of length at most L that joins u to v that contains an intermediate vertex of type at least $R/2 - 2\log\log R$, then there must be a path of length at most $L/2$ either from u or from v to this vertex. Denote by $d_G(u, core)$ the graph distance of the vertex u to a vertex of type at least $R/2 - 2\log\log R$. The following lemma proves that almost all vertices are, in some sense, far away from vertices this type, immediately proving the lower bound.

Lemma 4.1. *Assume that $t_u \leq \log\log R$. For $\zeta > 0$, we have*

$$\mathbb{P}(d_G(u, core) \leq (\tau - \zeta^{1/2})\log R) = o(1).$$

We appeal to Lemma 2.2 on the event $\{d(u, core) \leq (\tau - \zeta^{1/2})\log R\}$. Clearly, this is a non-decreasing event in the sense that is used in that lemma. So, it suffices to prove Lemma 4.1 in the $\mathcal{P}_{\{u\},\emptyset}(N; \alpha, \nu)$ space.

To prove this statement, we keep track of the highest type in the neighbourhood of the vertex u. Let $N^{(0)}(u) = \{u\}$, $\theta_r^{(0)} = \theta_\ell^{(0)} = 0$. For $i \geq 0$, define $N^{(i)}(u)$ as the neighbours of vertices in $N^{(i-1)}(u)$ that are in clockwise direction of u and have relative angle greater than $\theta_\ell^{(i-1)}$ with u or that are in anticlockwise direction of u and have relative angle with u greater than $\theta_r^{(i-1)}$. Define $\theta_r^{(i)}$ as the maximum relative angle between u and any vertex in $N^{(i)}(u)$ that is in anticlockwise direction of u, setting it to $\theta_r^{(i-1)}$ if there is no such vertex. Similarly, define $\theta_\ell^{(i)}$ as the maximum relative angle between u and any vertex in $N^{(i)}(u)$ that is in clockwise direction of u, setting it to $\theta_\ell^{(i-1)}$ if there is no such vertex. This is the simultaneous breadth exploration process that will be defined in more detail in the next section.

Note that any vertex in $N^{(i)}(u)$ has graph distance i to u, but not every vertex of distance i is in $N^{(i)}(u)$. However, we claim that the process cannot leave a vertex that has type larger than the maximum type of any vertex in $N_i(u) := \bigcup_{j=0}^i N^{(j)}(u)$ and is within the sectors exposed undiscovered. For the sake of contradiction, assume that v is a vertex whose type is larger than

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the types of all vertices discovered in $N_i(u)$, but its angle with u satisfies $\theta_r^{(k-1)} < \vartheta_{u,v} \leq \theta_r^{(k)}$, for some $1 \leq k \leq i$. Then there are two vertices $v_{k-1} \in N^{(k-1)}(u)$ and $v_k \in N^{(k)}$ such that v is between them; that is, $\vartheta_{v,v_{k-1}} < 0 \leq \vartheta_{v,v_k}$. But the following holds (the second part will be used in the next section).

Claim 4.2. *Consider three vertices z, y and w, on $\mathcal{D}_R$ (in the hyperbolic plane with curvature -1), such that $d_H(z, w) < R$ and w is at the anticlockwise direction of z whereas y is between z and w. If $t_y > t_w$, then $d_H(y, z) < R$. Also, if $t_y > t_z$, then $d_H(y, w) < R$.*

Proof of Claim 4.2. This is the case as the point y' of type equal to that of y with $\theta_{y'w} = 0$ is still at distance less than R from z. If we move this clockwise towards z, the distance will remain smaller than R, as w will be at the anticlockwise side of y'. An analogous argument shows the second statement. $\square$

The first part of the above claim with v_{k-1}, v, v_k playing the role of z, y, w implies that v is adjacent to v_{k-1} and therefore should have been discovered and become a member of $N^{(k)}(u)$.

The above claim has also the following consequence. Denote by $t^{(i-1)}$ the maximum type of a vertex in $N_{i-1}(u)$. As every vertex in $N^{(i-1)}(u)$ is further in the anticlockwise or in the clockwise direction, in terms of relative angle from u, than all the vertices in $N_{i-2}(u)$, all vertices in $N^{(i)}(u)$ are either within (hyperbolic) distance R and in the clockwise direction of the point $p_\ell^{(i-1)}$ of type $t^{(i-1)}$ and of clockwise relative angle $\theta_\ell^{(i-1)}$ to u, or within (hyperbolic) distance R and in the anticlockwise direction of the point $p_r^{(i-1)}$ of type $t^{(i-1)}$ and of clockwise relative angle $\theta_r^{(i-1)}$ to u. Thus the highest type of a vertex in $N^{(i)}(u)$ is stochastically dominated from above by the highest type among all vertices that have hyperbolic distance at most R from a certain point of type $t^{(i-1)}$ (namely $p_r^{(i-1)}$ or $p_\ell^{(i-1)}$). Due to this we can bound the distribution function of $t^{(i)}$ from below using Fact 2.3. Let $\hat{t}^{(i)} := (1 + \delta + \zeta)^i t_u$, for any integer $i \geq 0$.

Claim 4.3. *For $i \geq 1$, assuming that $\hat{t}^{(i-1)} < \frac{R/2 - 2 \log \log R}{1 + \delta + \zeta}$, we have*

$$\mathbb{P}(t^{(i)} < (1+\delta+\zeta)\hat{t}^{(i-1)} \mid t^{(i-1)} < \hat{t}^{(i-1)}) \geq \exp\left(-\frac{2\nu}{(\alpha - 1/2)\pi} e^{-(\alpha-1/2)\zeta\hat{t}^{(i-1)}}\right).$$

Proof. By the assumption of the claim, if $t^{(i-1)} < \hat{t}^{(i-1)}$, then $t^{(i-1)} < (1/(1+\delta+\zeta))(R/2-2\log\log R) < (2\alpha-1)R/2$. Lemma 2.5 works for types t such that $t+t^{(i-1)} < R-c_0$ for a given constant c_0, so $t < R-(1/(1+\delta))R/2$ will do. Recall that $1/(1+\delta) = 2\alpha - 1$, so $t < R(3/2 - \alpha)$ is sufficient. But $3/2-\alpha > 1/(2\alpha)$, and so if we take $\hat{t} = R/(2\alpha)+\omega(N)$, for some sufficiently slowly growing function $\omega(N)$, we are able to use Lemma 2.5 for points of type at most $\hat{t}$. The first part of Corollary 2.7 implies that the expected number of vertices of type at least $\hat{t}$ in $\mathcal{G}_{\{u\},\emptyset}(N;\alpha,\nu)$ is $o(1)$.

As discussed above, the event where $t^{(i)} \le (1+\delta+\zeta)\hat{t}^{(i-1)}$ has no smaller probability than the event that a vertex of type $\hat{t}^{(i-1)}$ has no neighbour of type at least $\hat{t}^{(i)}$. Thus by Fact 2.3 and Lemma 2.5, for $\varepsilon > 0$ small enough so that $(1+2\varepsilon)\alpha < 1$ we have

$$\mathbb{P}\left(t^{(i)} < (1+\delta+\zeta)\hat{t}^{(i-1)} \mid t^{(i-1)} < \hat{t}^{(i-1)}\right)$$

$$\ge \exp\left(-N\int_{(1+\delta+\zeta)\hat{t}^{(i-1)}}^{\hat{t}} \frac{4(1+\varepsilon)}{2\pi}e^{1/2(t+\hat{t}^{(i-1)}-R)}\alpha e^{-\alpha t}dt + o(1)\right)$$

$$\ge \exp\left(-\frac{2(1+2\varepsilon)\alpha\nu}{\pi}e^{\frac{\hat{t}^{(i-1)}}{2}}\int_{(1+\delta+\zeta)\hat{t}^{(i-1)}}^{\infty} e^{(1/2-\alpha)t}dt\right)$$

$$\ge \exp\left(-\frac{2(1+2\varepsilon)\alpha\nu}{\pi}e^{\frac{\hat{t}^{(i-1)}}{2}}\frac{1}{\alpha-1/2}e^{(1/2-\alpha)(1+\delta+\zeta)\hat{t}^{(i-1)}}\right)$$

$$\ge \exp\left(-\frac{2(1+2\varepsilon)\alpha\nu}{\pi}e^{\frac{\hat{t}^{(i-1)}}{2}}\frac{1}{\alpha-1/2}e^{(-1/2+(1/2-\alpha)\zeta)\hat{t}^{(i-1)}}\right)$$

$$\ge \exp\left(-\frac{2\nu}{(\alpha-1/2)\pi}e^{-(\alpha-1/2)\zeta\hat{t}^{(i-1)}}\right),$$

as $(\alpha - 1/2)(1 + \delta) = 1/2$. $\qquad\square$

We repeatedly apply this bound to bound the distance from the core. Assume that $t_u = \log\log R$. Denote by $\mathcal{U}$ the event that if we explore as above the neighbours u for every $i < (\tau - \zeta^{1/2})\log R$ we have $t^{(i)} < \hat{t}^{(i)}$.

Claim 4.4. *Assume that $t_u = \log\log R$. For $\zeta > 0$ small enough (depending on α), the event $\mathcal{U}$ has probability $1-o(1)$ and after the steps are completed the maximum type reached is less than $R/2 - 2\log\log R$, if N is sufficiently large.*

Proof. On this event, after executing the $(\tau - \zeta^{1/2})\log R$ steps we have

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reached type less than

$$(1 + \delta + \zeta)^{(\tau - \zeta^{1/2}) \log R} \log \log R = e^{\log(1+\delta+\zeta)(\tau-\zeta) \log R} \log \log R$$
$$\leq R^{(\log(1+\delta)+\zeta)(\tau-\zeta^{1/2})} \log \log R$$
$$= R^{(\tau^{-1}+\zeta)(\tau-\zeta^{1/2})} \log \log R$$
$$= R^{1-\tau^{-1}\zeta^{1/2}+\tau\zeta-\zeta^{3/2}} \log \log R = o(R/2 - 2 \log \log R).$$

Moreover, we are able to apply Claim 4.3 repeatedly for this number of steps and deduce that $\mathcal{U}$ has probability

$$\mathbb{P}(\mathcal{U}) \geq \prod_{i=0}^{(\tau-\zeta^{1/2}) \log R} \exp\left(-\frac{2\nu}{(\alpha-1/2)\pi} e^{-(\alpha-1/2)\zeta(1+\delta+\zeta)^i \log \log R}\right)$$

$$\geq \prod_{i=0}^{(\tau-\zeta^{1/2}) \log R} \left(1 - \frac{2\nu}{(\alpha-1/2)\pi} e^{-(\alpha-1/2)\zeta(1+\delta+\zeta)^i \log \log R}\right)$$

$$\geq 1 - \sum_{i=0}^{(\tau-\zeta^{1/2}) \log R} \frac{2\nu}{(\alpha-1/2)\pi} e^{-(\alpha-1/2)\zeta(1+\delta+\zeta)^i \log \log R}$$

$$\geq 1 - \frac{4\nu}{(\alpha-1/2)\pi} e^{-(\alpha-1/2)\zeta \log \log R} = 1 - o(1).$$

$\square$

Proof of Lemma 4.1. Fact 2.4 implies that increasing the type of a vertex will keep all edges intact, so any path will stay a path if we increase the type of one of its vertices. Thus by a simple coupling argument we have that $\mathbb{P}(d(u, core) \leq d|t_u) \leq \mathbb{P}(d(u, core) \leq d|t'_u)$ for $t_u \leq t'_u$. We can thus assume that $t_u = \log \log R$. By Claim 4.4, a.a.s. executing $(\tau - \zeta^{1/2}) \log R$ steps yields maximum type that is less than $R/2 - 2 \log \log R$, so

$$\mathbb{P}(d(u, core) \leq (\tau - \zeta^{1/2}) \log R) = o(1).$$

$\square$

5   Proof of Theorem 1.2

Here, we consider the case where $\alpha > 1$. In this case, the main result in [3] implies that all components contain at most sublinear number of

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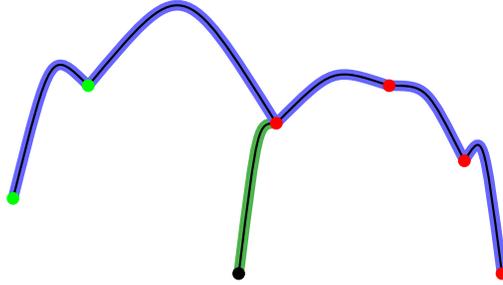
Figure 2: Example of an umbrella.

vertices. More precisely, we show that a.a.s. all components contain at most $N^{1/\alpha}$ vertices (up to a poly-logarithmic factor). In fact, there are many components of polynomial size (as there are many vertices of polynomial degree which do not belong to the same component).

To prove Theorem 1.2, for any given vertex we explore a path that in some sense traverses its component. We show that almost all vertices are close to such a *spanning path*, which itself is short. This results in short distances for most pairs of vertices which belong to the same component.

Note that since $\alpha > 1$, a.a.s. there is no component whose convex hull contains the origin. In fact, components are included in a section of the disc spanning $o(1)$ of all angles. Due to this, it creates no ambiguity to talk of clockwise and anticlockwise directions in a component.

Definition 5.1. *We call a path $P = v_1, \ldots, v_\ell$ in a component C a spanning path of C if v_1 is the vertex of C that is farthest in clockwise and v_ℓ is the vertex of C that is farthest in anticlockwise direction.*

An umbrella *U with root vertex v is a spanning path P of the component of v together with a path connecting v to P. The* size *of the umbrella U is the maximum among the distances of v from the two endpoints of the associated spanning path.*

Note that any vertex in C that is above a spanning path P of C is directly connected to one of the vertices of P by Fact 2.4. Since there is no restriction on the length of the paths, if v is on some spanning path P, then P is an umbrella with root v.

The following follows immediately as the vertices of a component that are to the farthest in clockwise and anticlockwise direction are always in a

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spanning path:

Corollary 5.2. *If P and P' are spanning paths of the same component, then $P \cap P' \neq \emptyset$.*

This fact allows us to do the following: Given any pair of vertices u and v in the same component, construct a u-v-path by traversing the umbrella U_u of u until the first vertex z that is on the umbrella U_v of v is reached. Then $uU_u z U_v v$ is a path connecting u and v. Thus the following lemma is key to the proof of Theorem 1.2.

Lemma 5.3. *Let $\varepsilon > 0$. For a vertex v of $\mathcal{G}(N; \alpha, \nu)$, a.a.s. there is an umbrella for v of size at most $\log^{1+\varepsilon} \log N$.*

For the proof of this lemma we define the *simultaneous breadth exploration process* starting at a vertex v similar to the one that we introduced in [3]. Here, we keep track of two sets of vertices V_ℓ and V_r, which both start out as $\{v\}$. Roughly speaking, we update the two sets adding the neighbours of the current sets that are located in the clockwise and anticlockwise direction from the "current" vertices, respectively. If there are no neighbours that are farther in the clockwise direction of V_r and no neighbours that are farther in the anticlockwise direction of V_ℓ, then the process stops. We define the process starting at vertex v as the following steps:

(i) Let $V_\ell^{(0)} = V_r^{(0)} = \{v\}$ and let $i := 1$.

(ii) Let $V_\ell^{(i)}$ be the set of vertices not in $V_\ell^{(i-1)} \cup V_r^{(i-1)}$ that are neighbours of some vertex in $V_\ell^{(i-1)} \cup V_r^{(i-1)}$ and are in the clockwise direction of every vertex in $\bigcup_{j=0}^{i-1} \{V_\ell^{(j)} \cup V_r^{(j)}\}$. We define similarly the set $V_r^{(i)}$ as the set of vertices not in $V_\ell^{(i-1)} \cup V_r^{(i-1)}$ that are neighbours of some vertex in $V_\ell^{(i-1)} \cup V_r^{(i-1)}$ and are in the anticlockwise direction of every vertex in $\bigcup_{j=0}^{i-1} \{V_\ell^{(j)} \cup V_r^{(j)}\}$.

(iii) If $V_\ell^{(i)} = \emptyset = V_r^{(i)}$, then stop. Otherwise, let $i := i + 1$ and go to step (ii).

We call a repetition of steps (ii) and (iii) a *round*. To prove Lemma 5.3, we show that this process yields an umbrella and bound the number of steps needed until completion.

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Lemma 5.4. *If the simultaneous breadth exploration process starting at a vertex v stops after k rounds, then there is an umbrella for v that has size at most k.*

Proof. Let $C(v)$ denote the connected component that v belongs to. Let $V_i' = \bigcup_{j=0}^{i}\{V_\ell^{(j)} \cup V_r^{(j)}\}$, that is, the set of vertices discovered up to round i. We denote by v_ℓ' the vertex in V_i' with the largest relative angle with v in the clockwise direction. We let $\theta_\ell^{(i)}$ be this angle and let $t_\ell^{(i)}$ be the type of this vertex. Similarly, let v_r' be the vertex of V_i' that is the farthest in the anticlockwise direction, and let $\theta_r^{(i)}$ and $t_r^{(i)}$ denote its angle and type. Note that there is an edge between some vertex v_ℓ in V_{i-1}' to the vertex v_ℓ' in $V_\ell^{(i)}$ and also an edge between some vertex $v_r \in V_{i-1}'$ and the vertex v_r'.

We now claim that if the process stops at round k, then the vertices $\hat{v}_r$ and $\hat{v}_\ell$ that are the farthest to the anticlockwise and clockwise direction of $C(v)$ belong to V_{k-1}'. Note that $V_\ell^{(k)} = V_r^{(k)} = \emptyset$, so $V_{k-1}' = V_k'$. Assume this is not the case, so without loss of generality $\hat{v}_r \notin V_{k-1}'$. As v and $\hat{v}_r$ are in the same component, there is a path P from v to $\hat{v}_r$. Let w be the first vertex on P that is outside the range of angles from $\theta_\ell^{(k-1)}$ to $\theta_r^{(k-1)}$. Since $\hat{v}_r$ is the vertex that is farthest in the anticlockwise direction and $\hat{v}_r \notin V_k'$ this vertex must exist. Let u be the predecessor of w on P. We cannot have $u \in V_k'$ as otherwise w, being farther in the clockwise or anticlockwise direction than any other vertex in V_k', must also be in V_k' by the choice made in step (ii). There exists an $i < k$ and two *adjacent* vertices x and y such that x has been discovered at round $i-1$ and y has been discovered at round i and u is between x and y. Now, if $t_u \geq t_y$, then by Claim 4.2 (x, u, y playing the role of w, y, z) it follows that u is adjacent to x as well. If $t_u < t_y$, then again Claim 4.2 implies that y is adjacent to w. Hence, in either case w would have been discovered by round $i+1$, whereby $w \in V_r^{(i+1)} \cup V_\ell^{i+1} \subseteq V_k'$; a contradiction.

So both $\hat{v}_\ell$ and $\hat{v}_r$ are in V_k'. Note that every vertex in $V_\ell^{(i)} \cup V_r^{(i)}$ has a neighbour in $V_\ell^{(i-1)} \cup V_r^{(i-1)}$, so we can find a paths P_ℓ and P_r of length at most k from $\hat{v}_\ell$ to v and from $\hat{v}_r$ to v, respectively. Together, possibly deleting redundant subpaths in $v_\ell P_\ell v P_r v_r$, we have an umbrella for v of size at most k. □

We are now ready to prove Lemma 5.3

Proof of Lemma 5.3. We aim to bound the number of rounds it takes for

the simultaneous breadth exploration process started at some vertex v to stop. By Corollary 2.7, it would be sufficient to consider a variation of the simultaneous breadth exploration process where we expose only those vertices that have type at most $R/(2\alpha) + \omega(N)$, for some slowly growing function $\omega(N) \to \infty$. We will use the same notation for the parameters of the process as in the unmodified process.

Let T denote the stopping time of this process. Without loss of generality, assume that $V_\ell^{(i)}, V_r^{(i)} \neq \emptyset$ for $i = 1, \ldots, T-1$. Define V_i', $\theta_\ell^{(i)}$ and $\theta_r^{(i)}$ as in the previous proof (but for the modified process). Unlike the last proof, let $t_\ell^{(i)}$ and $t_r^{(i)}$ be the maximum types of vertices in $V_\ell^{(i)}$ and $V_r^{(i)}$, respectively, and they are set to 0, if the corresponding set contains no vertices. Let $t_i = \max\{t_\ell^{(i)}, t_r^{(i)}\}$. Let $p_\ell^{(i)}$ be the point of type t_i and angle $\theta_\ell^{(i)}$ in the clockwise direction from v. Similarly, let $p_r^{(i)}$ be the point of type t_i and angle $\theta_r^{(i)}$ in the anticlockwise direction from v.

Claim 5.5. *We have $V_\ell^{(i+1)} \subset T_\varepsilon^+(p_\ell^{(i)})$ and $V_r^{(i+1)} \subset T_\varepsilon^+(p_r^{(i)})$.*

Proof of Claim 5.5. Let p be a point that is within hyperbolic distance R from $u \in V_\ell^{(i)} \cup V_r^{(i)}$ and satisfies $\vartheta_{p,v} > \theta_\ell^{(i)}$. Let u' be the point of type $t_\ell^{(i)}$, which has $\theta_{u,u'} = 0$.

Note that $\vartheta_{p,p_\ell^{(i)}} \leq \vartheta_{p,u}$. Since $p \in T_\varepsilon^+(u)$, we have $\vartheta_{p,u} \leq 2(1 + \varepsilon)\frac{\nu}{N}e^{\frac{t_p+t_u}{2}}$. As $t_u \leq t_{u'} = t_\ell^{(i)}$, it follows that $\vartheta_{p,u} \leq 2(1 + \varepsilon)\frac{\nu}{N}e^{\frac{t_p+t_\ell^{(i)}}{2}}$. In other words, $p \in T_\varepsilon^+(p_\ell^{(i)})$. Thereby, $V_\ell^{(i+1)} \subset T_\varepsilon^+(p_\ell^{(i)})$.

The proof that $V_r^{(i+1)} \subset T_\varepsilon^+(p_r^{(i)})$ is analogous. $\qquad\square$

The above claim implies that the highest type of a vertex in $V_\ell^{(i+1)}$, which we denoted by $t_\ell^{(i)}$, is stochastically dominated by the highest type among the vertices in $\left\{p \in T_\varepsilon^+(p_\ell^{(i)}) \ : \ \vartheta_{p,p_\ell^{(i)}} > 0, \ t_p < R/(2\alpha) + \omega(N)\right\}$. Similarly, the highest type of a vertex in $V_r^{(i+1)}$, which we denoted by $t_r^{(i)}$ is stochastically dominated by the highest type among the vertices in $\left\{p \in T_\varepsilon^+(p_r^{(i)}) \ : \ \vartheta_{p,p_r^{(i)}} < 0, \ t_p < R/(2\alpha) + \omega(N)\right\}$. Let $T_\ell(p_\ell^{(i)})$ and $T_r(p_r^{(i)})$ denote these two sets.

Thus, t_{i+1} is stochastically bounded from above by the largest type in $T_\ell(p_\ell^{(i)}) \cup T_r(p_r^{(i)})$. In turn, this is stochastically bounded from above by the maximum type of a vertex in $T_\ell(p^{(i)}) \cup T_r(p^{(i)})$ for a point $p^{(i)}$ of type $t_i = \max\{t_\ell^{(i)}, t_r^{(i)}\}$. We shall proceed with the estimation of the cdf of the latter random variable.

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Observe first that Claim 5.5 implies that for all $0 < i \leq T$ we have $V_i' \subset \bigcup_{j=0}^{i-1}\{T_\varepsilon^+(p_\ell^{(j)}) \cup T_\varepsilon^+(p_r^{(j)})\}$, assuming that $p_\ell^{(0)}, p_r^{(0)}$ are both set to the point of $\mathcal{D}_R$ where v is located. Let $\mathcal{N}_i$ be the set of vertices that belong to V_i'. For a vertex $u \in V_N \setminus V_i'$, the distribution on $\mathcal{D}_R$ is uniform (within the plane of curvature $-\alpha^2$) on the subset of $\mathcal{D}_R$ that excludes the union of the balls of radius R around each vertex in V_i'. Recall that $\text{Area}_\alpha(\cdot)$ denotes the area of a measurable subset of $\mathcal{D}_R$ on the hyperbolic plane of curvature $-\alpha^2$. By Lemma 2.5 and the above observation, the area of the latter is at most $\sum_{j=0}^{i-1} \text{Area}_\alpha(T_\varepsilon^+(p_\ell^{(j)}) \cup T_\varepsilon^+(p_r^{(j)}))$. But for each j, the angle that is spanned by $T_\varepsilon^+(p_\ell^{(j)}) \cup T_\varepsilon^+(p_r^{(j)})$ is proportional to $e^{R/(2\alpha)-R+\omega(N)} = o(1)$. Thus, if $i < R$, then we have $\sum_{j=0}^{i-1} \text{Area}_\alpha(T_\varepsilon^+(p_\ell^{(j)}) \cup T_\varepsilon^+(p_r^{(j)})) = o(\text{Area}_\alpha(\mathcal{D}_R))$.

Using this, we conclude that the conditional probability that a vertex $u \in V_N \setminus \mathcal{N}_i$ belongs to $T_\varepsilon^+(p^{(i)})$ and has type t_u that satisfies $t \leq t_u < R/(2\alpha) + \omega(N)$ is at most

$$\int_t^{\frac{R}{2\alpha}+\omega(N)} \frac{4(1+\varepsilon)}{2\pi} e^{\frac{t_i+t'-R}{2}} \frac{\alpha \sinh(\alpha(R-t'))}{\cosh(\alpha R)(1-o(1))} dt'$$

$$\leq \frac{2\alpha(1+2\varepsilon)}{\pi} e^{\frac{t_i-R}{2}} \int_t^{\frac{R}{2\alpha}+\omega(N)} e^{t'/2} \frac{e^{\alpha(R-t')}}{2\cosh(\alpha R)(1-o(1))} dt'$$

$$\leq \frac{2\alpha(1+3\varepsilon)}{\pi} e^{\frac{t_i-R}{2}} \int_t^{\frac{R}{2\alpha}+\omega(N)} e^{\left(\frac{1}{2}-\alpha\right)t'} dt'$$

$$= \frac{2\alpha\nu(1+3\varepsilon)}{\pi} \frac{e^{t_i/2}}{N} \int_t^{\frac{R}{2\alpha}+\omega(N)} e^{\left(\frac{1}{2}-\alpha\right)t'} dt' < \frac{4\alpha\nu(1+3\varepsilon)}{\pi(2\alpha-1)} \frac{e^{t_i/2}}{N} e^{\left(\frac{1}{2}-\alpha\right)t},$$

for N sufficiently large. Therefrom, the conditional probability that *none* of the vertices in $V_N \setminus \mathcal{N}_i$ satisfies this is at least

$$\left(1 - \frac{4\alpha\nu(1+3\varepsilon)}{\pi(2\alpha-1)} \frac{e^{t_i/2}}{N} e^{\left(\frac{1}{2}-\alpha\right)t}\right)^{|V_N \setminus \mathcal{N}_i|} > \left(1 - \frac{4\alpha\nu(1+3\varepsilon)}{\pi(2\alpha-1)} \frac{e^{t_i}/2}{N} e^{\left(\frac{1}{2}-\alpha\right)t}\right)^N$$

$$> \exp\left(-D_{\alpha,\nu,\varepsilon} e^{\frac{t_i}{2}-(\alpha-1/2)t}\right),$$

(5.1)

for some $D_{\alpha,\nu,\varepsilon} > 0$ and any N sufficiently large.

Therefore, for $i < R$ the random variable $\max\{t_\ell^{(i+1)}, t_r^{(i+1)}\}$ conditional on the history of the process up to step i is stochastically dominated by a random variable that follows the Gumbel distribution. The expectation of

the latter is

$$\frac{t_i + 2\ln(2D_{\alpha,\nu,\varepsilon})}{2\alpha - 1} + \frac{2\gamma}{2\alpha - 1},$$

where γ is Euler's constant. Therefore, the following inequality holds:

$$\mathbb{E}\left[t_{i+1}|\mathcal{F}_i\right] \leq \frac{t_i + 2\ln(2D_{\alpha,\nu,\varepsilon})}{2\alpha - 1} + \frac{2\gamma}{2\alpha - 1},$$

where $\mathcal{F}_i$ denotes the sub-σ-algebra generated by the process up to step i. There exists a constant $U_{\alpha,\nu,\varepsilon} > 0$ such that when $t_i > U_{\alpha,\nu,\varepsilon}$, we have

$$\mathbb{E}\left[t_{i+1}|\mathcal{F}_i\right] \leq \frac{t_i + 2\ln(2D_{\alpha,\nu,\varepsilon})}{2\alpha - 1} + \frac{2\gamma}{2\alpha - 1} < \frac{\alpha}{2\alpha - 1}t_i =: \lambda_\alpha t_i < t_i. \quad (5.2)$$

On the other hand, (5.1) implies that if $t_i \leq U_{\alpha,\nu,\varepsilon}$, then

$$\mathbb{P}(t_{i+1} = 0) \geq p > 0, \quad (5.3)$$

for some positive constant p.

With these tools, we can bound the stopping time T of the process. Let $[T_1^{(s)}, T_2^{(s)} \wedge R]$ denote the sth interval of indices in which the process stays above $U_{\alpha,\nu,\varepsilon}$. By (5.2), for $T_1^{(s)} < i \leq T_2^{(s)} \wedge R$ the process (t_i) is a supermartingale with decay rate at most λ_α.

Claim 5.6. *For any $\varepsilon' > 0$*

$$\mathbb{P}((T_2^{(s)} \wedge R) - T_1^{(s)} \geq \log_{1/\lambda_\alpha}^{1+\varepsilon'} R) = o(1).$$

Proof of Claim 5.6. Let $S := \log_{1/\lambda_\alpha}^{1+\varepsilon'} R$ and let $T^{(s)} := T_2^{(s)} \wedge R$. Note that $\mathbb{E}\left[t_{i \wedge T^{(s)}} \mid \mathcal{F}_{T_1^{(s)}}\right] \leq \lambda_\alpha^{i \wedge T^{(s)} - T_1^{(s)}} t_{T_1^{(s)}} \leq \lambda_\alpha^{i \wedge T^{(s)} - T_1^{(s)}} R$. Let A be the event $\{T^{(s)} > S + T_1^{(s)}\}$. If $\omega \in A$, then $\lambda_\alpha^{(S+T_1^{(s)}(\omega)) \wedge T^{(s)}(\omega) - T_1^{(s)}} t_{T_1^{(s)}}(\omega) < \lambda_\alpha^S R = o(1)$. By the definition of the conditional expectation, we deduce that $\mathbb{E}\left[t_{(S+T_1^{(s)}) \wedge T^{(s)}}\mathbf{1}_A\right] = o(1)$ and since $\mathbb{E}\left[t_{(S+T_1^{(s)}) \wedge T^{(s)}}\mathbf{1}_A\right] > U_{\alpha,\nu,\varepsilon}\mathbb{P}(A)$, we finally deduce that $\mathbb{P}(A) = o(1)$. $\qquad \square$

Now, the length of the (discrete) interval $(T_2^{(s)}, T_1^{(s+1)} \wedge T \wedge R)$ is stochastically bounded from above by a geometric random variable that has parameter at least p.

We call the union of these intervals an *epoch*, that is, we call an epoch

the interval $[T_1^{(s)}, T_1^{(s+1)} \wedge T \wedge R)$, for some $s > 0$. By the above claim, for any $\varepsilon' > 0$, with probability $1 - o(1)$, we have $(T_2^{(s)} \wedge R) - T_1^{(s)} \leq \log_{1/\lambda_\alpha}^{1+\varepsilon'} R$. Additionally, the stochastic upper bound on the interval $(T_2^{(s)}, T_1^{(s+1)})$ implies that this is at most $\log_{1/\lambda_\alpha}^{\varepsilon'} R$ with probability $1 - o(1)$. Hence, with probability $1 - o(1)$ an epoch lasts for at most $\log_{1/\lambda_\alpha}^{1+2\varepsilon'} R$ steps. Finally, since every epoch has probability at least p to be the final one, it follows that the process hits 0 within $\log_{1/\lambda_\alpha}^{1+3\varepsilon'} R$ steps with probability $1 - o(1)$. In other words, a.a.s. we have $T \leq \log_{1/\lambda_\alpha}^{1+3\varepsilon'} R$.

$\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad\qquad$ $\square$

Using the previous lemmas we prove Theorem 1.2.

Proof of Theorem 1.2. Let $0 < \varepsilon' < \varepsilon$. Let V' be the set of vertices in $\mathcal{G}(N; \alpha, \nu)$ that have an umbrella of size at most $\log^{1+\varepsilon'} \log N$. By Lemma 5.3 we have $|V'| = (1 - o(1))N$ a.a.s. For any $u, v \in V'$, if they are in the same component, by Corollary 5.2 the umbrellas are not disjoint. Thus there is a u-v-path of length at most $|U_u| + |U_v| \leq 2 \log^{1+\varepsilon'} \log N < \log^{1+\varepsilon} \log N$ for N large enough. $\qquad\qquad\qquad$ $\square$

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